

G Simplifications on [/]

Theorem 11. $\boxed{\text{If } \begin{array}{l} P \text{ is } v\text{-closed} \\ z \notin \text{Var}(P) \end{array} \text{ then } P[z/y] \equiv P\{[z/y]\} \wedge \text{is}(z)}$

Proof (Th. 11). From the generalized variable renaming theorem Th. 6, $\llbracket P\{[z/y]\} \rrbracket = \{s, h \mid s^\bullet, h \in \llbracket P \rrbracket\}$,
 $\llbracket P\{[z/y]\} \wedge \text{is}(z) \rrbracket = \{s, h \mid [s \mid y \rightarrow s(z)], h \in \llbracket P \rrbracket\}$,
which is $\llbracket P[z/y] \rrbracket$. \square

Some more simplifications:

Remember that we have in Lemma 1 $\llbracket E\{E'/x\} \rrbracket^s = \llbracket E \rrbracket^{[s|x \rightarrow \llbracket E' \rrbracket^s]}$ if $\llbracket E' \rrbracket^s$ exists.

- $\boxed{(E_1 = E_2)[E'/x] \equiv (E_1\{E'/x\} = E_2\{E'/x\}) \wedge \text{is}(E')} :$

$$\begin{aligned} & \llbracket (E_1\{E'/x\} = E_2\{E'/x\}) \wedge \text{is}(E') \rrbracket \\ &= \{s, h \mid \llbracket E_1\{E'/x\} \rrbracket^s = \llbracket E_2\{E'/x\} \rrbracket^s \text{ and } \llbracket E' \rrbracket^s \text{ exists}\} \\ &= \{s, h \mid \llbracket E_1 \rrbracket^{[s|x \rightarrow \llbracket E' \rrbracket^s]} = \llbracket E_2 \rrbracket^{[s|x \rightarrow \llbracket E' \rrbracket^s]}\} \\ &= \{s, h \mid \llbracket E_1 \rrbracket^{[s|x \rightarrow \llbracket E' \rrbracket^s]} = \llbracket E_2 \rrbracket^{[s|x \rightarrow \llbracket E' \rrbracket^s]}\} \\ &= \{s, h \mid [s \mid x \rightarrow \llbracket E' \rrbracket^s], h \in \llbracket E_1 = E_2 \rrbracket\} \\ &= \llbracket (E_1 = E_2)[E'/x] \rrbracket \end{aligned}$$

- $\boxed{(E \mapsto E_1, E_2)[E'/x] \equiv (E\{E'/x\} \mapsto E_1\{E'/x\}, E_2\{E'/x\}) \wedge \text{is}(E')} :$

$$\begin{aligned} & \llbracket (E\{E'/x\} \mapsto E_1\{E'/x\}, E_2\{E'/x\}) \wedge \text{is}(E') \rrbracket \\ &= \{s, h \mid \text{dom}(h) = \{\llbracket E\{E'/x\} \rrbracket^s\} \text{ and } h(\llbracket E\{E'/x\} \rrbracket^s) = \langle \llbracket E_1\{E'/x\} \rrbracket^s, \llbracket E_2\{E'/x\} \rrbracket^s \rangle \\ & \quad \text{and } \llbracket E' \rrbracket^s \text{ exists}\} \\ &= \{s, h \mid \text{dom}(h) = \{\llbracket E \rrbracket^{[s|x \rightarrow \llbracket E' \rrbracket^s]}\} \\ & \quad \text{and } h(\llbracket E \rrbracket^{[s|x \rightarrow \llbracket E' \rrbracket^s]}) = \langle \llbracket E_1 \rrbracket^{[s|x \rightarrow \llbracket E' \rrbracket^s]}, \llbracket E_2 \rrbracket^{[s|x \rightarrow \llbracket E' \rrbracket^s]} \rangle\} \\ &= \{s, h \mid [s \mid x \mapsto \llbracket E' \rrbracket^s], h \in \llbracket E \mapsto E_1, E_2 \rrbracket\} \\ &= \llbracket (E \mapsto E_1, E_2)[E'/x] \rrbracket \end{aligned}$$

- $\boxed{\text{false}[E'/x] \equiv \text{false}} :$

$$\begin{aligned} & \llbracket \text{false}[E'/x] \rrbracket \\ &= \{s, h \mid [s \mid x \rightarrow \llbracket E' \rrbracket^s], h \in \llbracket \text{false} \rrbracket\} \\ &= \emptyset \\ &= \llbracket \text{false} \rrbracket \end{aligned}$$

- $\boxed{(P \Rightarrow Q)[E'/x] \equiv P[E'/x] \Rightarrow Q[E'/x]} :$

$$\begin{aligned} & \llbracket (P \Rightarrow Q)[E'/x] \rrbracket_\rho \\ &= \{s, h \mid [s \mid x \rightarrow \llbracket E' \rrbracket^s], h \in \llbracket P \Rightarrow Q \rrbracket_\rho\} \\ &= \{s, h \mid [s \mid x \rightarrow \llbracket E' \rrbracket^s], h \in ((\top \setminus \llbracket P \rrbracket_\rho) \cup \llbracket Q \rrbracket_\rho)\} \\ &= ((\top \setminus \{s, h \mid [s \mid x \rightarrow \llbracket E' \rrbracket^s], h \in \llbracket P \rrbracket_\rho\}) \cup \{s, h \mid [s \mid x \rightarrow \llbracket E' \rrbracket^s], h \in \llbracket Q \rrbracket_\rho\}) \\ &= ((\top \setminus \llbracket P[E'/x] \rrbracket_\rho) \cup \llbracket Q[E'/x] \rrbracket_\rho) \\ &= \llbracket P[E'/x] \Rightarrow Q[E'/x] \rrbracket_\rho \end{aligned}$$

