

Slit-slide-sew bijections for bipartite planar maps with prescribed degrees

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LaBRI





Introduction

Bipartite planar maps

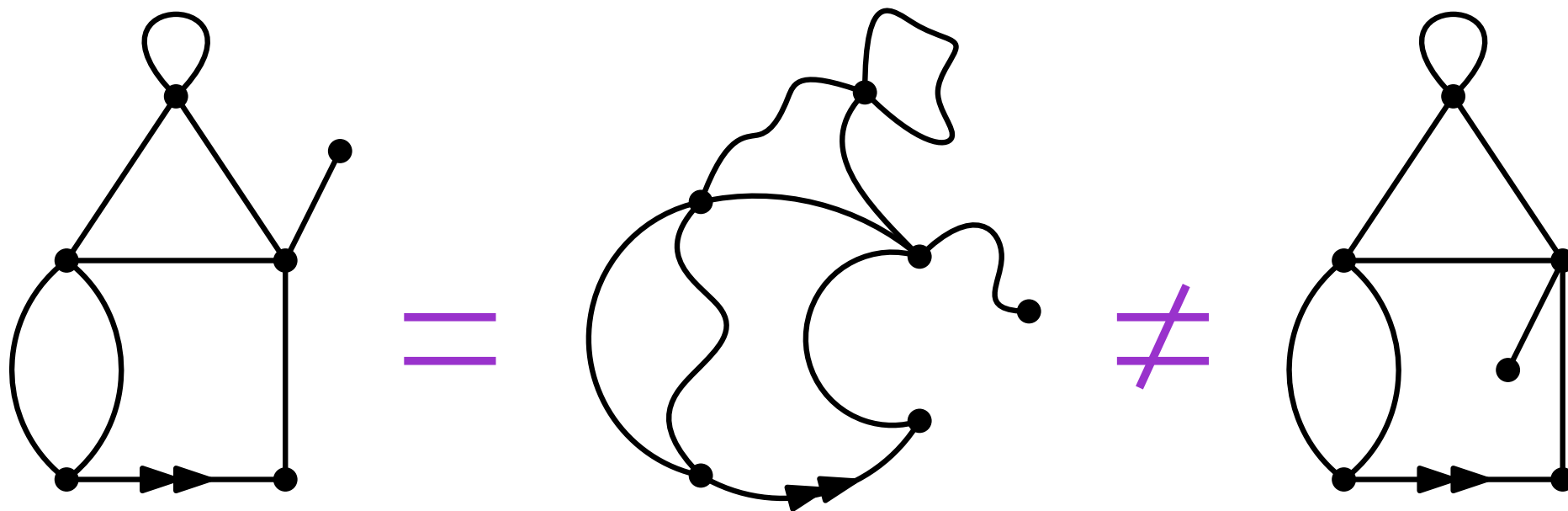
I- Enumeration of maps & Integrable Hierarchies

II- Bijections

Planar maps

A **planar map** is the embedding of a connected graph onto the sphere, up to orientation preserving homeomorphism.

Multi edges and loops are allowed.



Planar map = planar graph + cyclic ordering of the edges around each vertex.

All maps are **rooted**, i.e. an oriented edge is marked.

