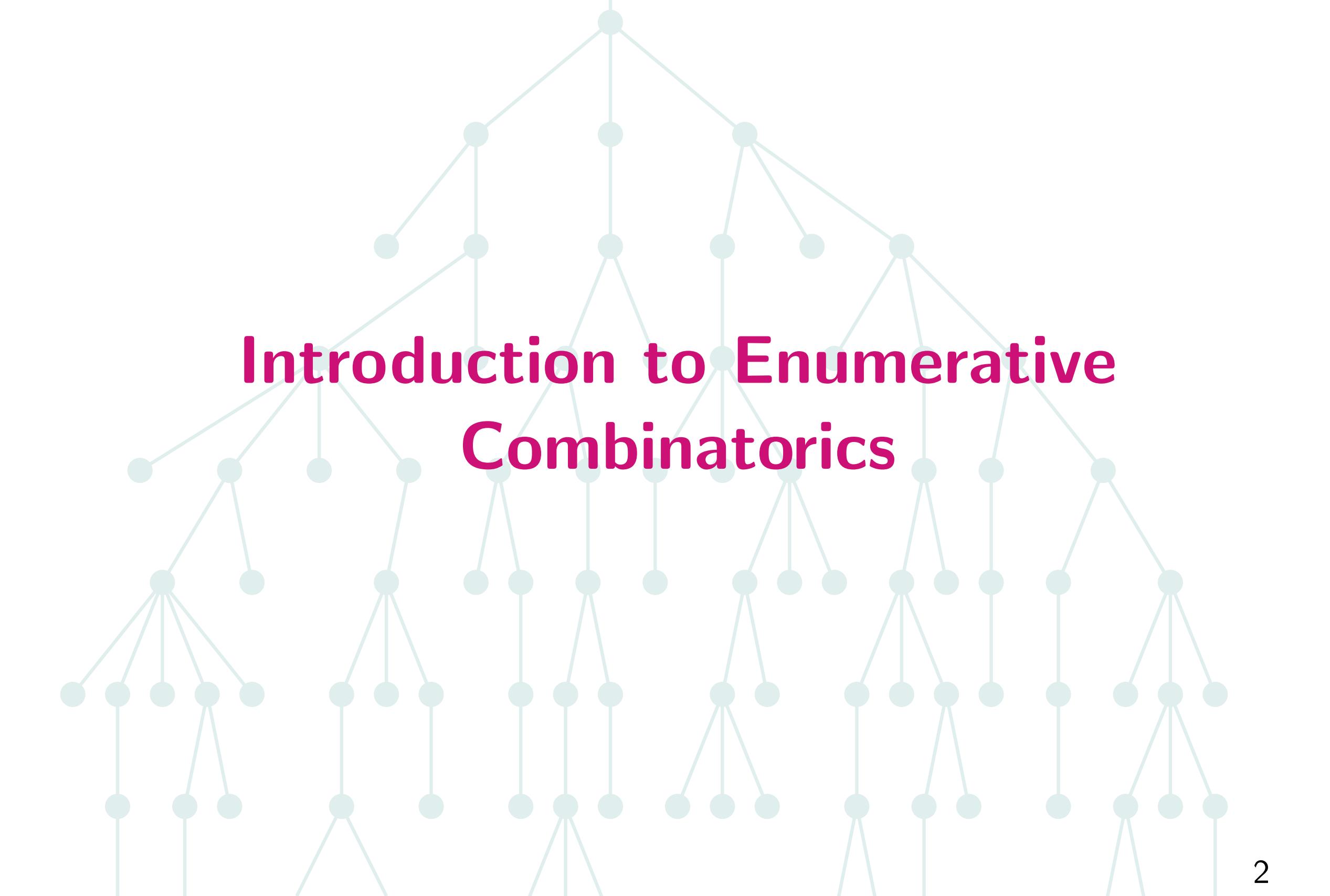


Gasket decomposition of 3-colored planar maps

Juliette Schabanel

LaBRI, Université de Bordeaux

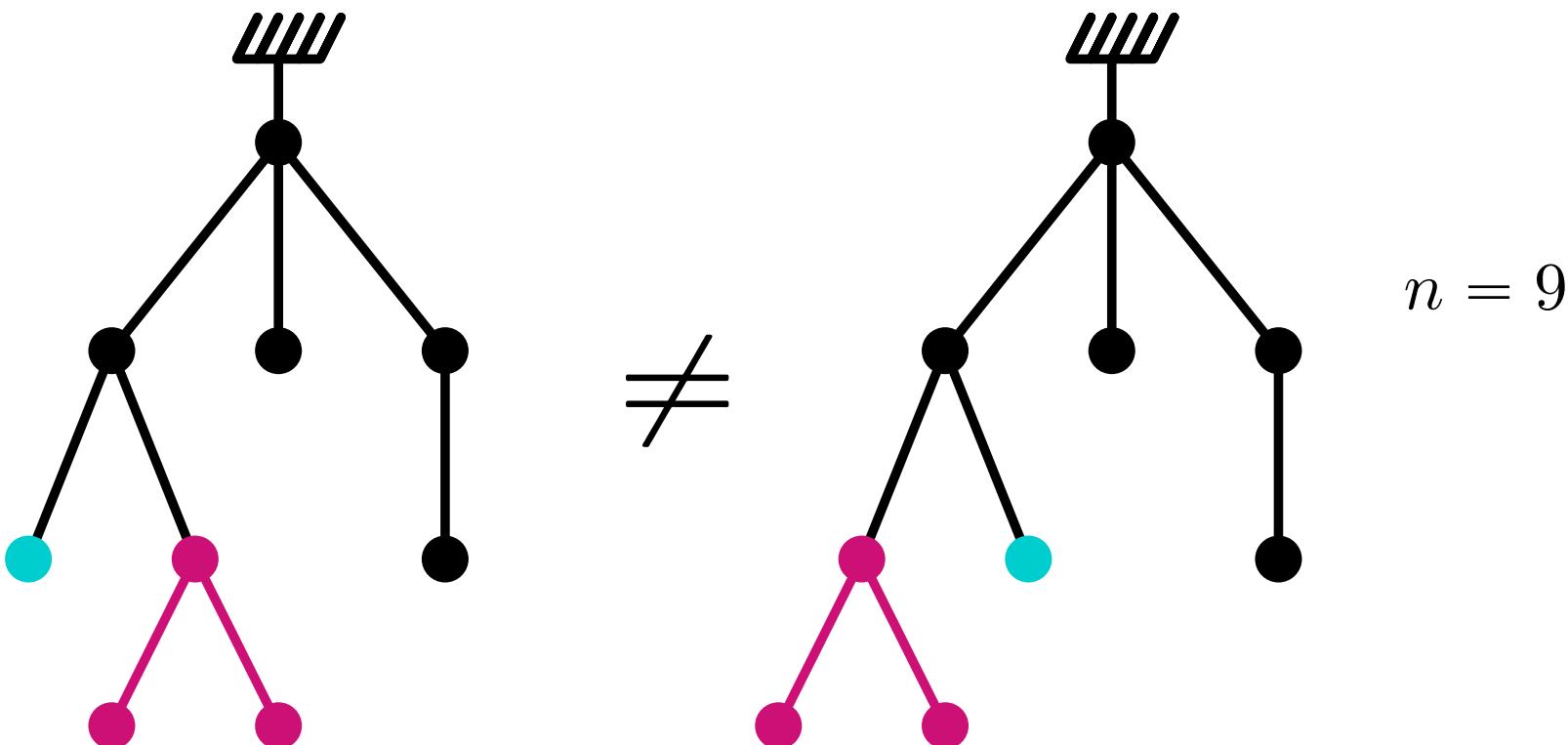


Introduction to Enumerative Combinatorics

First example: Trees

Plane tree = rooted tree with ordered children.

Size n = number of vertices.



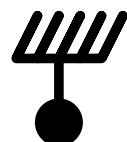
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size 1 :



$$t_1 = 1$$

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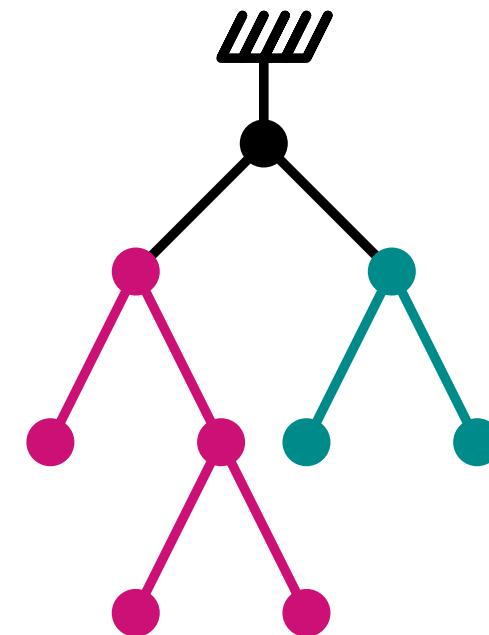
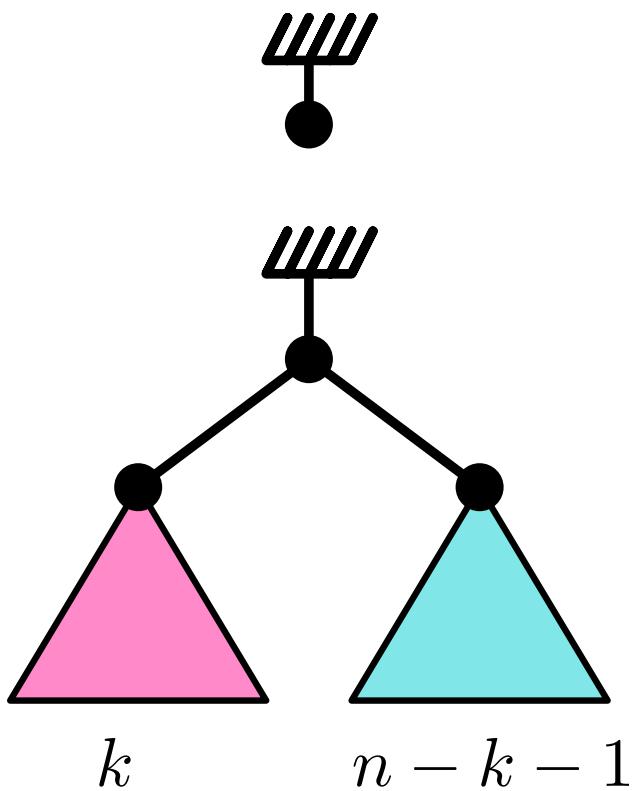
Binary trees: all vertices have 0 or 2 children.

size 1 :



$t_1 = 1$

size $n > 1$:



$n = 9$

$k = 5$

$n - k - 1 = 3$

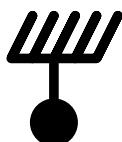
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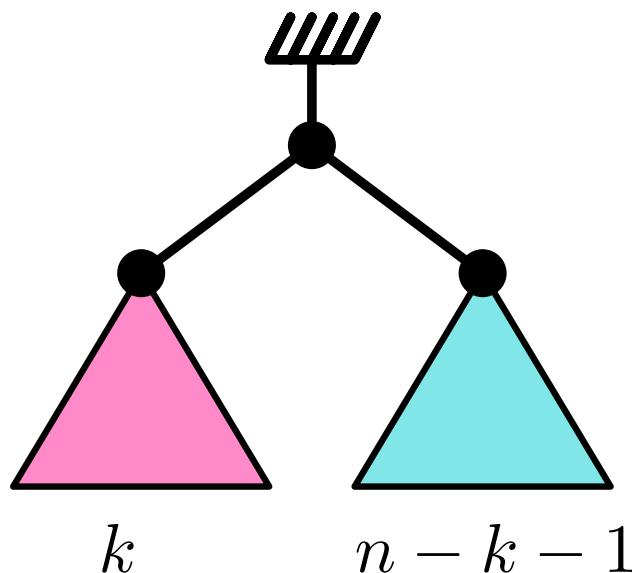
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$$t_n = \sum_{k=1}^{n-1} t_k t_{n-k-1}$$

Solution : $t_n = \frac{1}{n+1} \binom{2n}{n}$

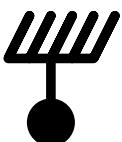
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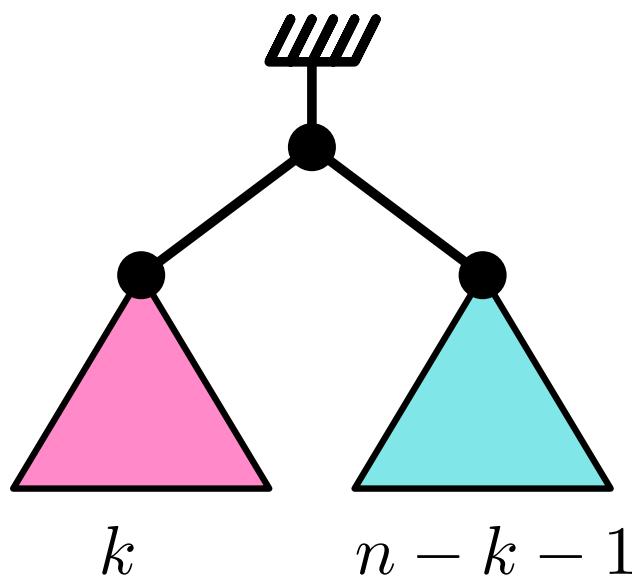
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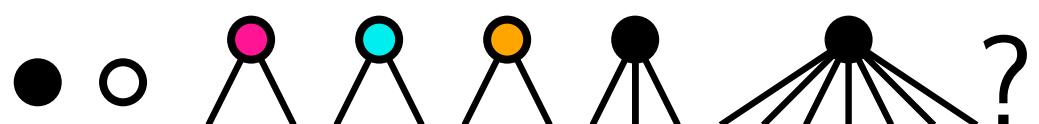
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Trees with nodes