- $\boxed{(\exists x.P)[E/x] \equiv (\exists x.P) \wedge is(E)} :$

$$\begin{aligned} & \llbracket (\exists x.P)[E/x] \rrbracket_\rho \\ &= \{s, h \mid [s \mid x \rightarrow \llbracket E \rrbracket^s], h \in \llbracket \exists x.P \rrbracket_\rho\} \\ &= \{s, h \mid \exists v.[s \mid x \rightarrow \llbracket E \rrbracket^s \mid x \rightarrow v], h \in \llbracket P \rrbracket_\rho\} \\ &= \{s, h \mid \exists v.[s \mid x \rightarrow v], h \in \llbracket P \rrbracket_\rho \text{ and } \llbracket E \rrbracket^s \text{ exists}\} \\ &= \llbracket (\exists x.P) \wedge is(E) \rrbracket_\rho \end{aligned}$$

- $\boxed{\begin{array}{l} \text{If } y \notin Var(E) \\ x \neq y \end{array} \text{ then } (\exists y.P)[E/x] \equiv \exists y.(P[E/x])} :$

First, notice that if $y \notin Var(E)$, then $\llbracket E \rrbracket^s = \llbracket E \rrbracket^{[s|y \rightarrow v]}$.

$$\begin{aligned} & \llbracket (\exists y.P)[E/x] \rrbracket_\rho \\ &= \{s, h \mid [s \mid x \rightarrow \llbracket E \rrbracket^s], h \in \llbracket \exists x.P \rrbracket_\rho\} \\ &= \{s, h \mid \exists v.[s \mid x \rightarrow \llbracket E \rrbracket^s \mid y \rightarrow v], h \in \llbracket P \rrbracket_\rho\} \\ &= \{s, h \mid \exists v.[s \mid y \rightarrow v \mid x \rightarrow \llbracket E \rrbracket^s], h \in \llbracket P \rrbracket_\rho\} \\ &= \{s, h \mid \exists v.[s \mid y \rightarrow v \mid x \rightarrow \llbracket E \rrbracket^{[s|y \rightarrow v]}], h \in \llbracket P \rrbracket_\rho\} \\ &= \{s, h \mid \exists v.[s \mid y \rightarrow v], h \in \llbracket P[E/x] \rrbracket_\rho\} \\ &= \llbracket \exists y.(P[E/x]) \rrbracket_\rho \end{aligned}$$

- $\boxed{\mathbf{emp}[E'/x] \equiv \mathbf{emp} \wedge is(E')} :$

$$\begin{aligned} &= \llbracket \mathbf{emp} \wedge is(E') \rrbracket \\ &= \{s, h \mid h = [] \text{ and } \llbracket E' \rrbracket^s \text{ exists}\} \\ &= \{s, h \mid [s \mid x \mapsto \llbracket E' \rrbracket^s], h \in \llbracket \mathbf{emp} \rrbracket\} \\ &= \llbracket \mathbf{emp}[E'/x] \rrbracket \end{aligned}$$

- $\boxed{(P * Q)[E'/x] \equiv P[E'/x] * Q[E'/x]} :$

$$\begin{aligned} &= \llbracket P[E'/x] * Q[E'/x] \rrbracket_\rho \\ &= \{s, h \mid \exists h_0, h_1. h_0 \# h_1, h_0.h_1 = h, s, h_0 \in \llbracket P[E'/x] \rrbracket_\rho \text{ and } s, h_1 \in \llbracket Q[E'/x] \rrbracket_\rho\} \\ &= \{s, h \mid \exists h_0, h_1. h_0 \# h_1, h_0.h_1 = h, [s \mid x \mapsto \llbracket E' \rrbracket^s], h_0 \in \llbracket P \rrbracket_\rho \text{ and } [s \mid x \mapsto \llbracket E' \rrbracket^s], h_1 \in \llbracket Q \rrbracket_\rho\} \\ &= \{s, h \mid [s \mid x \mapsto \llbracket E' \rrbracket^s], h \in \llbracket P * Q \rrbracket_\rho\} \\ &= \llbracket (P * Q)[E'/x] \rrbracket_\rho \end{aligned}$$

- $\boxed{(P \multimap Q)[E'/x] \equiv P[E'/x] \multimap Q[E'/x]} :$

$$\begin{aligned} &= \llbracket P[E'/x] \multimap Q[E'/x] \rrbracket_\rho \\ &= \{s, h \mid \forall h'. \text{ if } h' \# h \text{ and } s, h' \in \llbracket P[E'/x] \rrbracket_\rho \text{ then } s, h.h' \in \llbracket Q[E'/x] \rrbracket_\rho\} \\ &= \{s, h \mid \forall h'. \text{ if } h' \# h \text{ and } [s \mid x \mapsto \llbracket E' \rrbracket^s], h' \in \llbracket P \rrbracket_\rho \text{ then } [s \mid x \mapsto \llbracket E' \rrbracket^s], h.h' \in \llbracket Q \rrbracket_\rho\} \\ &= \{s, h \mid [s \mid x \mapsto \llbracket E' \rrbracket^s], h \in \llbracket P \multimap Q \rrbracket_\rho\} \\ &= \llbracket (P \multimap Q)[E'/x] \rrbracket_\rho \end{aligned}$$