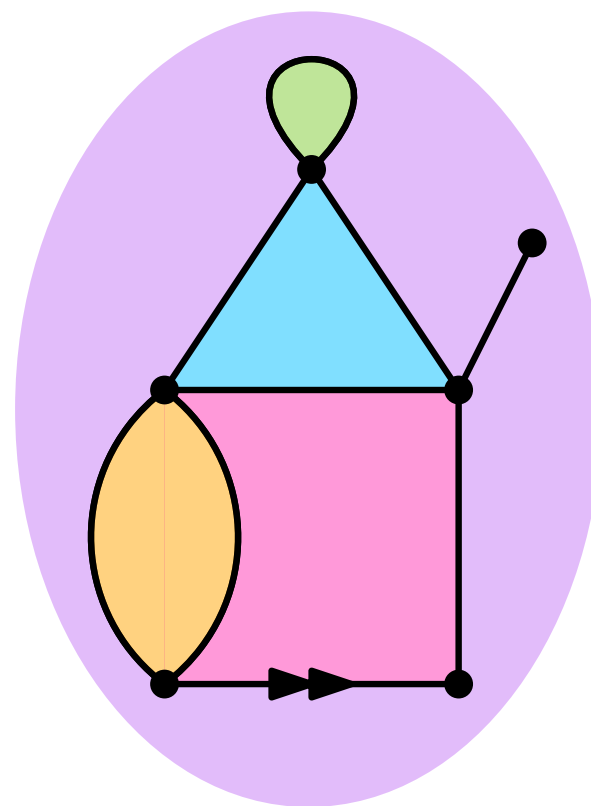
Map of genus $g > 0$ = graph embedded on a g -torus.

Planar maps

Planar map = planar graph + cyclic ordering of the edges around each vertex.

- **Vertices** and **edges** are inherited from the graph. (Counted by v and n)
- **Faces** are the connected components of the sphere minus the map. (Counted by f .)

Here: 6 vertices, 9 edges
et 5 faces.

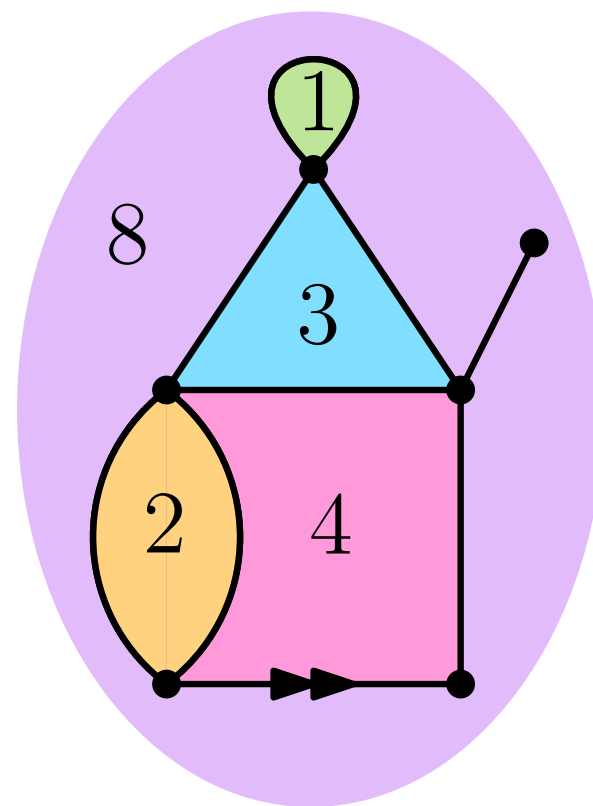
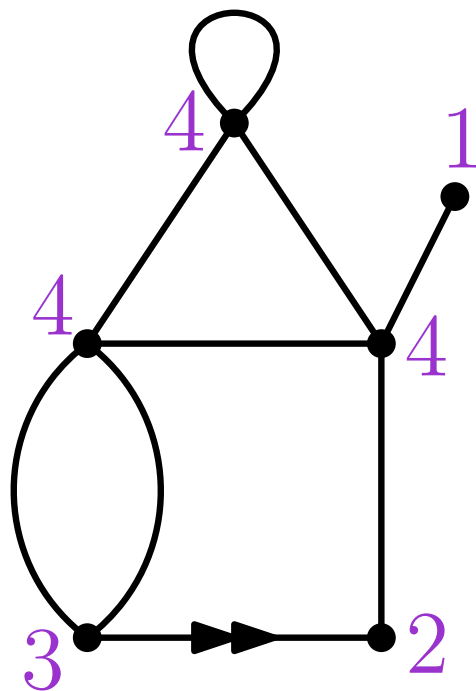