Generating functions

The **generating function** of a combinatorial class \mathcal{C} is the series

$$C(t) = \sum_{\gamma \in \mathcal{C}} t^{|\gamma|} = \sum_{n \geq 0} c_n t^n \in \mathbb{Q}[[t]].$$

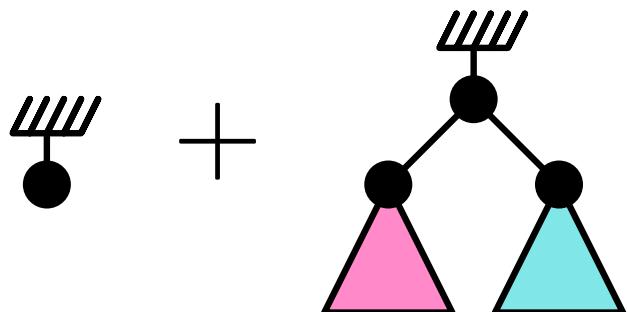
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Binary trees:



$$T(t) = t + tT(t)^2$$

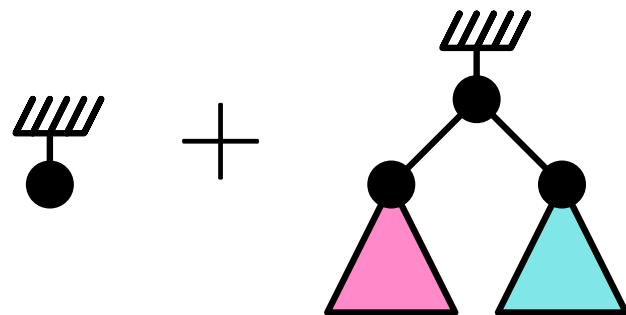
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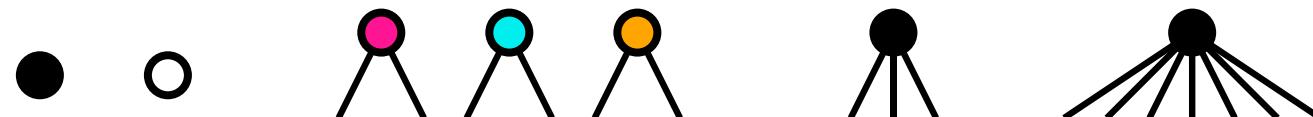
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Trees with nodes



$$T(t) = 2t + 3tT(t)^2 + tT(t)^3 + tT(t)^7$$

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What to do with it ?

- ▶ If we can solve the equation and get a closed form for $C(t)$, we can recover the c_n through $C^{(n)}(0) = n!c_n$.
- ▶ Use the equation to compute asymptotics of c_n .
- ▶ Build random generator (Boltzmann).

Generating functions

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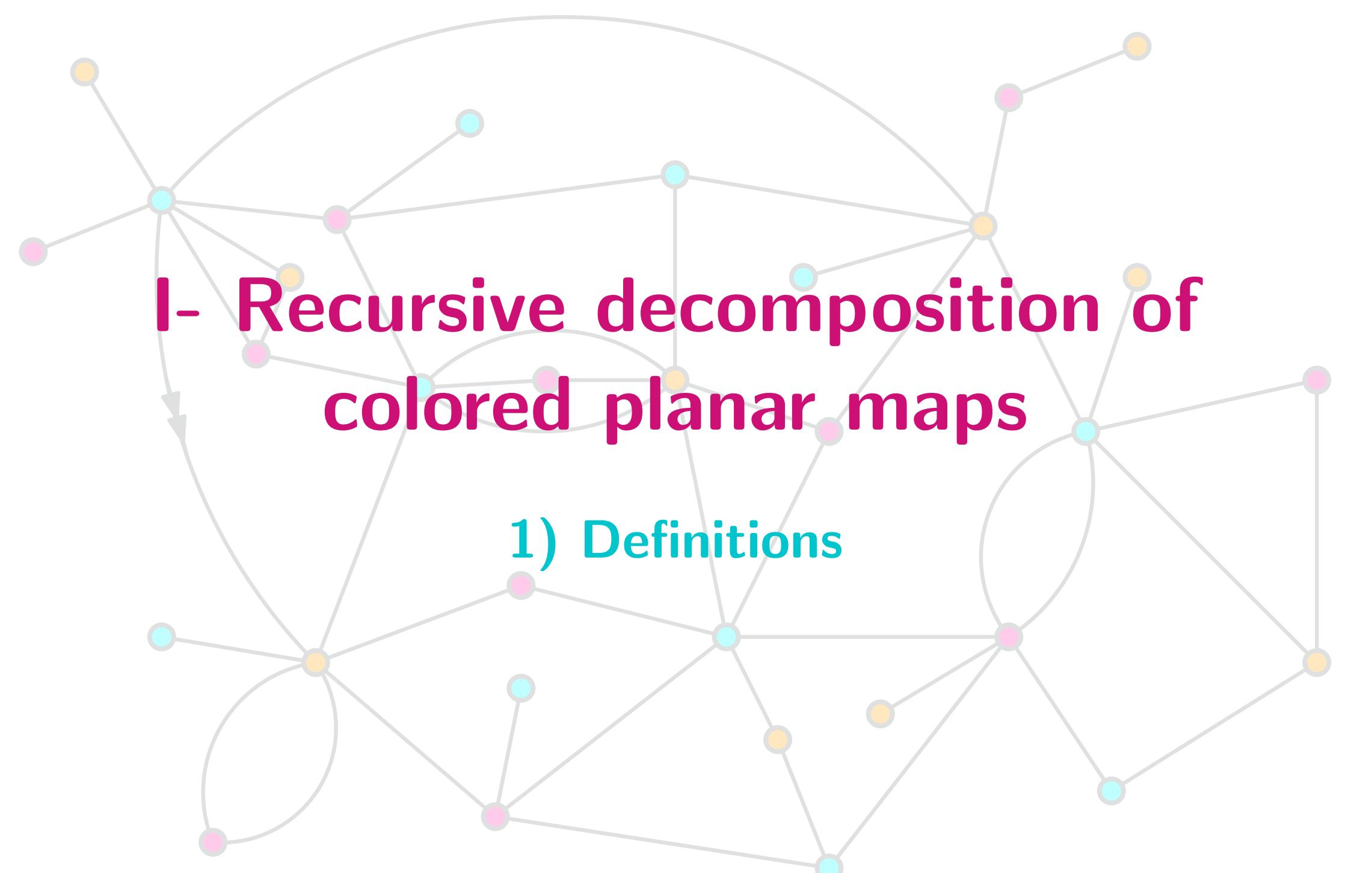
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The form of the equation gives information on the class:

- ▶ Rational function: $C(t) = P(t) + Q(t)C(t)$, P, Q polynomials.
 $\implies c_n \sim \kappa \lambda^n n^\alpha$, $\alpha \in \mathbb{N}$. $\mathcal{C} \approx$ regular language?
- ▶ Algebraic function: $P(t, C(t)) = 0$, $P(t, C)$ polynomial.
 $\implies c_n \sim \kappa \lambda^n n^\alpha$, $\alpha \in \mathbb{Q}$. $\mathcal{C} \approx$ non-ambiguous grammar (trees)?



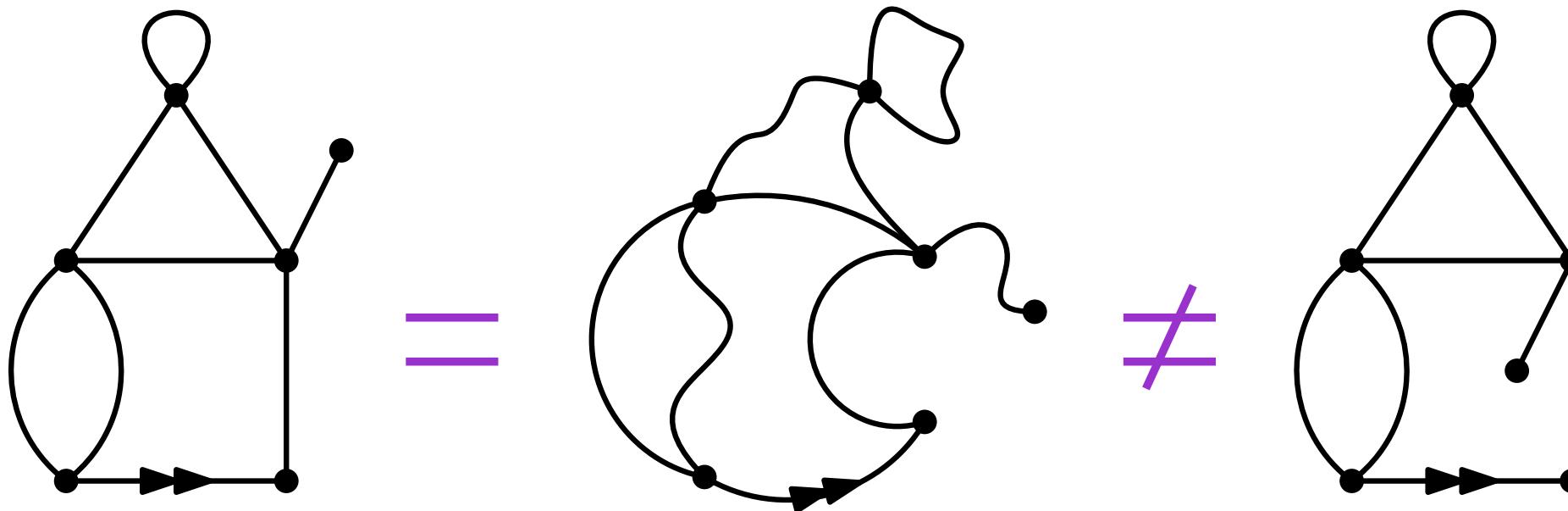
I- Recursive decomposition of colored planar maps

1) Definitions

Planar maps

A **planar map** is the embedding of a connected graph onto the sphere, up to orientation preserving homeomorphism.

Multi-edges and loops are allowed.



Planar map = planar graph + cyclic ordering of the edges around each vertex.

All maps are **rooted**, i.e. an oriented edge is marked.

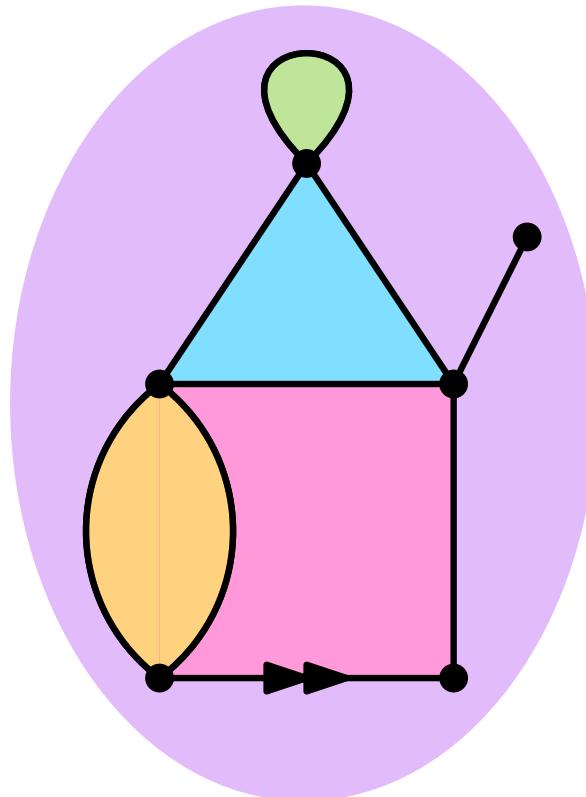
Planar maps

Planar map = planar graph + cyclic ordering of the edges around each vertex.

- **Vertices** and **edges** are inherited from the graph.
- **Faces** are the connected components of the sphere minus the map.

Size: number of edges n .

Here: 6 vertices, 9 edges
et 5 faces.



