Representation of Partial Traces

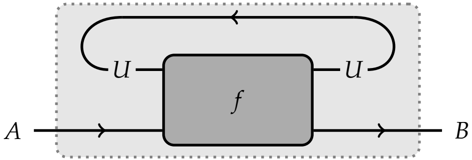
MFPS XXXI, Nijmegen, Holland

Marc Bagnol — *University of Ottawa*

Traces in symmetric monoidal categories

SMC: a category with an associative bifunctor $\otimes$, a unit object $\mathbf{1}$ and a family of isomorphisms $\sigma_{A,B} : A \otimes B \to B \otimes A$.

Trace (A. Joyal, R. Street, D. Verity): operation turning $f : A \otimes U \to B \otimes U$ into $\mathbf{Tr}^U[f] : A \to B$.



Understood as a *feedback along* U.

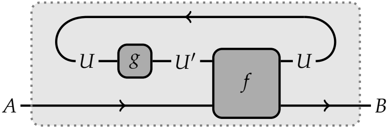
Ubiquitous structure in mathematics: linear algebra, topology, knot theory, computer science, proof theory…

Partial traces

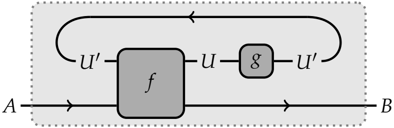
P. Scott & E. Haghverdi: axiomatization of partially-defined trace, capturing the idea of (partially defined) categorical feedback.

One example of partial traces axiom: sliding

$$\mathbf{Tr}^{U}\left[f(\mathrm{Id}_A \otimes g)\right] \leftrightharpoons \mathbf{Tr}^{U'}\left[(\mathrm{Id}_B \otimes g)f\right]$$



$\leftrightharpoons$



Partial traces and sub-categories

A straightforward way to build partial traces:

- Consider a totally traced category $\mathcal{D}$.
- Take any sub-SMC $\mathcal{C} \subseteq \mathcal{D}$.
- If $f : A \otimes U \to B \otimes U$ is in $\mathcal{C}$, it always has a trace $\mathbf{Tr}^U[f]$ in $\mathcal{D}$.
  ($\mathbf{Tr}^U[f]$ may or may not end up in $\mathcal{C}$)

  Define a partial trace $\widehat{\mathbf{Tr}}$ on $\mathcal{C}$ as:

$$\text{if } \mathbf{Tr}^U[f] \in \mathcal{C} \text{ then } \widehat{\mathbf{Tr}}^U[f] = \mathbf{Tr}^U[f], \text{ undefined otherwise}$$

Does any partial trace arise this way?

O. Malherbe, P. Scott, P. Selinger: representation theorem.

Allows intuitive diagrammatic reasoning also in the partially-defined case.

The representation theorem

More precisely: any partially traced category embeds in a totally traced one. We also have a universal property (free construction):

$$\begin{tikzcd} \mathcal{C} \arrow[r, "E_{\mathcal{C}}"] \arrow[dr, "F"'] & \mathbf{T}(\mathcal{C}) \arrow[d, dotted, "G"] \\ & \mathcal{D} \end{tikzcd}$$

(where $\mathcal{C}$ is partially traced, $\mathbf{T}(\mathcal{C})$ is the totally traced category in which it embeds, $\mathcal{D}$ is any other totally traced category, with F a traced functor from $\mathcal{C}$ to $\mathcal{D}$)

Original proof: intermediate partial version of the $\mathbf{Int}(\cdot)$ construction and "paracategories".

Contribution: a more direct and simplified proof.

The proof (I): the dialect construction

A generic construction $\mathbf{D}(\mathcal{C})$ on any monoidal category $\mathcal{C}$.

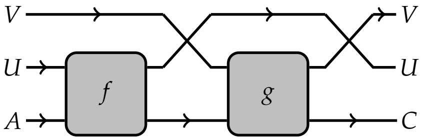
Basic idea: add a "state space" to morphisms.

A morphism from A to B in $\mathbf{D}(\mathcal{C})$ is a pair (f, U) with

- U an object of $\mathcal{C}$.
- $f : A \otimes U \to B \otimes U$ a morphism of $\mathcal{C}$.

When composing (f, U) and (g, V) the state spaces do not interact:

$(g, V) \circ (f, U) =$



(for the exerted eye: notice the similarity with composition in $\mathbf{Int}(\cdot)$ *categories)*

The proof (II): hiding and congruences

Hiding: given a partially traced $\mathcal{C}$ we can look at $\mathbf{D}(\mathcal{C})$ and define a *hiding* operation turning $(f, V) : A \otimes U \to B \otimes U$ into

$$\mathsf{H}^U[f, V] = (f, U \otimes V) : A \to B$$

$\mathsf{H}[\cdot]$ behaves a lot like a (total) trace.

Congruences: enforce the missing equations, for instance

$$(f, U \otimes V) \approx (\mathsf{Tr}^V[f], U) \text{ when } \mathsf{Tr}^V[f] \text{ is defined}$$

by considering the equivalence relation generated and setting $\mathbf{T}(\mathcal{C}) = \mathbf{D}(\mathcal{C})/\approx$ in which $\mathsf{H}[\cdot]$ induces a total trace, encompassing the original partial trace of $\mathcal{C}$.

The proof (III): a sketch

We can embed $\mathcal{C}$ in $\mathbf{T}(\mathcal{C})$ by setting $E_{\mathcal{C}}(f) = (f, \mathbf{1})$.

Is it really an embedding? We check that $(f, \mathbf{1}) \approx (g, \mathbf{1})$ implies $f = g$.

Because $\approx$ is freely generated, we can do it by induction on chains of elementary equivalences.

Universal property: we can close the diagram

$$\begin{tikzcd} \mathcal{C} \arrow[r, "E_{\mathcal{C}}"] \arrow[dr, "F"'] & \mathbf{T}(\mathcal{C}) \arrow[d, dotted, "G"] \\ & \mathcal{D} \end{tikzcd}$$

by setting $G(f, U) = \mathbf{Tr}^{FU}[Ff]$.

(well defined because $(f, U) \approx (g, V)$ implies $\mathbf{Tr}^{FU}(Ff) = \mathbf{Tr}^{FV}(Fg)$)

… Thank you for your attention !