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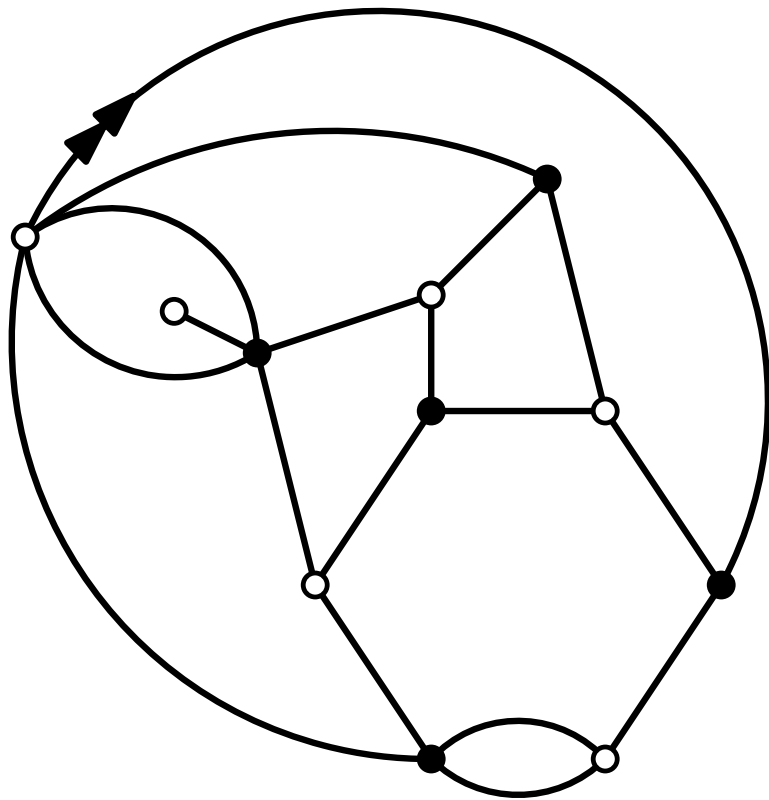
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Degree (of a vertex or face) = number of incident half edges.

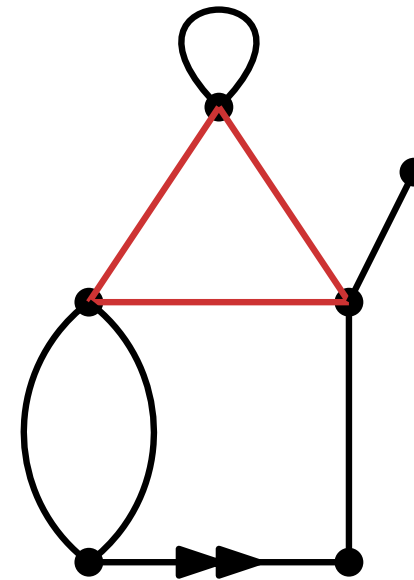


Bipartite maps

A planar map is **bipartite** if one can properly color its vertices in black and white $\iff$ all the faces have even degree.