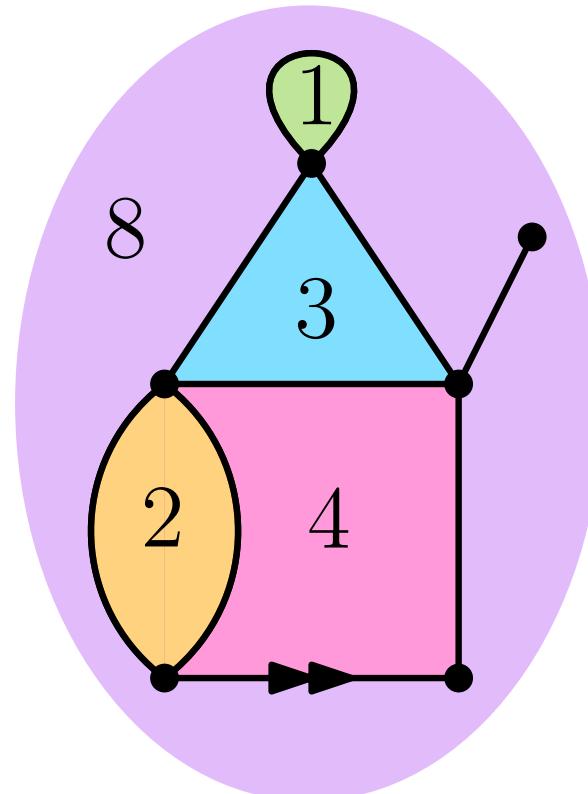
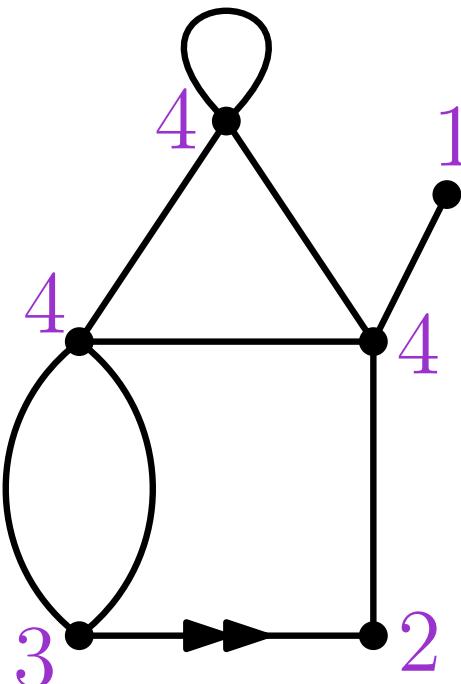
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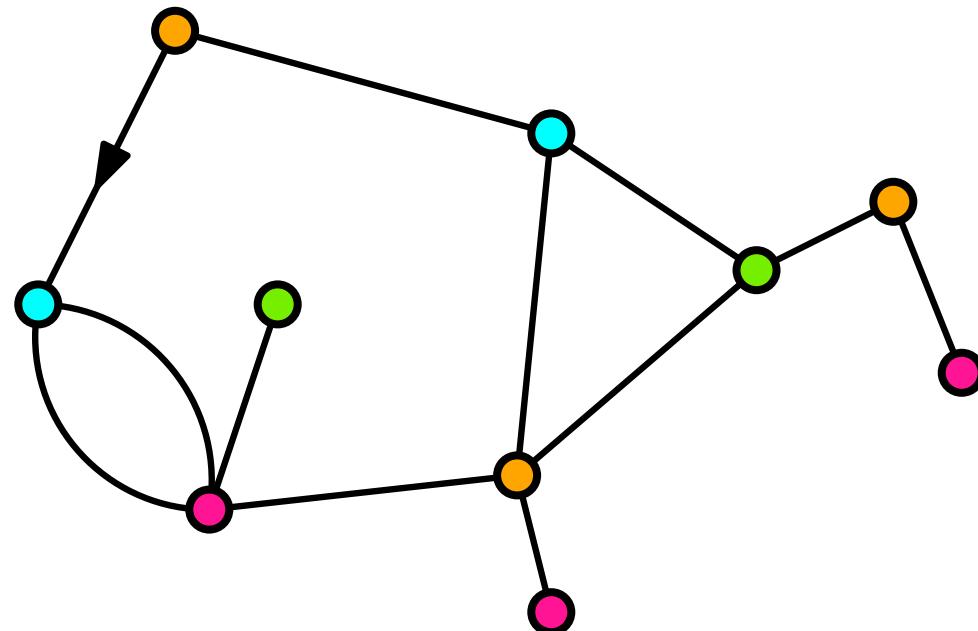
Size: number of edges n .

Degree (of a vertex or face) = number of incident half-edges.

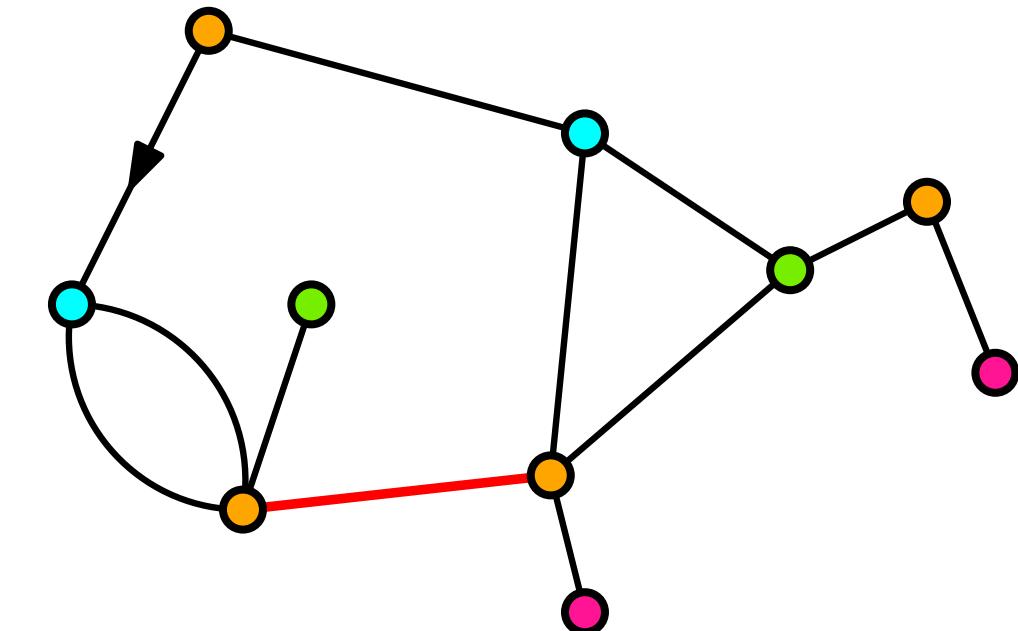


Colored maps

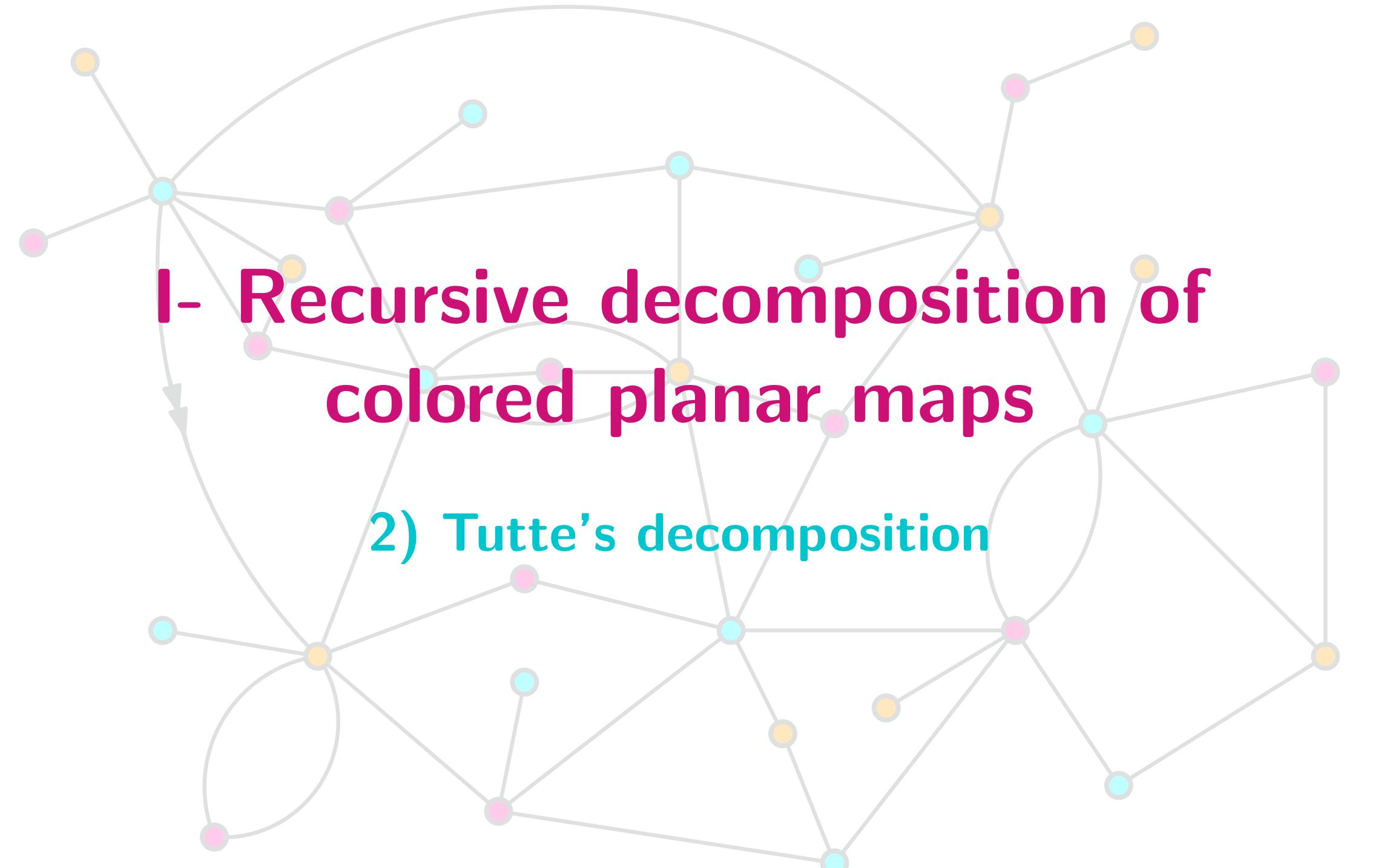
Colored map: assign a color to each vertex so that each pair of adjacent vertices receives different colors.



proper coloring



non-proper coloring

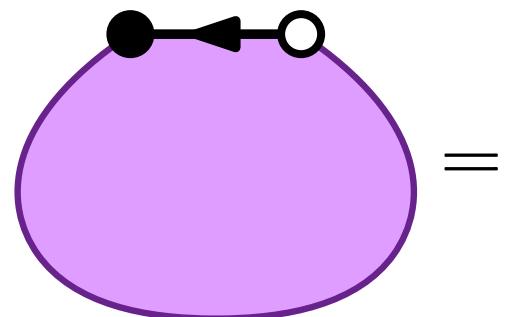


I- Recursive decomposition of colored planar maps

2) Tutte's decomposition

2-colored maps

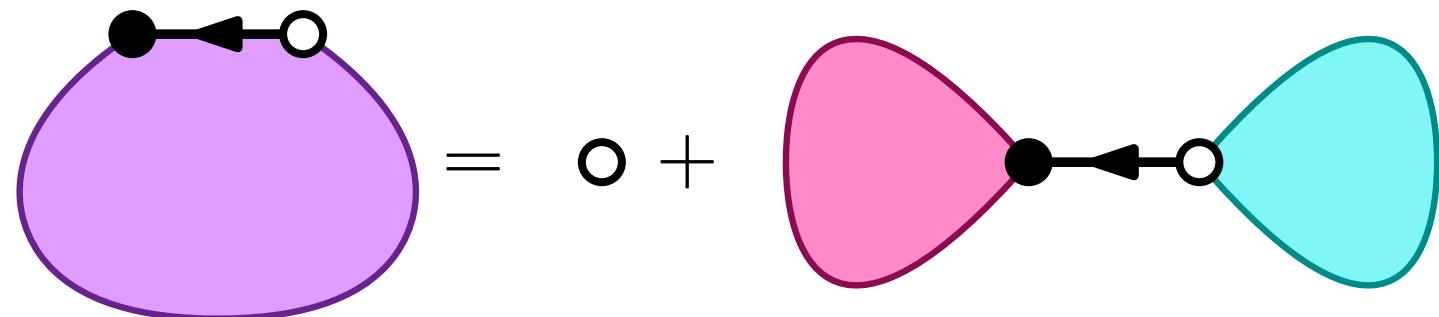
$B(t)$: generating function of bicolored planar maps counted by number of edges.



$$B(t) =$$

2-colored maps

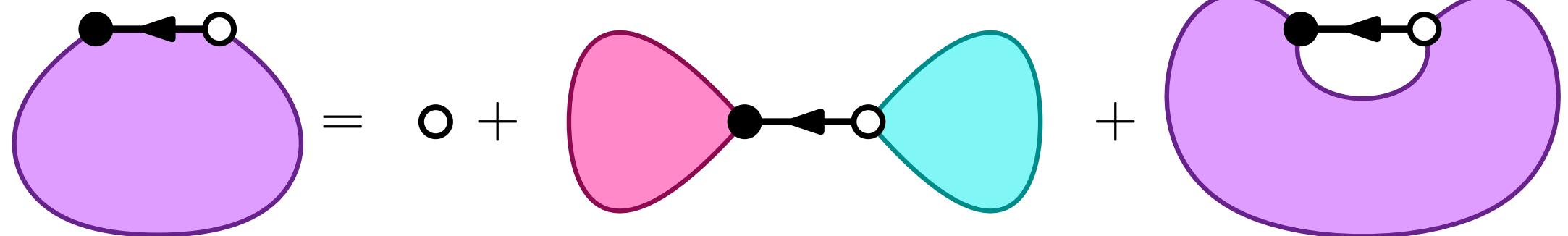
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$$B(t) = 1 + tB(t)^2$$

2-colored maps

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$$B(t) = 1 + tB(t)^2 + \text{??}$$

2-colored maps

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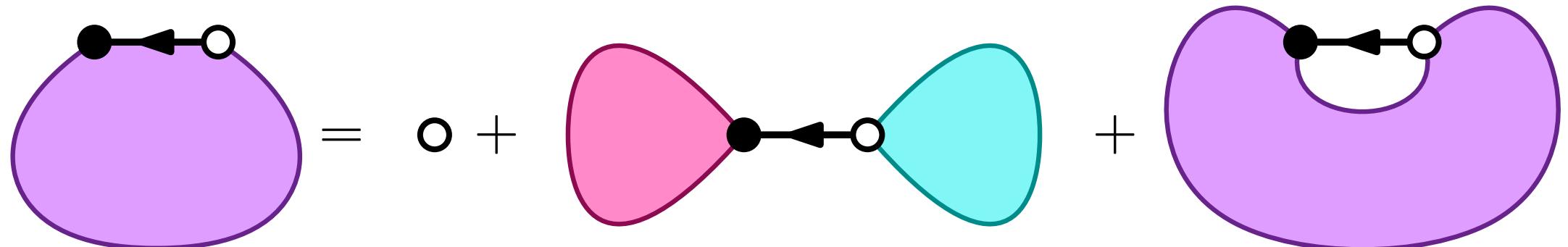
$$\text{Diagram: } \text{A purple circle with two black vertices connected by a curved edge.} = \text{O} + \text{Diagram: } \text{A purple oval with a pink oval on its left and a cyan oval on its right, connected by a black vertex with a curved edge.} + \text{Diagram: } \text{A purple oval with a black vertex on its left and a white vertex on its right, connected by a curved edge. The purple oval is surrounded by a dashed circle.} \\ \text{Equation: } B(t) = 1 + tB(t)^2 + \text{???}$$

2-colored maps

$B(t)$: generating function of bicolored planar maps counted by number of edges.

