

INTRODUCTION TO ALGEBRAIC OPERADS

EXERCISE SHEET 1: Homological algebra

1. CATEGORIES

EXERCISE 1 (Yoneda lemma). Given a small category \mathcal{C} , a functor $\mathcal{F} : \mathcal{C} \rightarrow \text{Set}$ and an object C of \mathcal{C} , prove that the map

$$\begin{aligned} \text{Nat}(\mathcal{C}(C, \cdot), \mathcal{F}) &\longrightarrow \mathcal{F}(C) \\ \tau &\longmapsto \tau_C(\text{id}_C) \end{aligned}$$

is a bijection. This result is known as the *Yoneda lemma*.

EXERCISE 2 (The 2-category Cat). We define a 2-category \mathcal{C} to be the following data:

- (i) A class of objects $\text{Ob}(\mathcal{C})$.
- (ii) For every object $X, Y \in \mathcal{C}$ a category $\mathcal{C}(X, Y)$ and for every $X \in \mathcal{C}$ an identity object id_X in $\mathcal{C}(X, X)$.
- (iii) For every objects $X, Y, Z \in \mathcal{C}$ a composition bifunctor $\mathcal{C}(Y, Z) \times \mathcal{C}(X, Y) \rightarrow \mathcal{C}(X, Z)$.

These data have to satisfy the same associativity and identity axioms as a standard category. The elements $f \in \mathcal{C}(X, Y)$ are called 1-morphisms and denoted $f : X \rightarrow Y$ and the elements $\tau \in \mathcal{C}(X, Y)(f, g)$ are called 2-morphisms and denoted $\tau : f \Rightarrow g$.

Prove that Cat is a 2-category whose 1-morphisms are functors and 2-morphisms are natural transformations.

EXERCISE 3 (Adjunction).

1. Given two functors $\mathcal{F}, \mathcal{G} : \mathcal{C} \rightarrow \mathcal{D}$, prove that a collection of bijections

$$\phi_{X,Y} : \mathcal{C}(\mathcal{F}(X), Y) \xrightarrow{\sim} \mathcal{D}(X, \mathcal{G}(Y))$$

defines a natural equivalence if and only if for all objects $X, X' \in \mathcal{C}$ and $Y, Y' \in \mathcal{D}$, and all morphisms $f : X \rightarrow X'$, $h : \mathcal{F}(X') \rightarrow Y$ and $g : Y \rightarrow Y'$ we have that

$$\phi_{X,Y'}(g \circ h \circ \mathcal{F}(f)) = \mathcal{G}(g) \circ \phi_{X',Y}(h) \circ f .$$

2. Given two adjoint functors $\mathcal{F} \dashv \mathcal{G}$ as in the previous question, prove that the diagram

$$\begin{array}{ccc} \mathcal{F}(X) & \xrightarrow{h} & Y \\ \downarrow \mathcal{F}(f) & & \downarrow g \\ \mathcal{F}(X') & \xrightarrow{h'} & Y' \end{array}$$

is commutative if and only if the diagram

$$\begin{array}{ccc} X & \xrightarrow{\phi_{X,Y}(h)} & \mathcal{G}(Y) \\ \downarrow f & & \downarrow \mathcal{G}(g) \\ X' & \xrightarrow{\phi_{X',Y'}(h')} & \mathcal{G}(Y') \end{array}$$

is commutative.

3. Prove that two functors $\mathcal{F} : \mathcal{C} \rightleftarrows \mathcal{D} : \mathcal{G}$ form an adjunction $\mathcal{F} \dashv \mathcal{G}$ if and only if there exists two natural transformations $\eta : \text{id}_{\mathcal{C}} \Rightarrow \mathcal{G} \circ \mathcal{F}$ and $\varepsilon : \mathcal{F} \circ \mathcal{G} \Rightarrow \text{id}_{\mathcal{D}}$ such that $\varepsilon \mathcal{F} \circ \mathcal{F} \eta = \text{id}_{\mathcal{F}}$ and $\mathcal{G} \varepsilon \circ \eta \mathcal{G} = \text{id}_{\mathcal{G}}$. The natural transformations η and ε are then respectively called the *unit* and the *counit* of the adjunction.

EXERCISE 4 (Monad). We define an endofunctor of a category \mathcal{C} to be a functor $\mathcal{C} \rightarrow \mathcal{C}$ and denote $\text{EndoFun}_{\mathcal{C}}$ the category of endofunctors of \mathcal{C} (see Exercise 2).

1. Prove that the composition \circ and the identity functor $\text{id}_{\mathcal{C}}$ endow the category $\text{EndoFun}_{\mathcal{C}}$ with a structure of strict monoidal category.

We define a monad to be a monoid in the monoidal category $\text{EndoFun}_{\mathcal{C}}$.

2. Prove that an adjunction $\mathcal{F} : \mathcal{C} \rightleftarrows \mathcal{D} : \mathcal{G}$ with unit η and counit ε gives rise to a monad structure on the endofunctor $\mathcal{G}\mathcal{F} : \mathcal{C} \rightarrow \mathcal{C}$ with multiplication $\mathcal{G}\varepsilon\mathcal{F}$ and unit η .