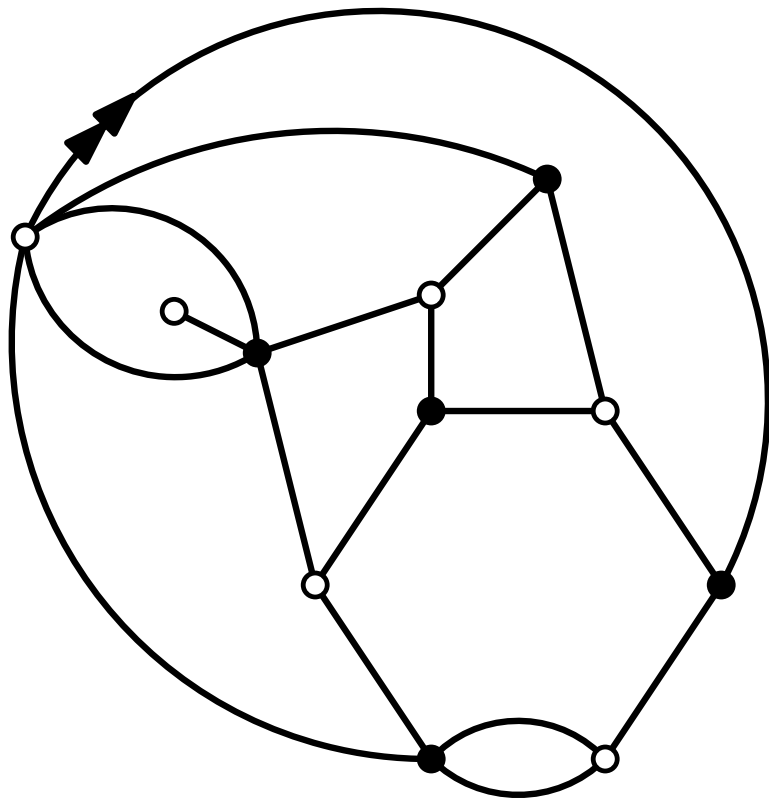
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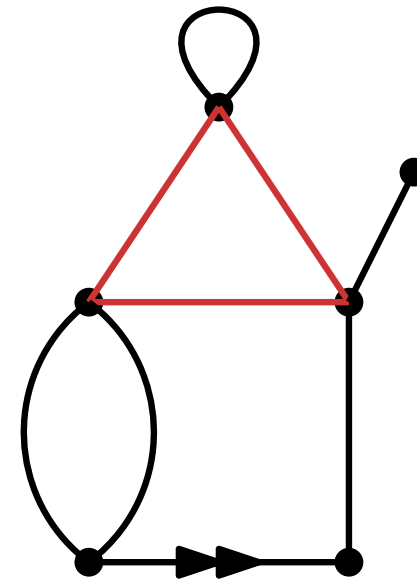
Not bipartite

