

Isogeny++: From ideals to modules

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Damien Robert

Équipe Canari, Inria Bordeaux Sud-Ouest



université
de BORDEAUX

Inria

From ideals to modules

- Lattices: RingLWE \rightarrow ModuleLWE
- Codes: Hamming metric \rightarrow Sum Rank metric
- Isogenies: Ideals \rightarrow Modules?
Dimension 1 \rightarrow Dimension g ?

- ☹ Increasing dimension in isogeny based cryptography is costly...
- 😊 Dimension 2 already provides a lot of flexibility (Kani...)

- **Open question:** is it worth it? (Beyond Kani!)

Ideals and isogenies: the oriented case

- $E/k, k = \mathbb{F}_q$, elliptic curve with a primitive orientation by a quadratic imaginary order
 $R = \mathbb{Z}[\sqrt{-\Delta}] \hookrightarrow \text{End}_k(E)$
- **Oriented isogeny:** $\phi : E_1 \rightarrow E_2$ that commutes with the orientations
- **Oriented kernel:** K stable by R
Unique R -orientation compatible on E/K with the quotient isogeny $E \rightarrow E/K$, and the isogeny is horizontal or ascending

Example: Frobenius orientation

- E/k with non trivial π_k -action: ordinary curves, supersingular curves over \mathbb{F}_p
- π_k -oriented isogenies = rational isogenies.
- **Invertible ideals I of $R \Leftrightarrow$** oriented horizontal isogenies $\phi_I : E \rightarrow E_I$
[Colò-Kohel 2020, Onuki 2020]
- $K = \text{Ker } \phi_I = E[I], \quad I = \{\alpha \in R \mid \alpha(K) = 0\}, \quad \deg \phi_I = N(I)$
- $\widetilde{\phi}_I = \phi_{\bar{I}} : E_I \rightarrow E$
- $I \simeq J \Leftrightarrow E_I \simeq E_J$
- Special case: p inert in R (can only happen for an orientation on a supersingular curve E/\mathbb{F}_{p^2})
- $\pi_p : E \rightarrow E^\sigma$ is not represented by an ideal
- $\rho_R(E)$ representation of R on the k -vector space $T_0(E)$
- an oriented isogeny $\phi : E \rightarrow E'$ comes from an ideal iff the representations $\rho_R(E)$ and $\rho_R(E')$ are equivalent.

Ideal and isogenies: the supersingular case

- Deuring correspondance
- maximal orders in $B_{p,\infty} =$ supersingular curves E/\mathbb{F}_{p^2} (up to quadratic twists and Galois conjugates)
- ideals = isogenies; $\deg \phi_I = N(I) := \text{nrd}(I)$

Ideal to isogeny: $I \Leftrightarrow E[I]$

- Easy if $\text{End}(E)$ known, $N(I)$ smooth and $N(I)$ -torsion accessible
- Many algorithms to handle the general case: KLPT, Eichler orders, refreshing the torsion, endomorphisms, Clapotis...
- Lots of research effort
- 😊 SQISign and variants

A general equivalence of category

- E_0/k primitively oriented by R quadratic imaginary ($Z(R) = R$)
- $E_0/k = \mathbb{F}_{p^2}$ with $R = \text{End}(E_0)$ maximal quaternionic order ($Z(R) = \mathbb{Z}$)

Theorem

There is an *antiequivalence of category* between the category of $Z(R)$ -oriented abelian varieties A k -isogenous to E_0^g (with the technical condition $\rho_{Z(R)}(A) \simeq \bigoplus_{i=1}^g \rho_{Z(R)}(E_0)$) and $Z(R)$ -oriented k -morphisms; and the category of finitely presented torsion free (left) R -modules M of rank g and R -module morphisms

[Waterhouse 1969], [Kani 2011], [Jordan, Keeton, Poonen, Rains, Shepherd-Barron, Tate 2018], [Page-R. 2023]

Alternative approaches to equivalences of category of abelian varieties via lifting to characteristic zero: Deligne, Howe, Marseglia...