⇒ add a **catalytic variable** recording the half-length of the boundary

$$\hookrightarrow B(t, y) = \sum_{\mathfrak{b} \in \mathcal{B}} t^{e(\mathfrak{b})} y^{b(\mathfrak{b})}.$$



$$B(t, y) = 1 + t y B(t, y)^2 + t y \sum_{p \geq 1} \left(\sum_{k \geq 1}^p y^{p-k} \right) [y^p] B(t, y)$$

$$\hookrightarrow B(t, y) = 1 + t y B(t, y)^2 + t y \frac{B(t, y) - B(t, 1)}{y - 1}.$$

2-colored maps

$B(t, y)$: generating function of bicolored planar maps, $t \rightarrow$ edges, $y \rightarrow$ half degree of the outer face.

$$B(t, y) = 1 + tyB(t, y)^2 + ty \frac{B(t, y) - B(t, 1)}{y - 1}$$

→ Can be solved with the quadratic method [Brown '60s].

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Theorem. [Bousquet-Mélou Jehanne '06] If $F(y) \equiv F(t, y) \in \mathbb{Q}[y][[t]]$ satisfies an equation

$$F(y) = P(y) + tQ(t, y, F(y), \Delta F(y), \dots, \Delta^k F(y))$$

where P, Q are polynomials and $\Delta F(y) = \frac{F(y) - F(1)}{y - 1}$

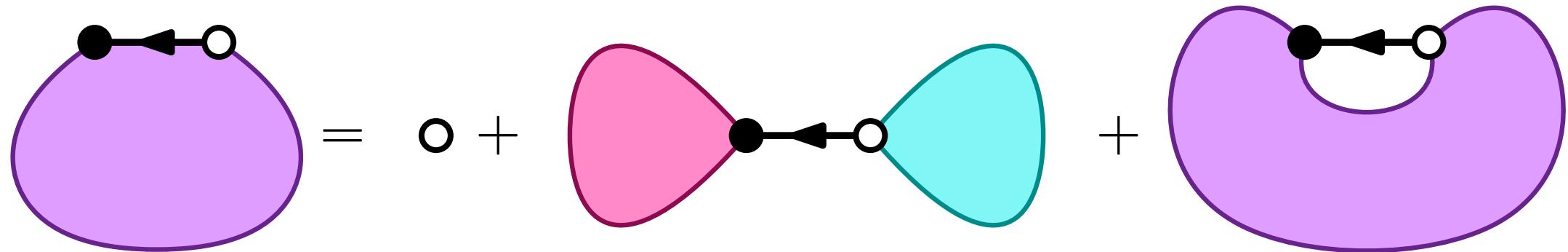
then $F(y)$ is an algebraic series.

→ $B(t, y)$ is algebraic.

+ Explicit bijection with a family of trees [Schaeffer '97].

3-colored maps

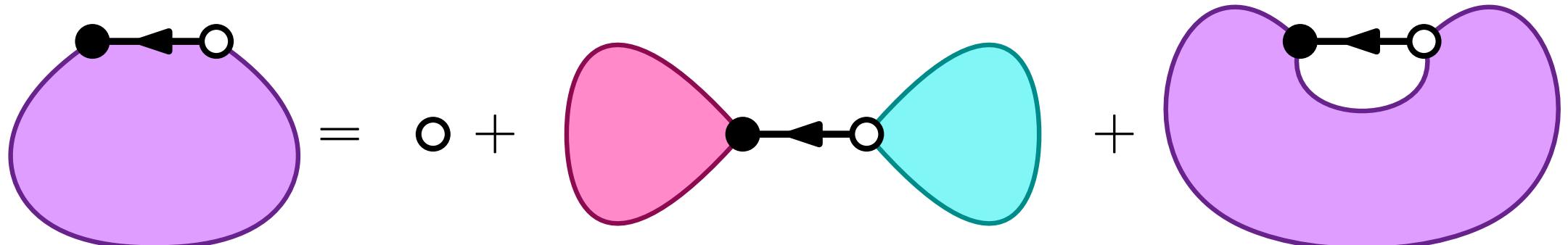
$M(t, y)$: generating function of 3-colored planar maps counted by number of edges (variable t) and degree of the outer face (variable y).



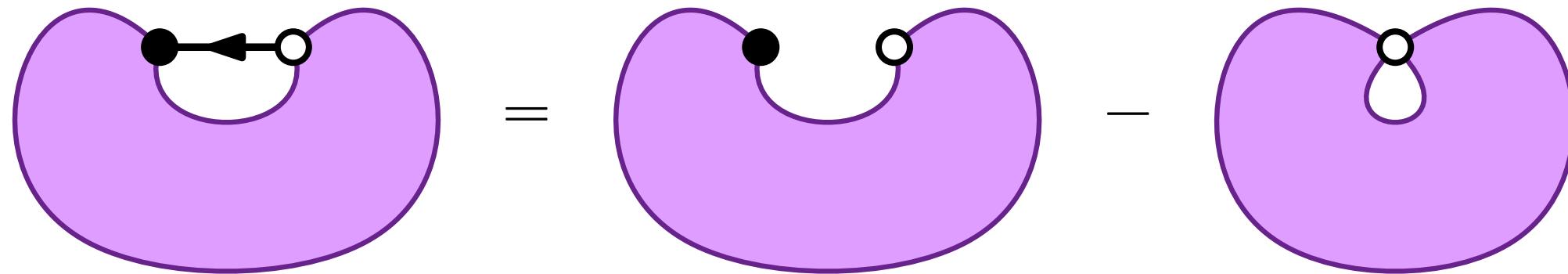
$$M(t, y) = 1 + 2ty^2 M(t, y)^2 + ty???$$

3-colored maps

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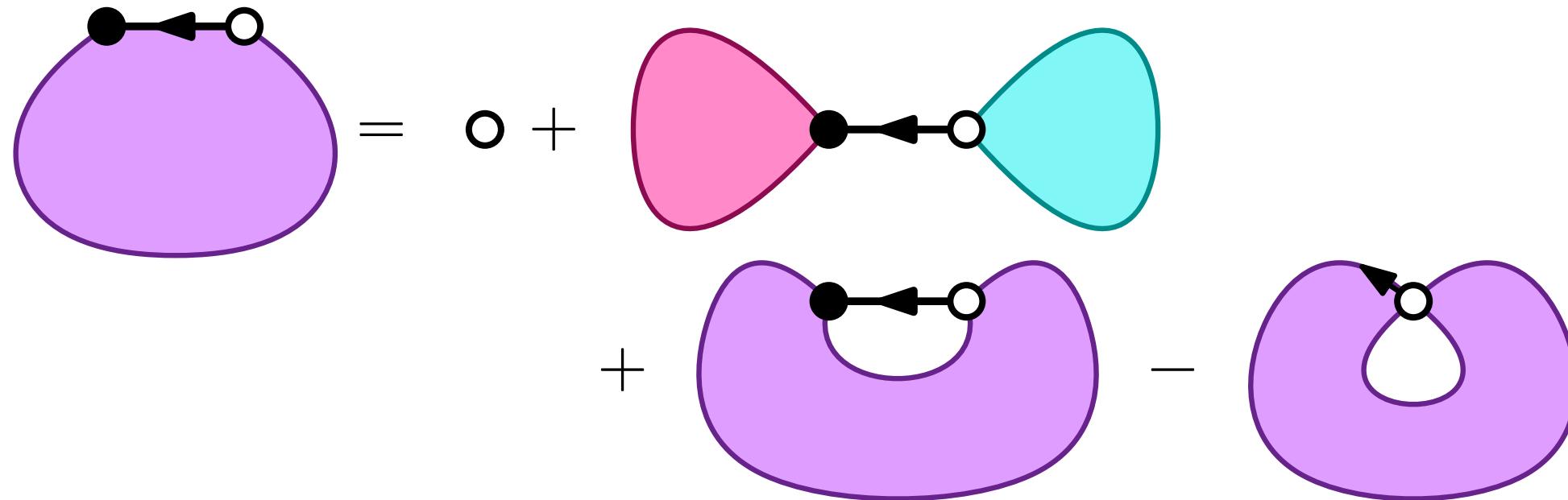
Solution:


$$= -$$

⇒ need a second catalytic variable recording the degree of the root vertex.

3-colored maps

$M(t, x, y)$: generating function of 3-colored planar maps counted by number of edges (variable t), degree of the outer face (variable y) and degree of the root vertex (variable x).



Theorem. [Bernardi Bousquet-Mélou '11]

$$\begin{aligned} M(x, y) = & 1 + xyt(2y - 1)M(x, y)M(1, y) - xytM(x, y)M(x, 1) \\ & - xyt \frac{xM(x, y) - M(1, y)}{x - 1} + xyt \frac{yM(x, y) - M(x, 1)}{y - 1}. \end{aligned}$$

3-colored maps

Equation for 3-colored planar maps:

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Not an algebraic/D-finite/... equation \implies no toolbox ?

We have 2 catalytic variables so Bousquet-Mélou and Jehanne's theorem does not apply.

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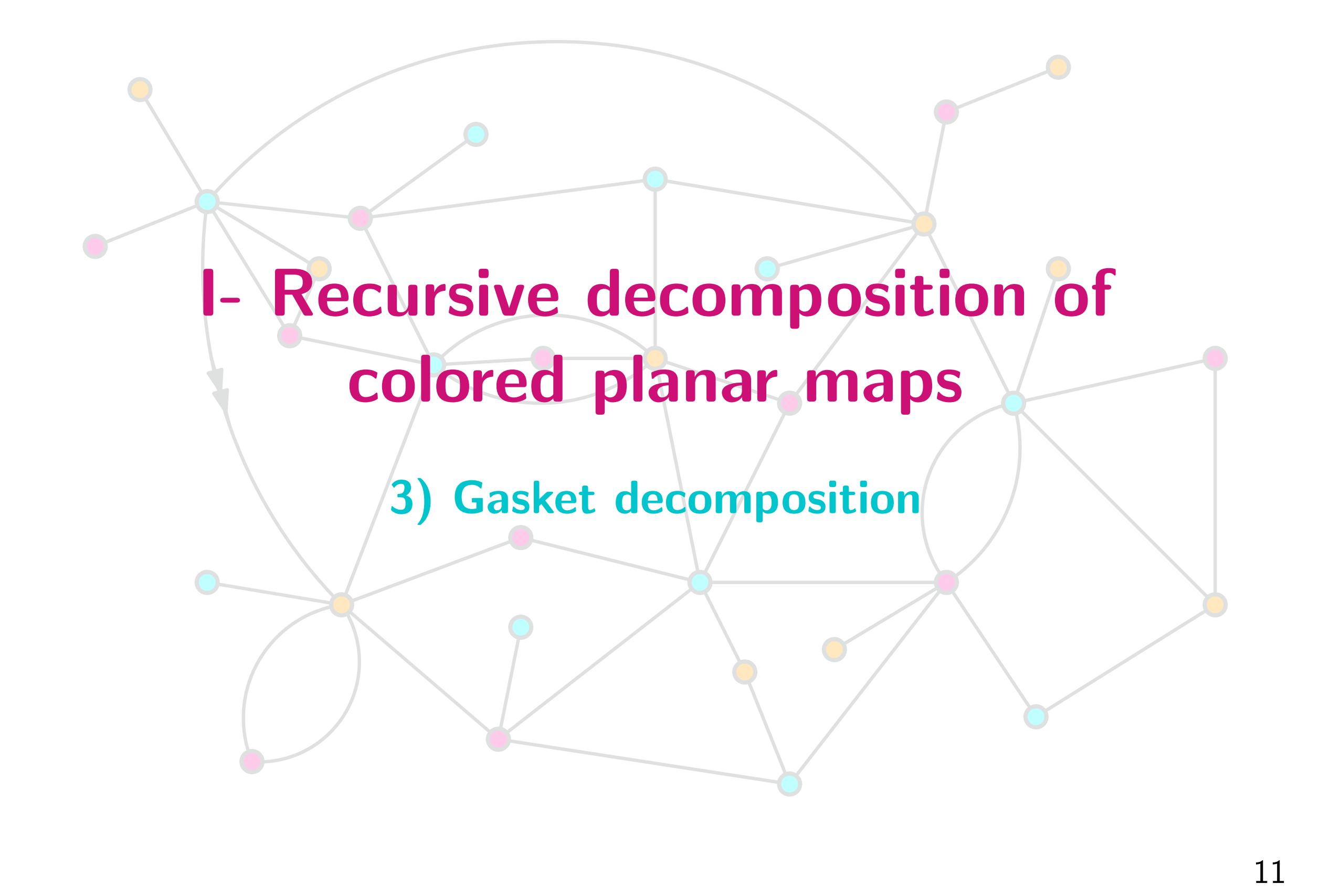
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Theorem. [Bernardi Bousquet-Mélou '11] The generating function of 3-colored planar map is algebraic and is equal to

$$M(1, 1) = \frac{(1 + 2S)(1 - 2S^2 - 4S^3 - 4S^4)}{(1 - 2S^3)^2}$$

with S the solution of $t = \frac{S(1 - S^3)}{(1 + 2S)^3}$ with constant term 0.

→ Obtained via a reduction to 1-catalytic equation.



