C Variable Renaming Theorem for $BI^{\mu\nu}$

Theorem 5. $\begin{array}{c|c} -P \ is \ v\text{-}closed \\ If \ -z \not\in Var(P) \\ -y \notin FV(P) \end{array} then \ P \equiv P\{[z/y]\} \ . \end{array}$

Lemma 1. $\llbracket E\{E'/x\} \rrbracket^s = \llbracket E \rrbracket^{[s|x \to \llbracket E' \rrbracket^s]}$ if $\llbracket E' \rrbracket^s$ exists

Proof. Lemma 1

 $- \begin{bmatrix} x \{E'/x\} \end{bmatrix}^s = \begin{bmatrix} E' \end{bmatrix}^s = \begin{bmatrix} x \end{bmatrix}^{[s|x \to [E']]^s]} \\ - \begin{bmatrix} y \{E'/x\} \end{bmatrix}^s = \begin{bmatrix} y \end{bmatrix}^{[s|x \to [E']]^s]} \\ - \begin{bmatrix} True \{E'/x\} \end{bmatrix}^s = \begin{bmatrix} True \end{bmatrix}^s = true = \begin{bmatrix} True \end{bmatrix}^{[s|x \to [E']]^s]} \\ - \begin{bmatrix} False \{E'/x\} \end{bmatrix}^s = \begin{bmatrix} False \end{bmatrix}^s = false = \begin{bmatrix} False \end{bmatrix}^{[s|x \to [E']]^s]} \\ - \begin{bmatrix} 42 \{E'/x\} \end{bmatrix}^s = \begin{bmatrix} 42 \end{bmatrix}^s = 42 = \begin{bmatrix} 42 \end{bmatrix}^{[s|x \to [E']]^s]} \\ - \begin{bmatrix} (E_1 \ op \ E_2) \{E'/x\} \end{bmatrix}^s = \begin{bmatrix} E_1 \{E'/x\} \ op \ E_2 \{E'/x\} \end{bmatrix}^s = \begin{bmatrix} E_1 \{E'/x\} \end{bmatrix}^s op \ \begin{bmatrix} E_2 \{E'/x\} \end{bmatrix}^s = \\ \begin{bmatrix} E_1 \end{bmatrix}^{[s|x \to [E']]^s]} op \ \begin{bmatrix} E_2 \end{bmatrix}^{[s|x \to [E']]^s]} = \begin{bmatrix} E_1 \ op \ E_2 \end{bmatrix}^{[s|x \to [E']]^s]}$

Lemma 2. $\llbracket E\{z/y\}\rrbracket^s = \llbracket E\rrbracket^{[s|y \to s(z)]}$ if $z \in dom(s)$

Proof. Lemma 2 By Lemma 1. \Box

Lemma 3. $\llbracket E\{E'/x\} \rrbracket^s = \llbracket E \rrbracket^s$ if $\llbracket E' \rrbracket^s$ doesn't exists but $\llbracket E\{E'/x\} \rrbracket^s$ does

Proof. Lemma 3 $[x{E'/x}]^s = [E']^s$ So $x \notin Var(E)$ and we directly have $E{E'/x} = E$

Lemma 4. $\llbracket E\{z/y\} \rrbracket^s = \llbracket E \rrbracket^s$ if $z \notin dom(s)$

Proof. Lemma 4 By Lemma 3. \Box

Let
$$s^{\bullet} \triangleq \begin{bmatrix} [s \mid y \to s(z)] \text{ if } z \in dom(s) \\ s & \text{ if } z \notin dom(s) \end{bmatrix}$$

Let ρ^{\bullet} be $[\forall X_v \in dom(\rho). X_v \to \{s, h \mid s^{\bullet}, h \in \rho(X_v)\}]$

Lemma 5. $[\![E\{z/y\}]\!]^s = [\![E]\!]^{s^{\bullet}}$

Proof. Lemma 5 By Lemma 2 and 4. \Box

Remember that in case P is E = E', $E \mapsto E_1, E_2$, false and emp we have $\forall \rho. \llbracket P \rrbracket_{\rho} = \llbracket P \rrbracket$ since they are v-closed, see Lemma 6.