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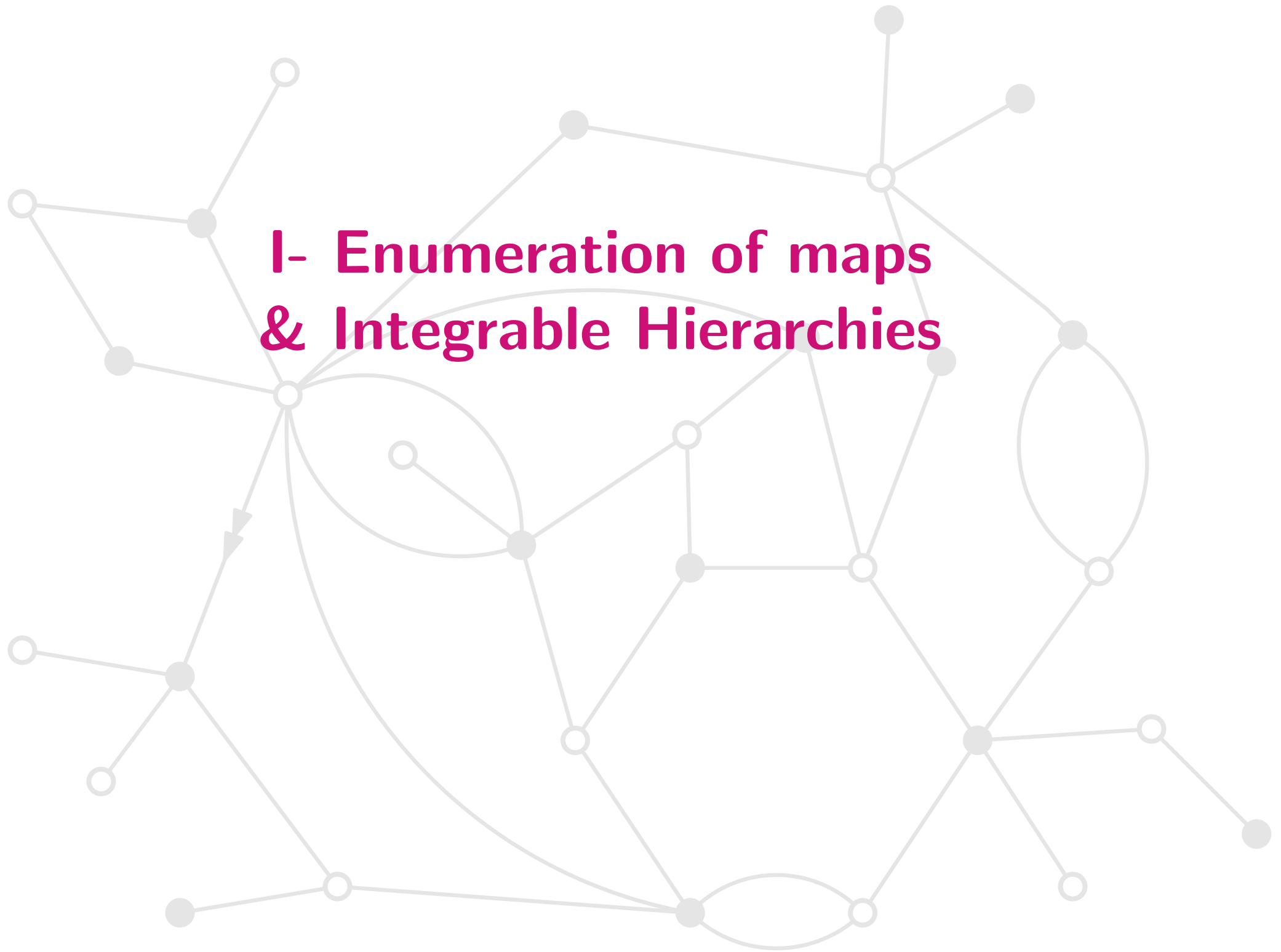
Bipartite



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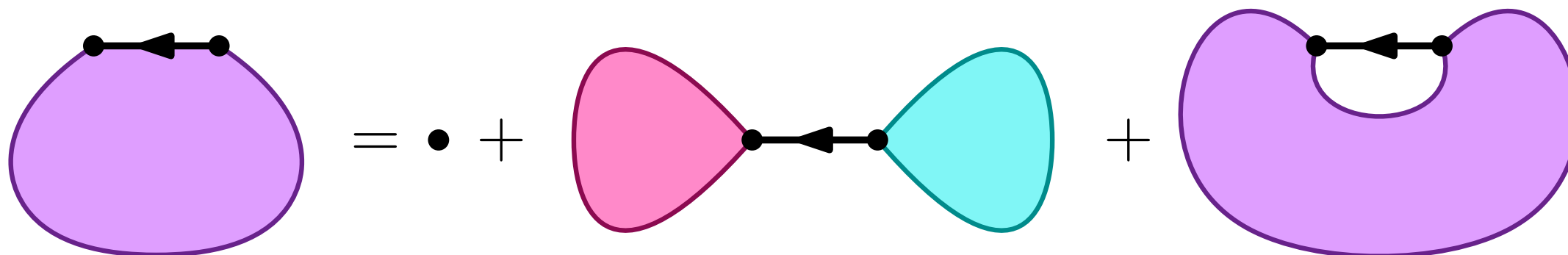
Canonical coloring: root edge from white to black. $\circ \rightarrow \bullet$