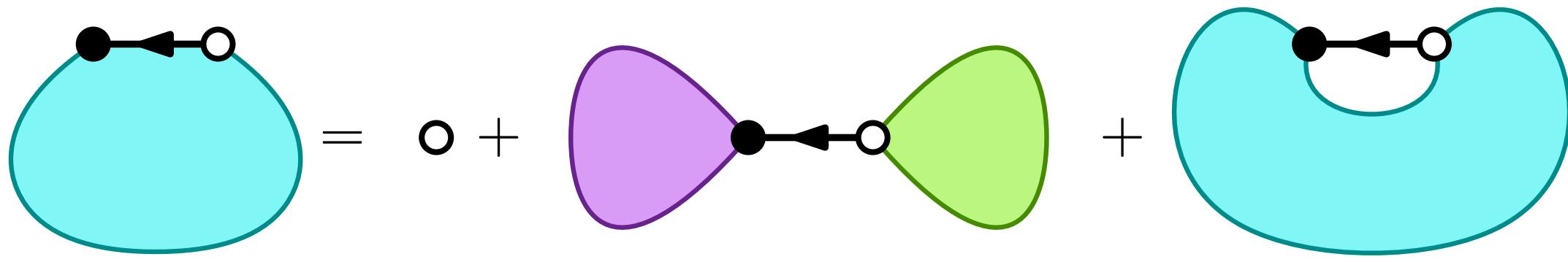
I- Recursive decomposition of colored planar maps

3) Gasket decomposition

Gasket decomposition

Idea: Use the structure from the coloring and get back to bipartite maps.

Let \mathcal{M} be the class of 3-colored planar maps with a **black and white outer face**. Let $M(t, y)$ be their generating function with t counting the edges and y half the degree of the outer face.

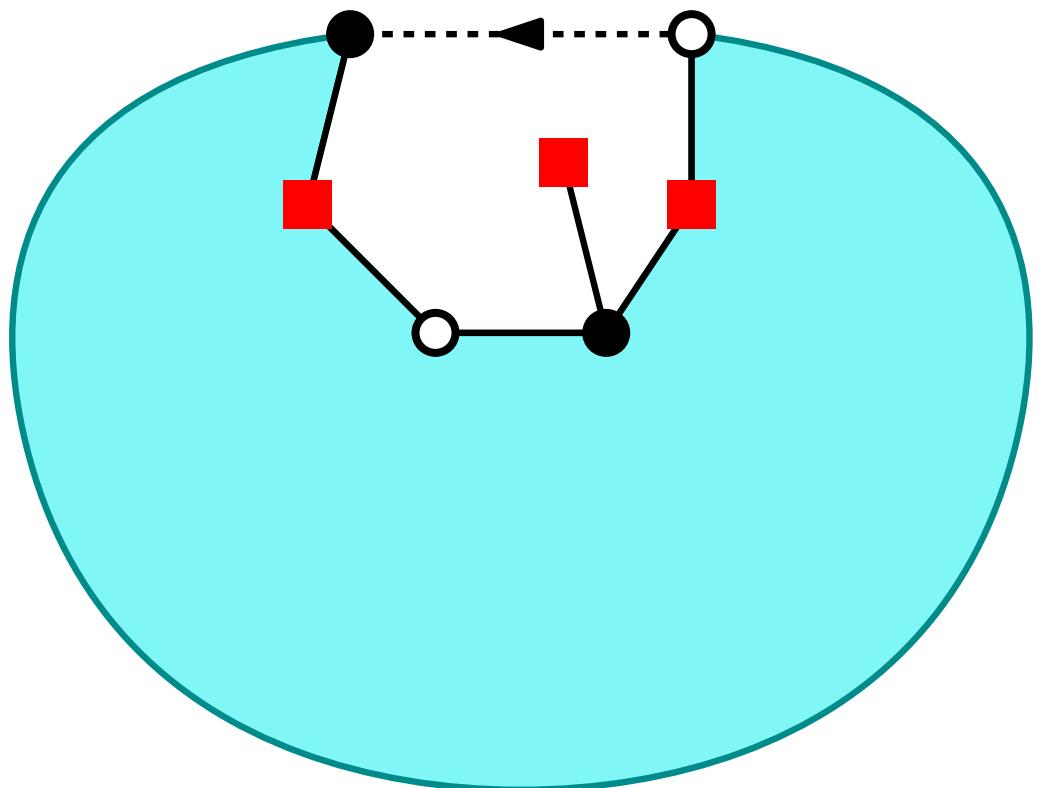


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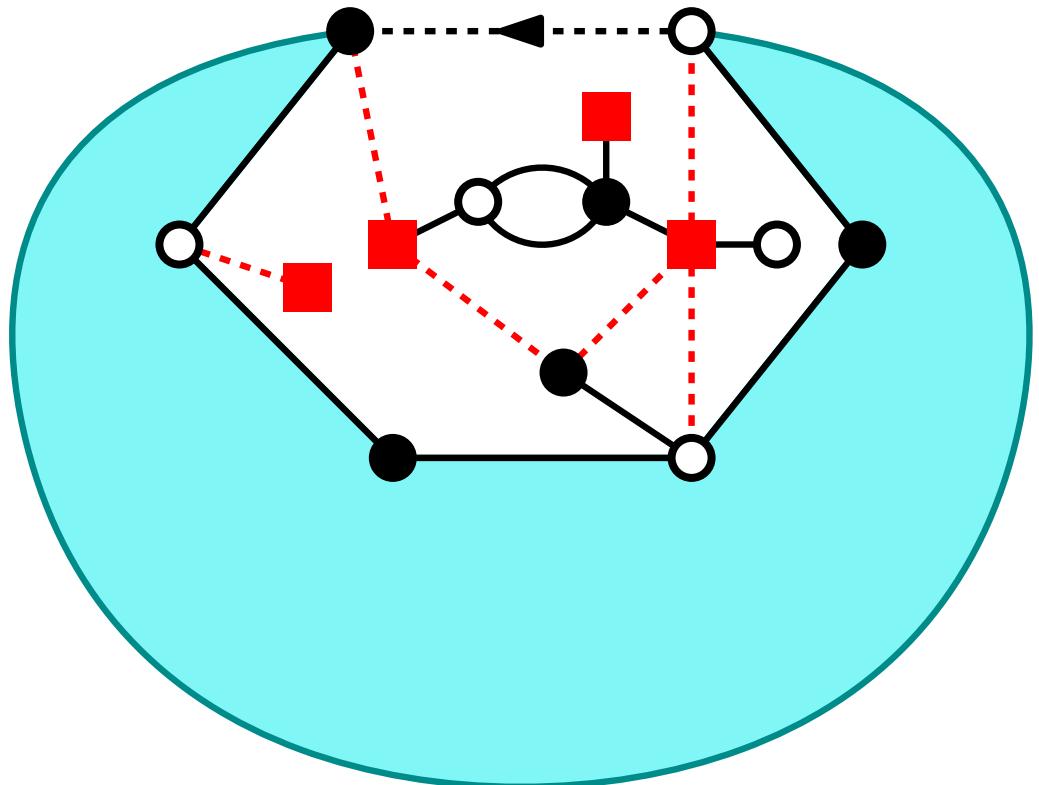
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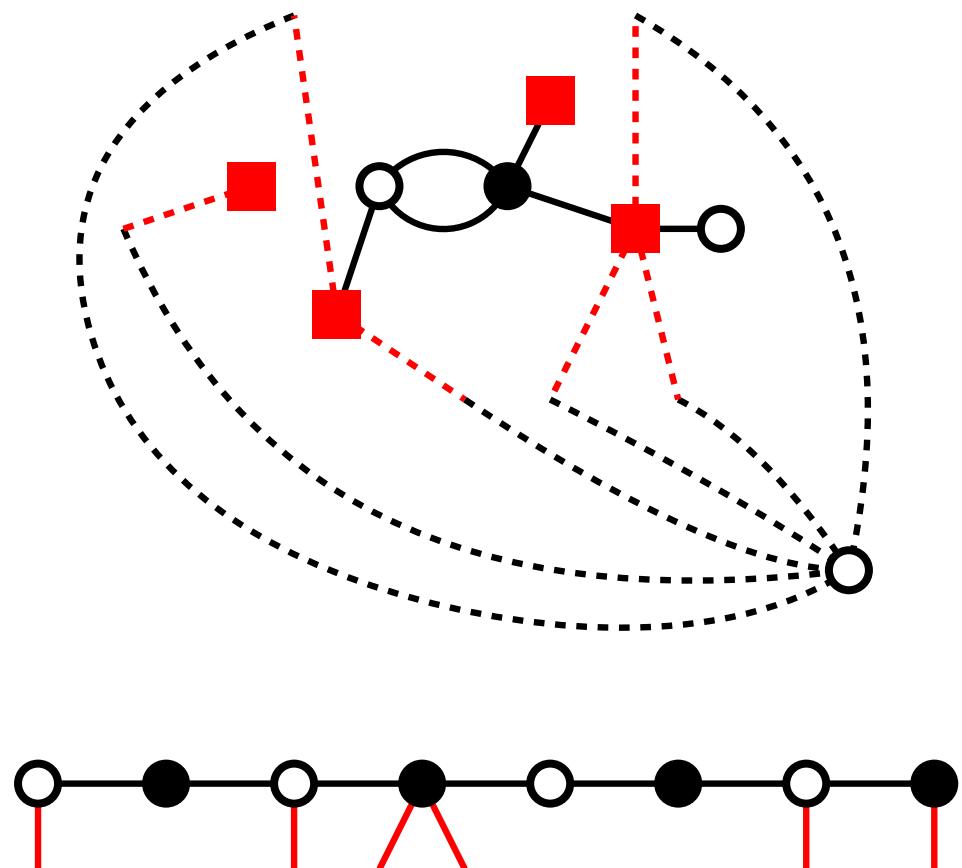
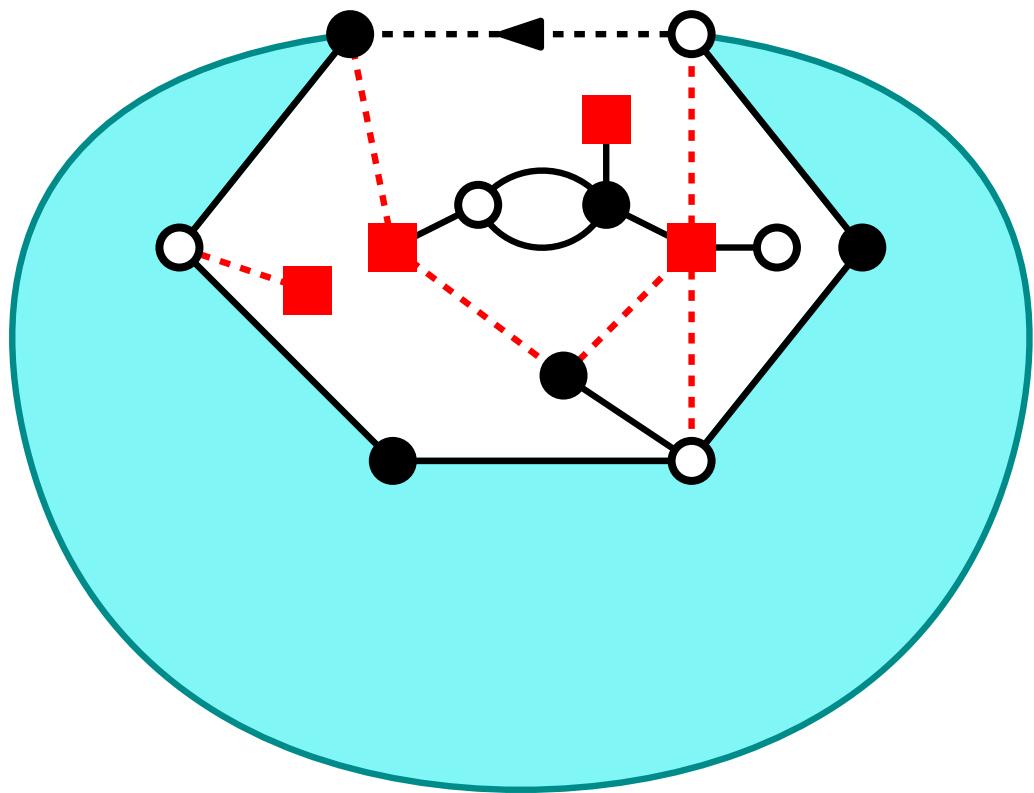
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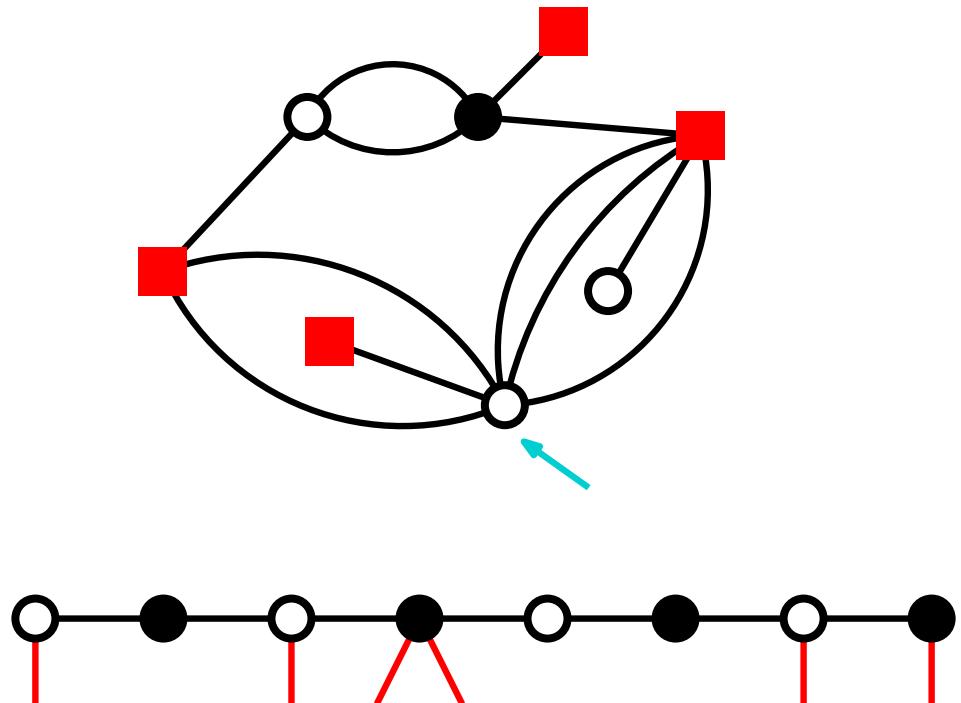
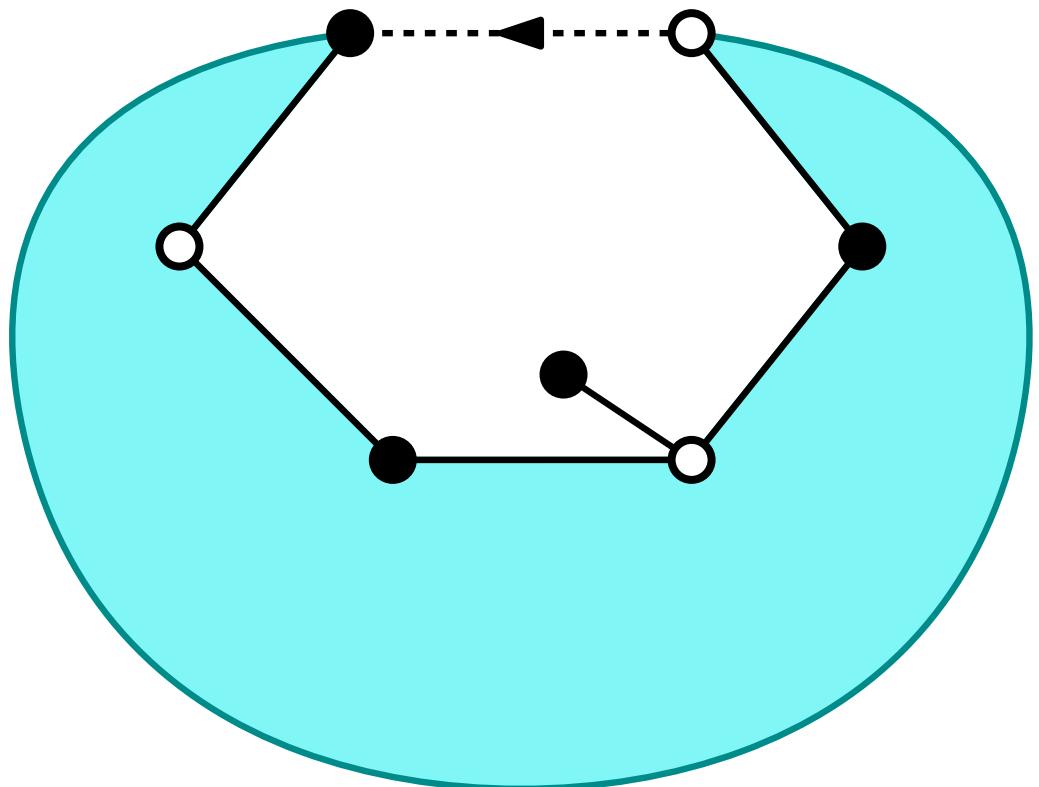
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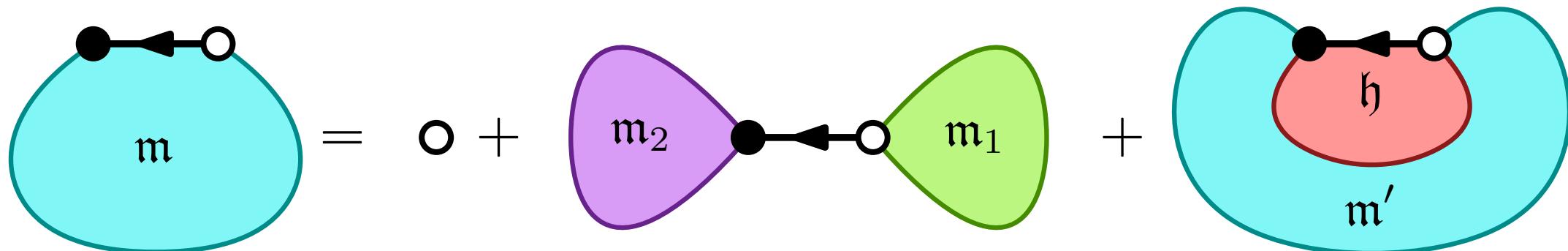
3-colored map where all neighbors of the root vertex are red.
+ connection caterpillar.