Corollary

- principal polarisation $\lambda_A : A \rightarrow \widehat{A} = \mathfrak{a}$ unimodular Hermitian R -form H_A on M_A
- N -isogeny $\phi : (A, \lambda_A) \rightarrow (B, \lambda_B) = N$ -similitude $\Phi : (M_B, H_B) \rightarrow (M_A, H_A)$:

$$\Phi^* H_A = N H_B$$

[Kirschmer, Narbonne, Ritzenthaler, R. 2021] (project started in 2011 with Christophe!)

The equivalence

Serre's generalised Ext and Tor functors: $\mathcal{F}(M) := \text{Ext}_R^1(M, E_0)$ E_0 "=" compact projective generator

Definition

If $R^m \rightarrow R^n \rightarrow M \rightarrow 0$ is a presentation of a R -module M , with corresponding matrix Φ , $\mathcal{F}(M) := \text{Ext}_R^1(M, E_0)$ is the kernel of the morphism $E_0^n \rightarrow E_0^m$ given by Φ^T and the R -orientation:

$$0 \rightarrow \mathcal{F}(M) \rightarrow E_0^n \rightarrow E_0^m$$

\mathcal{F} is a **contravariant exact functor** from f.p. R -modules to proper group schemes over k

- Ideals: $\mathcal{F}(R/I) \simeq E_0[I]$, $\mathcal{F}(I) \simeq E_0/E_0[I]$
- Abelian varieties: If M is torsion free of rank g , $A = \mathcal{F}(M)$ is an abelian variety of rank g
- Duality: $A^\vee \simeq \mathcal{F}(M^\vee)$
- Torsion: $A[n] \simeq \text{Ext}_R^1(M, E_0[n])$
- Rational points: $A(k') \simeq \text{Ext}_R^1(M, E_0(k'))$

Inverse map: $A \mapsto \text{Hom}_{Z(R)}(A, E_0)$: module of (oriented) morphisms from A to E_0

Warmup: ideals

The oriented case:

- $\mathcal{F}(R) = E_0$, so $\phi_I : E_0 \rightarrow E_I$ corresponds to $I \rightarrow R$
- Canonical unimodular Hermitian form on I :

$$H_I(x, y) = \frac{x\bar{y}}{N(I)}$$

- The inclusion $(I, H_I) \subset (R, H_R)$ is a $N(I)$ -similitude
- Handles ascending isogenies: I not invertible (the R -orientation needs not be primitive on E_I)

The supersingular case ($R = O_0$):

- Maximal orders \Leftrightarrow left O_0 -ideals
- To an order O we associated a connecting (O_0, O) -ideal
- To a left O_0 -ideal I we associate the right order $O_R(I)$

N.B.: could use duality to get an equivalence of categories, but contravariance is more practical

Modules to abelian varieties

- $R^m \rightarrow R^n \rightarrow M \rightarrow 0$ presentation of M
- $0 \rightarrow A \hookrightarrow E_0^n \rightarrow E_0^m$ co-presentation of $A = \mathcal{F}(M)$

Example: $I = (\alpha, \beta)$, with syzygys of rank 1: $u\alpha + v\beta = 0$

$$R \xrightarrow{(u,v)^T} R^2 \xrightarrow{(\alpha,\beta)} I \subset R \quad \Leftrightarrow \quad E_0 \twoheadrightarrow E_I \hookrightarrow E_0^2 \rightarrow E_0$$

- $E_0 \rightarrow E_0^2, P \mapsto (\alpha P, \beta P)$ has kernel $E_0[I]$, so the image is isomorphic to E_I
- $E_I \hookrightarrow E_0^2$ is also given by the kernel of $E_0^2 \rightarrow E_0, (P, Q) \mapsto uP + vQ$

Module to explicit abelian variety:

- Find a nice N -similitude $(M, H_M) \hookrightarrow (R^g, \bigoplus_{i=1}^g H_R)$
- Convert to $E_0^g \twoheadrightarrow A_M$