Remember $P \equiv Q$ iff $\forall \rho$. either $(\llbracket P \rrbracket_{\rho} \text{ and } \llbracket Q \rrbracket_{\rho} \text{ do not exist})$ either $\llbracket P \rrbracket_{\rho} = \llbracket Q \rrbracket_{\rho}$.

Proof (Th. 5). By Th. 6, $\llbracket P\{[z/y]\} \rrbracket_{\rho^{\bullet}} = \{s, h \mid s^{\bullet}, h \in \llbracket P \rrbracket_{\rho}\}$ in case $nodep(z, \rho)$ and $z \notin Var(P)$, then $\llbracket P\{[z/y]\} \rrbracket = \{s, h \mid s^{\bullet}, h \in \llbracket P \rrbracket\}$ if $z \notin Var(P)$.

Which is if $z \notin dom(s)$ then $s, h \in [\![P\{[z/y]\}]\!]$ iffs, $h \in [\![P]]$ and if $z \in dom(s)$ we have $s, h \in [\![P\{[z/y]\}]\!]$ iff $[s \mid y \to s(z)], h \in [\![P]]$.

Then since $y \notin FV(P)$ with the stack extension theorem 4 we have nodep(y, P)and so $s, h \in$ so $[s \mid y \to s(z)], h \in [\![P]\!]$ iffs, $h \in [\![P]\!]$.

We then have $\llbracket P\{[z/y]\} \rrbracket = \llbracket P \rrbracket$ which is what we wanted since P is v-closed (see Lemma 6). \Box

Theorem 6. If $\begin{array}{c} -nodep(z,\rho) \\ -z \notin Var(P) \end{array}$ then $\llbracket P\{[z/y]\} \rrbracket_{\rho^{\bullet}} = \{s,h \mid s^{\bullet},h \in \llbracket P \rrbracket_{\rho}\}$

Proof (Th. 6). THE PROOF IS ONLY MADE IF ALL THE lfp and gpf are for MONOTONIC FUNCTIONS. But this is the case for the wlp and sp formulas and for the example in the paper.