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Let \mathcal{H} be the class of 3-colored planar maps where all the **neighbors** of the **root vertex** are **red**. Let $H(t, x)$ be their generating function with t counting the edges and x the degree of the root vertex.

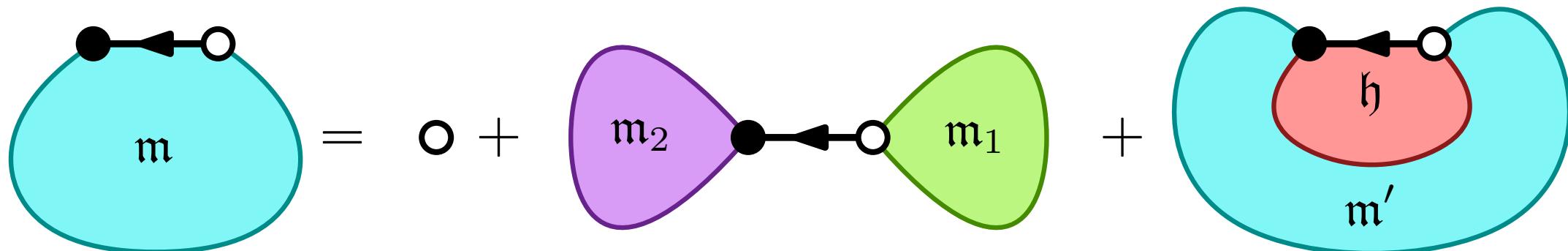


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Let \mathcal{H} be the class of 3-colored planar maps where all the **neighbors** of the **root vertex** are **red**. Let $H(t, x)$ be their generating function with t counting the edges and x the degree of the root vertex.



$$M(t, y) = 1 + ty \mathcal{M}(t, y)^2 + ty \sum_{p \geq 1} \left(\sum_{k \geq 1}^p y^{p-k} h_k \right) [y^p] M(t, y)$$

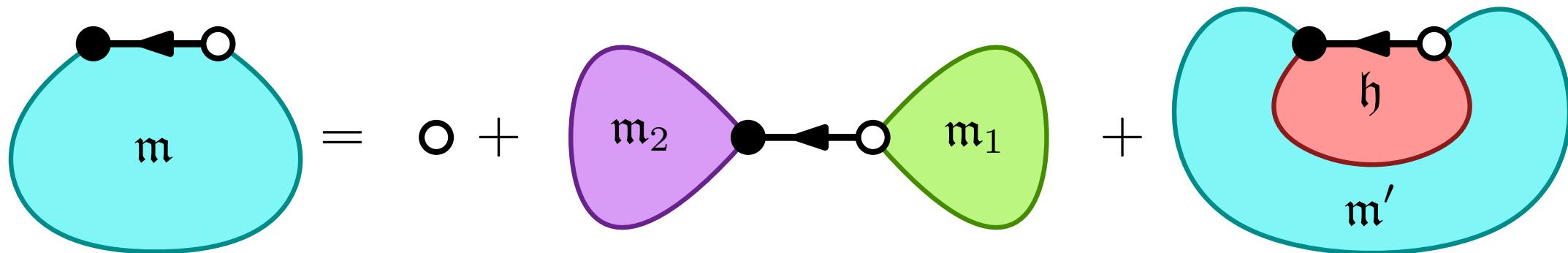
$$h_k = \sum_{d \geq 0} \binom{2k + d - 1}{2k - 1} [x^d] H(t, x)$$

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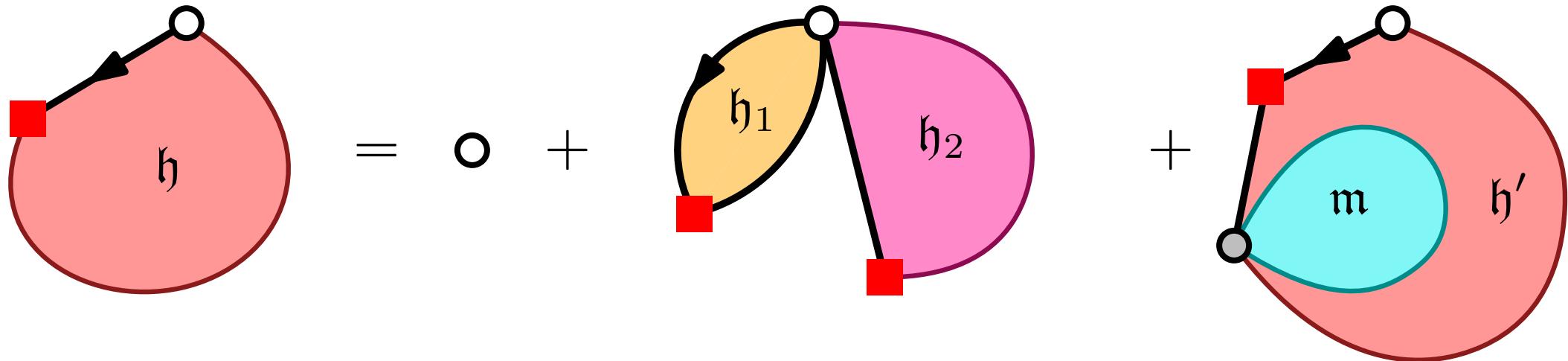
$$M(t, y) = 1 + ty\mathcal{M}(t, y)^2 + ty[y^{\geq 0}](\mathcal{M}(t, y) - 1)\tilde{H}(t, \bar{y}), \quad \bar{y} = 1/y$$

$$\tilde{H}(t, y) = \sum_{k \geq 1} y^k h_k = \frac{1}{2} \left[\frac{\sqrt{y}}{1 - \sqrt{y}} H\left(t, \frac{1}{1 - \sqrt{y}}\right) - \frac{\sqrt{y}}{1 + \sqrt{y}} H\left(t, \frac{1}{1 + \sqrt{y}}\right) \right].$$

Gasket decomposition

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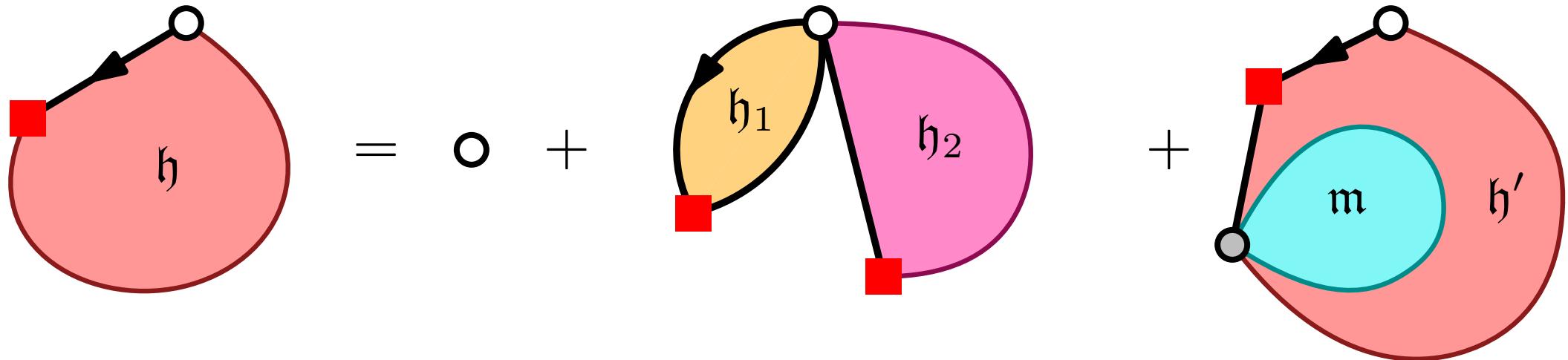
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Let \mathcal{H} be the class of 3-colored planar maps where all the **neighbors** of the **root vertex** are **red**. Let $H(t, x)$ be their generating function with t counting the edges and x the degree of the root vertex.



$$H(t, x) = 1 + txH(t, x)^2 + 2tx[x^{\geq 0}](H(t, x) - 1)\widetilde{M}(t, \bar{x}).$$

$$\widetilde{M}(t, x) = \sum_{i \geq 1} x^i \sum_{p \geq 0} \binom{2p + i - 1}{i - 1} [y^p] M(t, y) = \frac{x}{(1 - x)^2} M\left(t, \frac{1}{(1 - x)^2}\right),$$

Gasket decomposition

Let \mathcal{M} be the class of 3-colored planar maps with a **black and white outer face**.

Let \mathcal{H} be the class of 3-colored planar maps where all the **neighbors** of the **root vertex** are **red**. Let $H(t, x)$ be their generating function with t counting the edges and x the degree of the root vertex.