☞ There are unimodular Hermitian R -modules such that no such N -similitude exist for any N , c.f. the arithmetic obstructions in [Kirschmer, Narbonne, Ritzenthaler, R. 2021]

Abelian variety to module:

- Find n morphisms $\phi_i : A \rightarrow E_0$ whose kernels intersect trivially

Example: a double path $E_I \rightarrow E_0!$

- Find the R -lattice of relations on the ϕ_i

Find relations by testing on points of smooth order. Each relation reduces the tentative module M_A . Use the principal polarisation on A as a stop criterion (pairings). N.B.: Explicit endomorphisms on $E_0 \Leftrightarrow$ abstract endomorphisms.

- $A \hookrightarrow E_0^n \rightarrow E_0^m$ gives M_A

Similitudes to isogenies

Module morphism to morphism of abelian varieties:

$$\begin{array}{ccccccc} R^{m_1} & \longrightarrow & R^{n_1} & \twoheadrightarrow & M_1 & \longrightarrow & 0 \\ \uparrow & & \uparrow & & \uparrow & & \\ \vdots & & \vdots & & & & \\ R^{m_2} & \longrightarrow & R^{n_2} & \twoheadrightarrow & M_2 & \longrightarrow & 0 \end{array}$$

$$\begin{array}{ccccccc} 0 & \longrightarrow & A_1 & \hookrightarrow & E_0^{n_1} & \longrightarrow & E_0^{m_1} \\ & & \downarrow & & \vdots & & \vdots \\ 0 & \longrightarrow & A_2 & \hookrightarrow & E_0^{n_2} & \longrightarrow & E_0^{m_2} \end{array}$$

R^n is a projective module, so we can lift module maps. The commutative diagram allows to find the kernel of $A_1 \rightarrow A_2$.

- N -similitudes $\Leftrightarrow N$ -isogenies
- $\phi : A_1 \rightarrow A_2 \Leftrightarrow (M_2, H/N) \subset (M_1, H)$
An isogeny is an epimorphism (with finite kernel) so corresponds to a monomorphism (=inclusion) of modules (with finite cokernel)
- $\text{Ker } \phi = A_1[M_2]$
- Equivalence **practical** if N smooth and the N -torsion on E_0 is accessible
- **Open question for the general case:** ModuleKLPT?

Cryptographic applications?

[Page-R. 2023]

- **Clapotis**: Class group Action in Polynomial Time via Sesquilinear forms
- Original motivation for this work: “new” ModuleKLPT algorithm for $M = I \oplus \bar{I} \subset R \oplus R$

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[Page-R. 2023]

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- **Clapoti**: bypass the equivalence of category by just using Kani... again...



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Help needed! Any interesting cryptographic application of modules?

Strong assumption: we can extend in dimension g all our algorithmic tools and security assumptions from dimension 1 to dimension g

- **ModuleSQISign**: Short signatures for oriented isogenies?
- **ModuleSIDH**: combining torsion noise and oriented commutative group action for **key exchange**?

The isogeny graph of oriented isogenies in higher dimension

- M torsion free of rank g : $M \simeq R^{g-1} \oplus I$ Assume R maximal for simplicity
- $A \simeq E_0^{g-1} \times E_I$
- # $\text{Cl}(R)$ isomorphism classes of **non-polarised** R -oriented abelian varieties R -isogenous to E_0^g
- **Polarisations** add supersingular like graph complexity if $g > 1$ ($\text{End}_R(E_0^g) = M_g(R)$)
- **Universal group action**: $I \cdot (M, H_M) = (IM, H_M/N(I)) \subset (M, H_M)$
- $I \cdot A = A_I := A/A[I]$
- Intuition: multiplication by $[n] \Rightarrow$ multiplication by $[I]$
- Multiple orbits; linked together by oriented isogenies (which are not multiplication by $[I]$)

Example:

- E_0/\mathbb{F}_p supersingular and $g = 2$: graph of supersingular abelian surfaces isogeneous to E_0^2 over \mathbb{F}_p and \mathbb{F}_p -rational isogenies
The graph contains the Weil restriction $W_{\mathbb{F}_{p^2}/\mathbb{F}_p} E$ of supersingular elliptic curves over \mathbb{F}_{p^2} (these are neither Jacobians nor product of curves over \mathbb{F}_p).
- Conjecture: $\approx p^{3/2}$ nodes
- Universal group action from $\text{Cl}(\mathbb{Z}[\sqrt{-p}])$
- If $\ell = \ell$ splits in R , $A[\ell] = A[I] \oplus A[\bar{I}] \Rightarrow$ action by I and \bar{I}
(+ $\ell + 1$ other R -oriented ℓ -isogenies?)

ModuleSQISign: Short signatures for oriented isogenies?

- $\phi_I : E_0 \rightarrow E_I$
- Recovering I from $(E_0, E_I) \Leftrightarrow$ recovering the module associated to $E_0 \times E_I$
 $E_0 \times E_I$ is doubly oriented!

\Rightarrow SQISign like protocol in dimension 2 (👁️ not SQISign2d!)

$$\begin{array}{ccc} E_0 \times E_0 & \xrightarrow{\phi_{sec}} & E_0 \times E_I \\ \downarrow \phi_{com} & & \downarrow \phi_{chl} \\ A & \xleftarrow{\phi_{resp}} & B \end{array}$$

- **Soundness:** check that the response is not R -backtracking through the challenge
We want an R -endomorphism on $E_0 \times E_I$ which does not come from R !
- **ZK:** the commitment should probably not be R -backtracking either
- Needs a generalised ModuleTolsogeny for the response

ModuleSIDH: Noisy group action key exchange?

- Commutative group action on a supersingular like graph
- ⇒ Mask the torsion in a SIDH-like key exchange by using this commutative group action (like M-SIDH but using $[I]$ rather than $[n]$)
- ⇒ Hide the commutative group action in a CSIDH-like key exchange by adding a SIDH-like torsion exchange

$$\begin{array}{ccccc}
 A_0 & \xrightarrow{\phi_a} & A_{a_1} & \xrightarrow{[a]} & (A_{a_2}, [a] \circ \phi_a(A_0[N_B])) \\
 \downarrow \phi_b & & \downarrow \phi'_b & & \downarrow \phi''_b \\
 A_{b_1} & \xrightarrow{\phi'_a} & A_{a_1, b_1} & \xrightarrow{[a]} & A_{a_2, b_1} \\
 \downarrow [b] & & \downarrow [b] & & \downarrow [b] \\
 (A_{b_2}, [b] \circ \phi_b(A_0[N_A])) & \xrightarrow{\phi''_a} & A_{a_1, b_2} & \xrightarrow{[a]} & A_{a_2, b_2}
 \end{array}$$

- ϕ_a : oriented N_A -isogeny; ϕ_b : oriented N_B -isogeny
- Speed up trick: do a standard SIDH key exchange over \mathbb{F}_{p^2} , take Weil restriction to \mathbb{F}_p , apply group action of $\text{Cl}(\mathbb{Z}[\sqrt{-p}])$ in dimension 2
- Size: $p = 4\lambda$ (or 6λ ?); $J(A_{a_2})$: $3 \log_2(p)$; torsion on deterministic R -basis: $4 \log_2(p)$ (or $3 \log_2 p$ using pairings?)
Total: $6 \log_2 p = 24\lambda$ (vs $3.5 \log_2 p$ for SIDH)

Conclusion: the module equivalence of category

- The module equivalence of category is often more natural than the ideal one: clear distinction of objects and morphisms
 - Many algorithmic operations already done in dimension 1 (e.g., double path to E_0) come from the module interpretation
 - Unified framework to handle the oriented and supersingular case (still modules, but different rings)
- ⇒ Can keep track of forgetting the orientation or Weil restrictions purely at the module level
- Generalizes to higher dimension
 - Also able to keep track of level structure
 - The current methods implicitly use the conductor square and excision to embed level structure information via suborders of conductor divisible by the level, but that's arguably less natural
 - New cryptographic protocols?