We will prove by structural induction on P: (recall that for $E1 = E2, E \mapsto E_1, E_2, \texttt{false}, \texttt{emp } \forall \rho. \llbracket P \rrbracket_{\rho} = \llbracket P \rrbracket)$ $- [[(E_1 = E_2)\{[z/y]\}]]$ $= [\![(E_1\{z/y\} = E_2\{z/y\}]\!]$ $= \{s, h \mid [E_1\{z/y\}]^s = [E_2\{z/y\}]^s \}$ = $\{s, h \mid [E_1]^{s^{\bullet}} = [E_2]^{s^{\bullet}} \}$ $= \{s, h \mid s^{\bullet}, h \in [\![E_1 = E_2]\!]\}$ $- [[(E \mapsto E_1, E_2)\{[z/y]\}]]$ $= \llbracket E\{z/y\} \mapsto E_1\{z/y\}, E_2\{z/y\} \rrbracket$ $= \{s, h \mid dom(h) = \{ [\![E\{z/y\}]\!]^s \} \text{ and } h([\![E\{z/y\}]\!]^s) = \langle [\![E_1\{z/y\}]\!]^s, [\![E_2\{z/y\}]\!]^s \rangle$ $= \{s, h \mid dom(h) = \{\llbracket E \rrbracket^{s^{\bullet}}\} \text{ and } h(\llbracket E \rrbracket^{s^{\bullet}}) = \langle \llbracket E_1 \rrbracket^{s^{\bullet}}, \llbracket E_2 \rrbracket^{s^{\bullet}} \rangle \}$ $= \{s, h \mid s^{\bullet}, h \in \llbracket E \mapsto E_1, E_2 \rrbracket \}$ - $[false{[z/y]}]$ = [false] $= \emptyset$ $= \{s, h \mid s^{\bullet}, h \in \llbracket \texttt{false} \rrbracket$ - If $\llbracket P \Rightarrow Q \rrbracket_{\rho}$ exists then $\llbracket (P \Rightarrow Q) \{ [z/y] \} \rrbracket_{\rho} \bullet$ $= \llbracket P\{[z/y]\} \Rightarrow Q\{[z/y]\} \rrbracket_{\rho} \bullet$ $= (\top \setminus \llbracket P\{[z/y]\} \rrbracket_{\rho^{\bullet}}) \cup \llbracket Q\{[z/y]\} \rrbracket_{\rho^{\bullet}}$ $= (\top \setminus \{s, h \mid s^{\bullet}, h \in \llbracket P \rrbracket_{\rho}\}) \cup \{s, h \mid s^{\bullet} \in \llbracket Q \rrbracket_{\rho}\} (ind.hyp.)$ $= \{s, h \mid s^{\bullet}, h \in (\top \setminus \llbracket P \rrbracket_{\rho}) \cup \llbracket Q \rrbracket_{\rho}\}$ $= \{s, h \mid s^{\bullet}, h \in \llbracket P \Rightarrow Q \rrbracket_{\rho} \}$ $- [[(\exists x. P)\{[z/y]\}]]_{\rho}$ • when $x \neq y$ $= [\exists x. (P\{[z/y]\})]_{\rho}$ $= \{s, h \mid \exists v [s \mid x \to v], h \in \llbracket P\{[z/y]\} \rrbracket_{\rho^{\bullet}}\}$ $= \{s, h \mid \exists v.[s \mid x \to v], h \in \{s, h \mid s^{\bullet}, h \in [\![P]\!]_{\rho}\}\} \text{ (ind.)}$ $= \{s, h \mid \exists v [s \mid x \to v]^{\bullet}, h \in \llbracket P \rrbracket_{\rho} \}$ $= \{s, h \mid \exists v [s^{\bullet} \mid x \to v], h \in \llbracket P \rrbracket_{\rho} \}$ (since $x \neq y, z$) $= \{s, h \mid s^{\bullet}, h \in \llbracket \exists x. P \rrbracket_{\rho}\}$ $- [[(\exists y. P)\{[z/y]\}]]_{\rho}$ $= [\exists z. (P\{[z/y]\})]_{\rho}$ $= \{s, h \mid \exists v [s \mid z \to v], h \in \llbracket P\{[z/y]\} \rrbracket_{\rho^{\bullet}}\}$ $= \{s, h \mid \exists v [s \mid z \to v], h \in \{s, h \mid s^{\bullet}, h \in \llbracket P \rrbracket_{\rho}\}\} \text{ (ind.)}$ $= \{s, h \mid \exists v [s \mid z \to v]^{\bullet}, h \in \llbracket P \rrbracket_{\rho} \}$ $= \{s, h \mid \exists v.[s \mid z \to v \mid y \to v], h \in \llbracket P \rrbracket_{\rho} \}$ $= \{s, h \mid \exists v [s \mid y \to v \mid z \to v], h \in \llbracket P \rrbracket_{\rho} \}$ (since $z \neq y$) $= \{s, h \mid \exists v [s \mid y \to v], h \in \llbracket P \rrbracket_{\rho} \}$ (from stack extension theorem) $= \{s, h \mid \exists v. [s^{\bullet} \mid y \to v], h \in \llbracket P \rrbracket_{\rho} \}$ $= \{s, h \mid s^{\bullet}, h \in \llbracket \exists y. P \rrbracket_{\rho}\}$