Theorem. [S. '26]

$$\begin{cases} M(t, y) = 1 + tyM(t, y)^2 + ty[y \geq 0](M(t, y) - 1)\tilde{H}(t, \bar{y}), \\ H(t, x) = 1 + txH(t, x)^2 + 2tx[x \geq 0](H(t, x) - 1)\tilde{M}(t, \bar{x}), \end{cases}$$

with

$$\tilde{M}(t, x) = \sum_{i \geq 1} x^i m_i = \frac{x}{(1-x)^2} M\left(t, \frac{1}{(1-x)^2}\right),$$

$$\tilde{H}(t, y) = \sum_{i \geq 1} y^i h_i = \frac{1}{2} \left[\frac{\sqrt{y}}{1-\sqrt{y}} H\left(t, \frac{1}{1-\sqrt{y}}\right) - \frac{\sqrt{y}}{1+\sqrt{y}} H\left(t, \frac{1}{1+\sqrt{y}}\right) \right].$$

$$P_M(S, y, M) = 4y^5(y-1)^2S^5(2S^3-1)^6M^6 + 24y^4(y-1)^2(1+2S)^3(2S^3-1)^5S^4M^5 - (1+2S)(2S^3-1)^4((64S^6+96S^5+88S^4+32S^3-2S+1)y^2 - (16S^4+456S^3+684S^2+334S+57)(1+2S)^2y + 56(1+2S)^5)S^3(y-1)y^3M^4 - 4(2S^3-1)^3(1+2S)^4((64S^6+96S^5+88S^4+32S^3-2S+1)y^2 - (16S^4+136S^3+204S^2+94S+17)(1+2S)^2y + 16(1+2S)^5)S^2(y-1)y^2M^3 - (1+2S)^2(2S^3-1)^2(2S(64S^6+96S^5+94S^4+50S^3+18S^2+4S+1)(S+1)^3y^4 + (1+2S)(6144S^{10}+20928S^9+33600S^8+32064S^7+18380S^6+5244S^5-162S^4-536S^3-87S^2+18S+5)y^3 - 2(384S^6+1120S^5+1914S^4+1604S^3+706S^2+171S+23)(1+2S)^5y^2 + (96S^4+616S^3+924S^2+414S+77)(1+2S)^7y - 36(1+2S)^{10})SyM^2 - 2(2S^3-1)(1+2S)^5(2S(64S^6+96S^5+94S^4+50S^3+18S^2+4S+1)(S+1)^3y^4 + (1+2S)(2048S^{10}+6592S^9+9536S^8+7488S^7+25008S^6+18364S^5+6000S^4-11894-6S+1)y^3 - 2(128S^6+224S^5+266S^4+164S^3+50S^2+11S+3)(1+2S)^5y^2 + (32S^4+72S^3+108S^2+38S+9)(1+2S)^7y - 4(1+2S)^{10})M - (1+2S)^3(S(4S^2+2S+1)^3(S+1)^6y^4 + 2(64S^6+96S^5+94S^4+50S^3+18S^2+4S+1)(S+1)^3(1+2S)^5y^3 + (1024S^9+3008S^8+3520S^7+1344S^6-1460S^5-2436S^4-1642S^3-616S^2-127S-12)(1+2S)^6y^2 - 2(64S^5+64S^4+14S^3-36S^2-34S-9)(1+2S)^{10}y + 8(2S^3-1)(1+2S)^{12})$$

$$P_H(S, \sqrt{\Delta}, x, H) = S^3x^5(2S^3-1)^4(2x-1)H^4 - 2S^2x^3(2x-1)(x-2)(1+2S)^3(2S^3-1)^3H^3 + Sx(1+2S)(2S^3-1)^2(2S(2S^3+4S^2+2S+1)x^4 - 2(8S^4+8S^3+12S^2+2S+1)(1+2S)^2x^3 + (8S^4-24S^3-36S^2-22S-3)(1+2S)^2x^2 + 10(1+2S)^5x - 4(1+2S)^5)H^2 - 2(1+2S)(2S^3-1)(S((8S^5+28S^4+38S^3+26S^2+10S+2)\sqrt{\Delta} + 16S^6+56S^5+76S^4+58S^3+28S^2+8S+1)x^4 - 2(16S^6+48S^5+86S^4+76S^3+34S^2+9S+1)(1+2S)^3x^3 + (20S^4+56S^3+84S^2+32S+7)(1+2S)^5x^2 - (8S^4+56S^3+84S^2+38S+7)(1+2S)^5x + 2(1+2S)^8)H(1+2S)^3((32S^8+(16S^6+64S^5+104S^4+90S^3+46S^2+14S+2)\sqrt{\Delta} + 176S^7+384S^6+452S^5+356S^4+182S^3+64S^2+15S+2)x^3 - 2(1+2S)(-64S^9-320S^8-656S^7+(8S^5+28S^4+38\sqrt{\Delta}S^3+26S^2+10S+2)\sqrt{\Delta} - 560S^6-72S^5+332S^4+354S^3+172S^2+44S+5)x^2 - 2(8S^7+96S^6+144S^5+60S^4-46S^3-80S^2-38S-7)(1+2S)^3x + 4(2S^3-1)(1+2S)^6)$$

III- Algebraicity proof

The strategy: Guess and Check

$$\begin{cases} M(t, y) = 1 + tyM(t, y)^2 + ty[y \geq 0](M(t, y) - 1)\tilde{H}(t, \bar{y}), \\ H(t, x) = 1 + txH(t, x)^2 + 2tx[x \geq 0](H(t, x) - 1)\tilde{M}(t, \bar{x}). \end{cases}$$

Idea: The equations allow us to recursively compute all the coefficients of the series M and H
⇒ unique solution in $\mathbb{Q}[y][[t]] \times \mathbb{Q}[x][[t]]$.

Guess and Check:

1. Guess polynomial equations satisfied by our series.
2. Prove that the solutions of the polynomial equations satisfy the original equations.

Guess

$$\begin{cases} M(t, y) = 1 + tyM(t, y)^2 + ty[y^{\geq 0}](M(t, y) - 1)\tilde{H}(t, \bar{y}), \\ H(t, x) = 1 + txH(t, x)^2 + 2tx[x^{\geq 0}](H(t, x) - 1)\tilde{M}(t, \bar{x}). \end{cases}$$

Guess:

1. Use the equations to compute the first (250) terms of the series M and H .
2. Guess polynomial equations with `algeqtoseries` (package gfun).
3. Determine which solution of the polynomial equation is the candidate.

→ We use the intermediate field $\mathbb{Q}(t, y) \rightarrow \mathbb{Q}(S, y) \rightarrow \mathbb{Q}(t, y, M(t, y))$

where $t = \frac{S(1 - S^3)}{(1 + 2S)^3}$ to have smaller equations.

→ Equations of degree 6 for $M(t, y)$ over $\mathbb{Q}(S, y)$ and of degree 8 for $H(t, x)$ over $\mathbb{Q}(S, x)$.

Guess

1. Use the equations to compute the first (250) terms of the series M and H .
2. Guess polynomial equations with `algeqtoseries` (package gfun).
3. Determine which solution of the polynomial equation is the candidate.

→ Equations of degree 6 for $M(t, y)$ over $\mathbb{Q}(S, y)$ and of degree 8 for $H(t, x)$ over $\mathbb{Q}(S, x)$.

$$P_M(S, y, M) = 4y^5(y-1)^2S^5(2S^3-1)^6M^6 + 24y^4(y-1)^2(1+2S)^3(2S^3-1)^5S^4M^5 - (1+2S)(2S^3-1)^4((64S^6+96S^5+88S^4+32S^3-2S+1)y^2 - (16S^4+456S^3+684S^2+334S+57)(1+2S)^2y + 56(1+2S)^5)S^3(y-1)y^3M^4 - 4(2S^3-1)^3(1+2S)^4((64S^6+96S^5+88S^4+32S^3-2S+1)y^2 - (16S^4+136S^3+204S^2+94S+17)(1+2S)^2y + 16(1+2S)^5)S^2(y-1)y^2M^3 - (1+2S)^2(2S^3-1)^2(2S(64S^6+96S^5+94S^4+50S^3+18S^2+4S+1)(S+1)^3y^4 + (1+2S)(6144S^{10}+20928S^9+33600S^8+32064S^7+18380S^6+5244S^5-162S^4-536S^3-87S^2+18S+5)y^3 - 2(384S^6+1120S^5+1914S^4+1604S^3+706S^2+171S+23)(1+2S)^5y^2 + (96S^4+616S^3+924S^2+414S+77)(1+2S)^7y - 36(1+2S)^{10})SyM^2 - 2(2S^3-1)(1+2S)^5(2S(64S^6+96S^5+94S^4+50S^3+18S^2+4S+1)(S+1)^3y^4 + (1+2S)(2048S^{10}+6592S^9+9536S^8+7488S^7+2508S^6-900S^5-1346S^4-600S^3-119S^2-6S+1)y^3 - 2(128S^6+224S^5+266S^4+164S^3+50S^2+11S+3)(1+2S)^5y^2 + (32S^4+72S^3+108S^2+38S+9)(1+2S)^7y - 4(1+2S)^{10})M - (1+2S)^3(S(4S^2+2S+1)^3(S+1)^6y^4 + 2(64S^6+96S^5+94S^4+50S^3+18S^2+4S+1)(S+1)^3(1+2S)^5y^3 + (1024S^9+3008S^8+3520S^7+1344S^6-1460S^5-2436S^4-1642S^3-616S^2-127S-12)(1+2S)^6y^2 - 2(64S^5+64S^4+14S^3-36S^2-34S-9)(1+2S)^{10}y + 8(2S^3-1)(1+2S)^{12})$$

Guess

1. Use the equations to compute the first (250) terms of the series M and H .
2. Guess polynomial equations with `algeqtoseries` (package gfun).
3. Determine which solution of the polynomial equation is the candidate.

→ Equations of degree 6 for $M(t, y)$ over $\mathbb{Q}(S, y)$ and of degree 8 for $H(t, x)$ over $\mathbb{Q}(S, x)$.

Both polynomial equations have only one solution that is a power series in t and they both have polynomial coefficients.

→ We denote them M_c and H_c .

Check

Check: Prove that M_c and H_c satisfy

$$\begin{cases} M_c(t, y) = 1 + ty M_c(t, y)^2 + ty[y \geq 0] (M_c(t, y) - 1) \tilde{H}_c(t, \bar{y}), \\ H_c(t, x) = 1 + tx H_c(t, x)^2 + 2tx[x \geq 0] (H_c(t, x) - 1) \tilde{M}_c(t, \bar{x}), \end{cases}$$

with $\tilde{M}_c(t, x) = \frac{x}{(1-x)^2} M_c\left(t, \frac{1}{(1-x)^2}\right)$

and $\tilde{H}_c(t, y) = \frac{1}{2} \left[\frac{\sqrt{y}}{1-\sqrt{y}} H_c\left(t, \frac{1}{1-\sqrt{y}}\right) - \frac{\sqrt{y}}{1+\sqrt{y}} H_c\left(t, \frac{1}{1+\sqrt{y}}\right) \right]$.

Check

Check: Prove that M_c and H_c satisfy

$$\begin{cases} M_c(t, y) = 1 + ty M_c(t, y)^2 + ty[y \geq 0] (M_c(t, y) - 1) \tilde{H}_c(t, \bar{y}), \\ H_c(t, x) = 1 + tx H_c(t, x)^2 + 2tx[x \geq 0] (H_c(t, x) - 1) \tilde{M}_c(t, \bar{x}), \end{cases}$$

with $\tilde{M}_c(t, x) = \frac{x}{(1-x)^2} M_c\left(t, \frac{1}{(1-x)^2}\right)$

and $\tilde{H}_c(t, y) = \frac{1}{2} \left[\frac{\sqrt{y}}{1-\sqrt{y}} H_c\left(t, \frac{1}{1-\sqrt{y}}\right) - \frac{\sqrt{y}}{1+\sqrt{y}} H_c\left(t, \frac{1}{1+\sqrt{y}}\right) \right]$.

1. Reformulate the equations:

$$\begin{cases} [y \geq 0] R_M(y) = 0, \\ [x \geq 0] R_H(x) = 0, \end{cases}$$

where $R_M(y) = (M_c(y) - 1) \tilde{H}_c(\bar{y}) - \frac{1}{ty} (M_c(y) - 1 - ty M_c(y)^2) \in \mathbb{Q}((\bar{y}))[[t]]$

and $R_H(x) = (H_c(x) - 1) \tilde{M}_c(\bar{x}) - \frac{1}{2tx} (H_c(x) - 1 - tx H_c(x)^2) \in \mathbb{Q}((\bar{x}))[[t]]$.

Check

Check: Prove that $[y^{\geq 0}]R_M(y) = [x^{\geq 0}]R_H(x) = 0$.

2. Compute minimal polynomials for R_M and R_H by eliminating M_c and H_c using resultants.

→ An equation of degree 12 for R_M over $\mathbb{Q}(S, y)$ and one of degree 12 for R_H over $\mathbb{Q}(S, x)$.

Unfortunately, the equations have solutions with positive powers

⇒ no “all solutions are negative” result

⇒ getting something on R will be hard.

Check

Check: Prove that $[y^{\geq 0}]R_M(y) = [x^{\geq 0}]R_H(x) = 0$.

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Unfortunately, the equations have solutions with positive powers

⇒ no “all solutions are negative” result

⇒ getting something on R will be hard.

3. Find intermediate field extensions.

$$\mathbb{Q}(S, y) \xrightarrow{3} \mathbb{Q}(S, L) \xrightarrow{2} \mathbb{Q}(S, K) \xrightarrow{2} \mathbb{Q}(S, y, R_M)$$

$$L = -\frac{1}{y} \frac{(1+2S)^3(S+2S^2+4S^3-L)(4S^3-L)}{S(1+2S+2S^2+4S^3-L)^2},$$

$$K = \frac{L}{4S(4S^2+2S+1)(1-K)}.$$

Check

Check: Prove that $[y^{\leq 0}]R_M(y) = [x^{\leq 0}]R_H(x) = 0$.