Torsion free f.p. R -modules

- In both cases: rank 1 torsion free module = ideal

Oriented case (R is a Bass ring)

- $M \simeq I_1 \oplus I_2 \oplus \cdots \oplus I_g$
- $R \subset O(I_1) \subset O(I_2) \subset \cdots \subset O(I_g)$
- $\det M = I_1 \cdot I_2 \cdots \cdot I_g$ invertible R_g -ideal
- The isomorphism class of M only depend on (R_1, \dots, R_g) and $\det M$
- Example: if all I_j are invertible in R ($\Leftrightarrow O(I_j) = R$),

$$M \simeq R^{g-1} \oplus I_1 \cdot I_2 \cdots \cdot I_g$$

Supersingular case

- $M \simeq R^g$ if $g > 1$

Non principal polarisations

- M torsion free, $V = M \otimes_{\mathbb{Z}} \mathbb{Q}, K = R \otimes_{\mathbb{Z}} \mathbb{Q}$
- HK -hermitian form on V
- Orthogonal: $M^{\#} = \{v \in V, H(\cdot, v) \subset R\}$

- H induces an isomorphism $M^{\#} \simeq M^{\vee}, m^{\#} \mapsto H(\cdot, m^{\#})$
- H is integral on $M^{\#} \Leftrightarrow M^{\#} \subset M$
- We then obtain a polarisation on $M^{\vee}: M^{\vee} \simeq M^{\#} \subset M$
- This gives a polarisation $\lambda : A \rightarrow A^{\vee}$ with kernel $\mathbb{F}(M/M^{\#})$
- The polarisation $n\lambda$ corresponds to H/n

- Principal polarisation: $M = M^{\#}$

Non R -backtracking isogenies

Non (partially) backtracking isogeny:

- $\phi : A \rightarrow B$ N -isogeny is non partially backtracking (nbt) $\Leftrightarrow \text{Ker } \phi$ of rank g
- $\phi_1 : A_1 \rightarrow A_2, \phi_2 : A_2 \rightarrow A_3$ nbt, then $\phi_2 \circ \phi_1$ nbt iff $\text{Ker } \phi_2 \cap \text{Ker } \widetilde{\phi_1} = 0$
- If $\phi_2 \circ \phi_1$ is nbt, ϕ_1, ϕ_2 is nbt
- If $\phi : A \rightarrow B$ nbt N -isogeny, and $N = \prod \ell_i$, ϕ uniquely decomposes as $\phi = \prod \phi_i$, with ϕ_i a ℓ_i -isogeny

Non R -backtracking isogeny: Assume all degrees prime to the conductor of R

- $\phi : A \rightarrow B$ is non R -backtracking iff it is nbt and does not come from the action of an ideal I
- If ϕ is nbt but comes from I , $\phi = \phi_2 \circ \phi_1$, then ϕ_i comes from I_i
- If ϕ nbt, it suffices to check that some subgroup $\text{Ker } \phi[\ell^e]$ is not induced by an ideal to know that ϕ is not R -backtracking

Combined with the following lemma, this gives a way to check that the response is not R -backtracking through the challenge for ModuleSQISign:

Lemma

$\phi_1 : A_1 \rightarrow A_2, \phi_2 : A_2 \rightarrow A_3, \phi_3 : A_3 \rightarrow A_4, \phi_4 : A_4 \rightarrow A_5$ such that $\phi_2 \circ \phi_1, \phi_3 \circ \phi_2$ and $\phi_4 \circ \phi_3$ are nbt. Then $\phi_4 \circ \phi_3 \circ \phi_2 \circ \phi_1$ is ℓ -nbt for each $\ell \mid \# \text{Ker } \phi_2 \wedge \# \text{Ker } \phi_3$, i.e. the ℓ -Sylow of its kernel is of rank g