$$\begin{split} &- [\exp\{[z/y]\}] \\ &= [\exp] \\ &= \{s, h \mid s^*, h \in [emp]\} \\ &= \{s, h \mid s^*, h \in [emp]\} \\ &= [(P\{[z/y]) * Q\{[z/y])]_{\rho^*} \\ &= [(P\{[z/y]) * Q\{[z/y])]_{\rho^*} \\ &= \{s, h \mid \exists h_0, h_1, h_0 \downarrow h_1 = h, s^*, h_0 \in [P\{[z/y]\}]_{\rho^*} \text{ and } s, h_1 \in [Q\{[z/y]\}]_{\rho^*}\} \\ &= \{s, h \mid \exists h_0, h_1, h_0 \downarrow h_1 = h, s^*, h_0 \in [P]_{\rho} \text{ and } s^*, h_1 \in [Q]_{\rho}\} \\ &= \{s, h \mid \exists h_0, h_1, h_0 \downarrow h_1 = h, s^*, h_0 \in [P]_{\rho} \text{ and } s^*, h_1 \in [Q\{[z/y]\}]_{\rho^*}\} \\ &= \{s, h \mid \exists h_0, h_1, h_0 \downarrow h_1 = h, s^*, h_0 \in [P]_{\rho} \text{ and } s^*, h_1 \in [Q\{[z/y]\}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P \to Q]_{\rho}\} \\ &= \{[P \to Q]\{[z/y]\}]_{\rho^*} \\ &= \{P, h \mid b^*, h \in P \to Q]_{\rho}\} \\ &= \{[X_v]_{\rho^*} \\ &= \{X_v]_{\rho^*} \\ &= [X_v]_{\rho^*} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [X_w]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [E^*]_{\rho^*}\}, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [E^*]_{\rho^*}\}, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [E^*]_{\rho^*}\}, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [E^*]_{\rho^*}\}, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}]_{\rho^*}\} \\ &= \{s, h \mid s^*, h \in [P_{\rho^*}$$

 $F \triangleq \lambda X. \llbracket P\{[z/y]\} \rrbracket_{[\rho^{\bullet}|X_v \to X]}$ $A \triangleq \mathrm{lfp}_{\emptyset}^{\subseteq} F$ $G \triangleq \lambda X. \llbracket P \rrbracket_{[\rho|X_v \to X]}$ $B \triangleq \operatorname{lfp}_{\emptyset}^{\subseteq} G$ $C \triangleq \{s, h \mid s^{\bullet}, h \in B\}$ We want then A = C. We know that B exists. Notice that $([\rho \mid X_v \to B])^{\bullet} = [\rho^{\bullet} \mid X_v \to C].$ First we prove that $A \subseteq C$, we prove it by proving that C = F(C). Since $B = \llbracket \mu X.P \rrbracket_{[\rho|X_v \to X]}$, by the stack extension theorem, we have nodep(z, B)so $nodep(z, [\rho \mid X_v \to B])$ and we can use the induction. By definition, B = G(B), so $B = \llbracket P \rrbracket_{[\rho|X_v \to B]}$ and then $C = \{s, h \mid s^{\bullet}, h \in \llbracket P \rrbracket_{[\rho|X_v \to B]}\}$ by induction we then have $C = \llbracket P\{z/y\} \rrbracket_{([\rho|X_v \to B])^{\bullet}}$, which is $C = \llbracket P\{z/y\} \rrbracket_{[\rho^{\bullet}|X_v \to C]}$, we then have C = F(C) and so A exists and $A \subseteq C$. Now we want to prove that $C \subseteq A$. Let D be the biggest set such that nodep(z, D) and $\{s, h \mid s^{\bullet}, h \in D\} \subseteq A$. By the stack extension theorem, $nodep(z, \llbracket P \rrbracket_{[\rho|X_v \to D]})$ and we can use the induction. Since we said we are working with monotonic functions, we have $F(\{s, h \mid$ $s^{\bullet}, h \in D\}) \subseteq F(A)$ since A is a fix point we have then $F(\{s, h \mid s^{\bullet}, h \in D\}) \subseteq A$ which is $\llbracket P\{[z/y]\} \rrbracket_{[\rho^{\bullet}|X_v \to \{s,h|s^{\bullet},h\in D\}]} \subseteq A$ by induction we have then $\{s, h \mid s^{\bullet}, h \in \llbracket P \rrbracket_{[\rho \mid X_v \to D]}\} \subseteq A$ which is $\{s, h \mid s^{\bullet}, h \in G(D)\} \subseteq A$ Then by construction of D as the biggest set we have $G(D) \subseteq D$ and then D is a postfixpoint of G and then $B \subseteq D$ and then $\{s, h \mid s^{\bullet}, h \in$ $B\} \subseteq \{s, h \mid s^{\bullet}, h \in D\}$ which is $C \subseteq A$. - If $\llbracket \nu X.P \rrbracket_{\rho}$ exists: $\llbracket (\nu X.P)\{[z/y]\} \rrbracket_{\rho} \bullet$ $= \llbracket \nu X.P\{[z/y]\} \rrbracket_{\rho} \bullet$ $= \operatorname{gfp}_{\emptyset}^{\subseteq} \lambda X. \llbracket P\{[z/y]\} \rrbracket_{[\rho^{\bullet}|X_v \to X]}$ see proof below $= \{s, h \mid s^{\bullet}, h \in \operatorname{gfp}_{\emptyset}^{\subseteq} \lambda X. \llbracket P \rrbracket_{[\rho|X_v \to X]} \}$ $= \{s, h \mid s^{\bullet}, h \in \llbracket \nu X. P \rrbracket_{\rho} \}$ Let $F \triangleq \lambda X. \llbracket P\{[z/y]\} \rrbracket_{[\rho^{\bullet}|X_v \to X]}$ $A \triangleq \mathrm{gfp}_{\emptyset}^{\subseteq} F$ $G \triangleq \lambda X . \llbracket P \rrbracket_{[\rho|X_v \to X]}$ $B \triangleq \mathrm{gfp}_{\emptyset}^{\subseteq} G$ $C \triangleq \{s, h \mid s^{\bullet}, h \in B\}$ We want then A = C. We know that B exists. Notice that $([\rho \mid X_v B])^{\bullet} = [\rho^{\bullet} \mid X_v \to C].$ First we prove that $A \supseteq C$, we prove it by proving that C = F(C).