3. Find intermediate field extensions.

$$\mathbb{Q}(S, y) \xrightarrow{3} \mathbb{Q}(S, L) \xrightarrow{2} \mathbb{Q}(S, K) \xrightarrow{2} \mathbb{Q}(S, y, R_M)$$

$$L = -\frac{1}{y} \frac{(1+2S)^3(S+2S^2+4S^3-L)(4S^3-L)}{S(1+2S+2S^2+4S^3-L)^2},$$

$$K = \frac{L}{4S(4S^2+2S+1)(1-K)}.$$

4. Solve the equation for R_M over $\mathbb{Q}(S, K)$.

→ $R_M(y) = \frac{B - C\sqrt{D}}{2A}$ with $A, B, C, D \in \mathbb{Q}[S, K]$.

5. Prove that L then K then R_M is y -negative.

6. Apply the same method to R_H .

Algebraicity

a.

$$\begin{cases} M(t, y) = 1 + ty \textcolor{teal}{M}(t, y)^2 + ty[y^{\geq 0}] (\textcolor{teal}{M}(t, y) - 1) \tilde{H}(t, \bar{y}), \\ H(t, x) = 1 + tx \textcolor{red}{H}(t, x)^2 + 2tx[x^{\geq 0}] (\textcolor{red}{H}(t, x) - 1) \tilde{M}(t, \bar{x}). \end{cases}$$

has a unique solution in $\mathbb{Q}[y][[t]] \times \mathbb{Q}[x][[t]]$.

b. There are polynomials $P_M(S, y, M)$ and $P_H(S, x, H)$ whose only power series solution satisfy the system.

Theorem. [S. '26] The generating function $M(t, y)$ of 3-colored planar maps with black and white border is algebraic over $\mathbb{Q}(t, y)$ of degree 24.

- ▶ Complete map of the subfields of $\mathbb{Q}(t, y, M(y))$.

$$\mathbb{Q}(t, y) \xrightarrow{4} \mathbb{Q}(S, y) \xrightarrow{3} \mathbb{Q}(S, L) \xrightarrow{2} \mathbb{Q}(S, y, M(y))$$

- ▶ All $[y^p]M(S, y)$ are rational fractions in S .

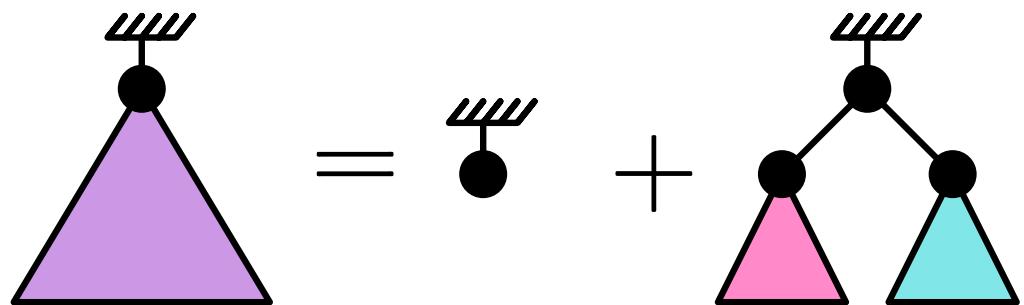


[Duchon, Flajolet,
Louchard, Schaeffer '04]

Exact size sampling

Principle: use the recursive decomposition to sample uniformly objects of a given size n .

For binary trees :



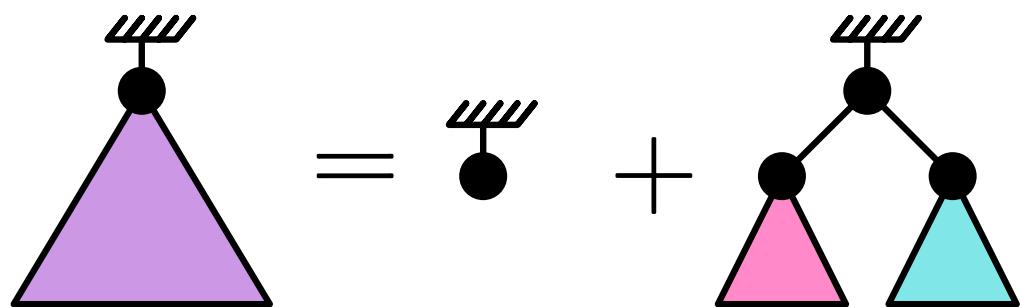
$$t_1 = 1$$
$$t_n = \sum_{k=1}^{n-1} t_k t_{n-k-1}$$

```
BinTreeSampler(n) =  
if n = 1 then •  
else draw k with  $\mathbb{P}(k) = t_k t_{n-k-1} / t_n$   
(BinTreeSampler(k), BinTreeSampler(n-1-k))
```

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```
BinTreeSampler(n) =  
  if n = 1 then •  
  else draw k with  $\mathbb{P}(k) = t_k t_{n-k-1}/t_n$  ← very expensive  
    (BinTreeSampler(k), BinTreeSampler(n-1-k))
```

- ▶ Need an explicit recurrence.
- ▶ Very slow for large size objects.

Boltzmann samplers

Principle: sample objects $\gamma \in \mathcal{C}$ according to the probability distribution

$$\mathbb{P}(\gamma) = \frac{t^{|\gamma|}}{C(t)}$$

with t a parameter, by following a recursive decomposition of \mathcal{C} .

For binary trees :

$$\begin{aligned} \text{Tree with root } &= \text{Single node } + \text{Tree with root } \\ T(t) &= t + tT(t)^2 \end{aligned}$$

$$\Gamma T(t) = \left\{ \begin{array}{ll} \text{Single node } & \text{probability } \frac{t}{T(t)} \\ \text{Tree with root } & \text{probability } tT(t) \\ \Gamma T(t) & \Gamma T(t) \end{array} \right.$$

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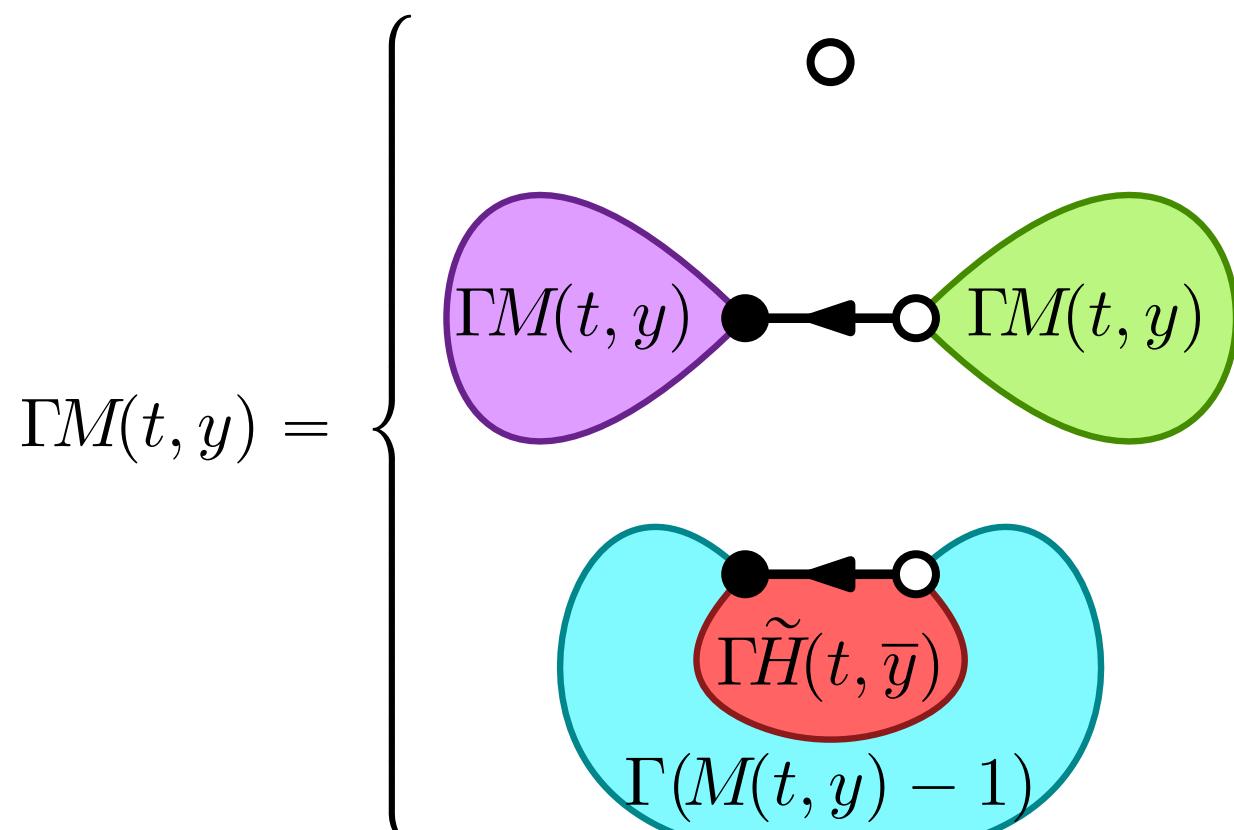
$$\begin{aligned}
 \text{Tree with root } \text{---} &= \text{Leaf } \text{---} + \text{Tree with root } \text{---} \\
 T(t) &= t + tT(t)^2
 \end{aligned}
 \quad \left| \quad \Gamma T(t) = \begin{cases} \text{Leaf } \text{---} & \text{probability } \frac{t}{T(t)} \\ \text{Tree with root } \text{---} & \text{probability } tT(t) \\ \text{Tree with root } \text{---} & \text{Tree with root } \text{---} \end{cases} \right.$$

- ▶ Unfixed size (adjust the value of t to aim a given size).
- ▶ Uniform distribution when conditioned by the size of the object.
- ▶ Can be built as long as we have a positive combinatorial equation.
- ▶ Very efficient for large objects.

Boltzmann sampler for 3-colored maps

- The catalytic variables becomes a second parameter: $\mathbb{P}(\mathfrak{m}) = \frac{t^{e(\mathfrak{m})} y^{b(\mathfrak{m})}}{M(t, y)}$.
- We use rejection to handle the $[y^{>0}]$ and $[x^{>0}]$.

$$\left\{ \begin{array}{l} M(t, y) = 1 + ty \textcolor{teal}{M}(t, y)^2 + ty[y^{>0}] (\textcolor{teal}{M}(t, y) - 1) \tilde{H}(t, \bar{y}), \\ H(t, x) = 1 + tx \textcolor{red}{H}(t, x)^2 + 2tx[x^{>0}] (\textcolor{red}{H}(t, x) - 1) \tilde{M}(t, \bar{x}). \end{array} \right.$$



$\Gamma H(t, x)$ similarly.

with probability $\frac{1}{M(t, y)}$

with probability $tyM(t, y)$

Otherwise.

Reject while the lengths of the borders of the maps are incompatible.

Boltzmann sampler for 3-colored maps

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$$\begin{cases} M(t, y) = 1 + ty \textcolor{teal}{M}(t, y)^2 + ty[y^{\geq 0}] (\textcolor{teal}{M}(t, y) - 1) \tilde{H}(t, \bar{y}), \\ H(t, x) = 1 + tx \textcolor{red}{H}(t, x)^2 + 2tx[x^{\geq 0}] (\textcolor{red}{H}(t, x) - 1) \tilde{M}(t, \bar{x}). \end{cases}$$

$$\Gamma \tilde{H}(t, \bar{y}) = \textcolor{red}{\Gamma H(t, x)} + \textcolor{pink}{\Gamma C(\bar{x}, \bar{y})}$$

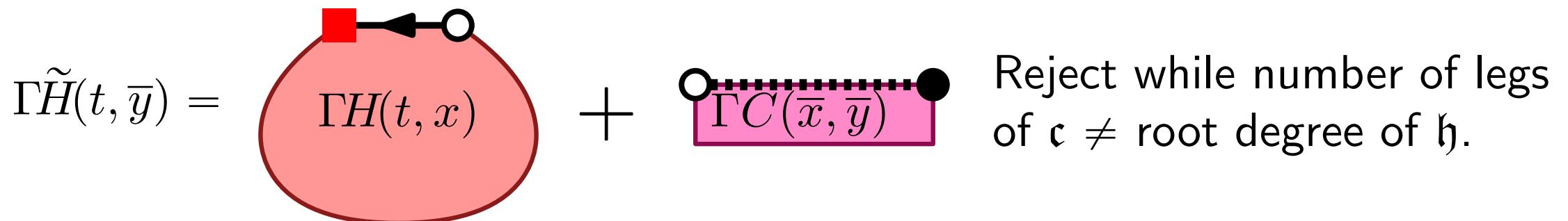
Reject while number of legs of $\mathfrak{c} \neq$ root degree of \mathfrak{h} .

C counts caterpillars of even length with $\bar{x} \rightarrow$ legs and $\bar{y} \rightarrow$ half length.

Boltzmann sampler for 3-colored maps

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C counts caterpillars of even length with $\bar{x} \rightarrow$ legs and $\bar{y} \rightarrow$ half length.

$$\Gamma C(\bar{x}, \bar{y}) = \begin{cases} \bullet & \text{with probability } \sqrt{\bar{y}} / A(\bar{x}, \bar{y}), \\ \bullet \dots \textcolor{pink}{\Gamma C(\bar{x}, \bar{y})} & \text{with probability } \sqrt{\bar{y}}, \\ \textcolor{pink}{\Gamma C(\bar{x}, \bar{y})} & \text{with probability } \bar{x}. \end{cases}$$

Reject if the length is odd.

Boltzmann sampler for 3-colored maps

Pretty image of a big random 3-colored map
that I'll have as soon as I find time to
implement this sampler.

Perspectives

- ▶ Adapt the gasket decomposition for others families of colored maps.
 - ↪ Colorful quadrangulations [Budd '25].
 - ↪ 4-colored triangulations.
- ▶ Others applications of the decomposition?
- ▶ Implement the samplers.

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Thank you !