An algebra over a monad (\mathcal{M}, μ, η) is defined to be an object C of \mathcal{C} together with a morphism $\gamma_A : \mathcal{M}(A) \rightarrow A$ making the following diagrams commute

$$\begin{array}{ccc} \mathcal{M}(\mathcal{M}(A)) & \xrightarrow{\mathcal{M}(\gamma_A)} & \mathcal{M}(A) \\ \downarrow \mu_A & & \downarrow \gamma_A \\ \mathcal{M}(A) & \xrightarrow{\gamma_A} & A \end{array} \quad \begin{array}{ccc} \text{id}_{\mathcal{C}}(A) & \xrightarrow{\eta_A} & \mathcal{M}(A) \\ & \searrow & \downarrow \gamma_A \\ & & A \end{array} .$$

3. Prove that the forgetful functor $\mathcal{U} : \text{Vect} \rightarrow \text{Set}$ has a left-adjoint \mathcal{F} and that an algebra structure on a set V over the monad $\mathcal{U}\mathcal{F}$ is then exactly a vector space structure on V .

2. HOMOLOGICAL ALGEBRA

EXERCISE 5 (Euler characteristic). A chain complex is said to be bounded if up to reindexing it is of the form

$$0 \longrightarrow C_n \longrightarrow \cdots \longrightarrow C_0 \longrightarrow 0 .$$

1. Prove that for a bounded chain complex C_* ,

$$\sum_{i=0}^n (-1)^i \dim(C_i) = \sum_{i=0}^n (-1)^i \dim(H_i(C)) .$$

This number is then called its *Euler characteristic* and denoted $\chi(C)$.

2. Prove that every short exact sequence of bounded chain complexes

$$0 \longrightarrow A_* \longrightarrow B_* \longrightarrow C_* \longrightarrow 0,$$

satisfies the relation $\chi(B) = \chi(A) + \chi(C)$.

EXERCISE 6 (Diagram chasing lemmas).

1. Prove that in a short exact sequence $0 \rightarrow A_* \rightarrow B_* \rightarrow C_* \rightarrow 0$ of chain complexes, if two of the three chain complexes are exact then so is the third.

2. Consider a commutative diagram of linear maps

$$\begin{array}{ccccccccc} A_1 & \longrightarrow & B_1 & \longrightarrow & C_1 & \longrightarrow & D_1 & \longrightarrow & E_1 \\ \downarrow a & & \downarrow b & & \downarrow c & & \downarrow d & & \downarrow e \\ A_2 & \longrightarrow & B_2 & \longrightarrow & C_2 & \longrightarrow & D_2 & \longrightarrow & E_2 \end{array}$$

whose rows are exact, and such that b and d are isomorphisms, a is surjective and e is injective. Prove that the linear map c is an isomorphism. This result is known as the *five lemma*.

3. Prove that in a commutative diagram of chain maps

$$\begin{array}{ccccccccc} 0 & \longrightarrow & A_* & \longrightarrow & B_* & \longrightarrow & C_* & \longrightarrow & 0 \\ & & \downarrow & & \downarrow & & \downarrow & & \\ 0 & \longrightarrow & D_* & \longrightarrow & E_* & \longrightarrow & F_* & \longrightarrow & 0 \end{array}$$

whose rows are exact, if two of the vertical morphisms are quasi-isomorphisms then so is the third.

EXERCISE 7 (Cone). Let C and D be two chain complexes and $f : C \rightarrow D$ be a chain map. We define the cone of f to be the chain complex $\text{Cone}(f)$ whose degree n part is $C_{n-1} \oplus D_n$ and whose differential is given by

$$\partial(c, d) := (-\partial_C c, \partial_D d - f(c)).$$

This is often written as

$$\partial_{\text{Cone}(f)} := \begin{pmatrix} -\partial_C & 0 \\ -f & \partial_D \end{pmatrix}.$$

1. Check that the above formula indeed defines a differential on $\text{Cone}(f)$.

2. Prove that the sequence

$$0 \longrightarrow D \longrightarrow \text{Cone}(f) \longrightarrow sC \longrightarrow 0$$

is exact and make the connecting homomorphism explicit.

3. Prove that f is a quasi-isomorphism if and only if $\text{Cone}(f)$ is exact.

4. Suppose that f and f' are two homotopic chain maps $C \rightarrow D$. Show that there exists a quasi-isomorphism $\text{Cone}(f) \rightarrow \text{Cone}(f')$.

We say that a chain complex C_* is contractible if the chain map id_C is null-homotopic, i.e. homotopic to the null chain map.

5. Prove that if f is a homotopy equivalence, then $\text{Cone}(f)$ is contractible.

A chain complex (C_*, ∂) is said to be split if there exists a collection of maps $s_n : C_n \rightarrow C_{n+1}$ such that $\partial_{n+1} s_n = \text{id}_{C_n}$.

6. Prove that a split chain complex is exact if and only if it is contractible.

7. Prove that $\text{Cone}(C) := \text{Cone}(\text{id}_C)$ is split exact.