Since $B = \llbracket \nu X.P \rrbracket_{[\rho|X_v \to X]}$, by the stack extension theorem, we have nodep(z, B) so $nodep(z, \lceil \rho \mid X_v \to B])$ and we can use the induction. By definition, B = G(B), so $B = \llbracket P \rrbracket_{[\rho|X_v \to B]}$ and then $C = \{s, h \mid s^{\bullet}, h \in [\![P]\!]_{[\rho|X_v \to B]}\}$ by induction we then have $C = [\![P\{z/y\}]\!]_{([\rho|X_v \to B])^{\bullet}}$, which is $C = [\![P\{z/y\}]\!]_{[\rho^{\bullet}|X_v \to C]}$, we then have C = F(C) and so A exists and $A \supseteq C$. Now we want to prove that $C \supseteq A$. Let D be the smallest set such that nodep(z, D) and $\{s, h \mid s^{\bullet}, h \in D\} \supseteq A$. By the stack extension theorem, $nodep(z, \llbracket P \rrbracket_{[\rho|X_v \to D]})$ and we can use the induction. Since we said we are working with monotonic functions, we have $F(\{s, h \mid$ $s^{\bullet}, h \in D\}) \supseteq F(A)$ since A is a fix point we have then $F(\{s, h \mid s^{\bullet}, h \in D\}) \supseteq A$ which is $\llbracket P\{\lfloor z/y \rfloor\} \rrbracket_{[\rho \bullet | X_v \to \{s,h|s \bullet,h \in D\}]} \supseteq A$ by induction we have then $\{s, h \mid s^{\bullet}, h \in \llbracket P \rrbracket_{[\rho|X_v \to D]}\} \supseteq A$ which is $\{s, h \mid s^{\bullet}, h \in G(D)\} \supseteq A$ Then by construction of D as the smallest set we have $G(D) \supseteq D$ and then D is a prefixpoint of G and then $B \supseteq D$ and then $\{s, h \mid s^{\bullet}, h \in$ $B\} \supseteq \{s, h \mid s^{\bullet}, h \in D\}$ which is $C \supseteq A$.

Lemma 6. If P is v-closed formula: $\forall \rho. \llbracket P \rrbracket_{\rho} = \llbracket P \rrbracket$.

Proof (Lemma 6). The simple case are direct from the definition of $[\![\cdot]\!]_{\rho}$, the others come by induction. \Box