I- Enumeration of maps & Integrable Hierarchies



Recursive approach

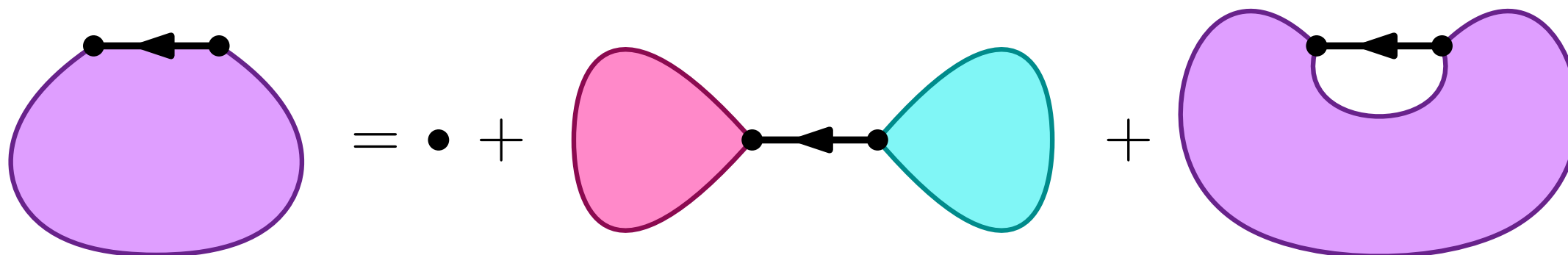
1) Find a decomposition of the map and write the associated equation on the generating series.



$$M(t, y) = 1 + ty^2 M(t, y)^2 + ty \frac{yM(t, y) - M(t, 1)}{y - 1}$$

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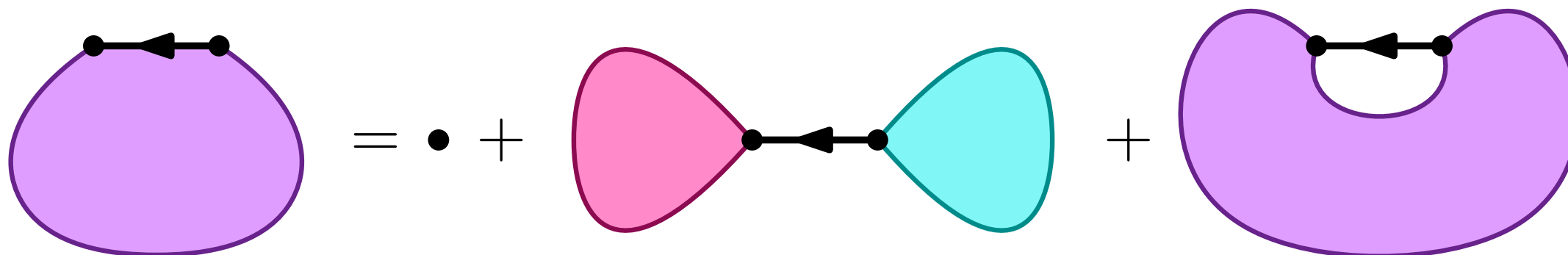
2) Solve the equation or manipulate it to get coefficients or asymptotics.

Theorem. [Tutte '63] The number of planar maps with n edges is

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Number of planar trees with n edges.

- ▶ Robust method.
- ▶ Bijections revealing hidden structure of maps. [Cori-Vauquelin-Schaeffer '98]

Integrable hierarchies



$$\partial_x(\partial_t u + u\partial_x u + \varepsilon^2 \partial_{xxt} u) + \lambda \partial_{yy} u = 0$$

2-Toda hierarchy: Extension of KP with two infinite sets of variables.

KP hierarchy:

Infinite set of PDEs with an infinite set of variables $(p_1, p_2, \dots)$.

Obtained by studying the symmetries of the Kadomtsev-Petviashvili equation.

$$F_{3,1} = F_{2,2} + \frac{1}{2}F_{1,1}^2 + \frac{1}{12}F_{1,1,1,1}$$

$$F_{4,1} = F_{3,2} + F_{1,1}F_{2,1} + \frac{1}{6}F_{1,1,1,2} \dots$$

Integrable hierarchies



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2-Toda hierarchy: Extension of KP with two infinite sets of variables.

Some families of solutions of the integrable hierarchies are known.

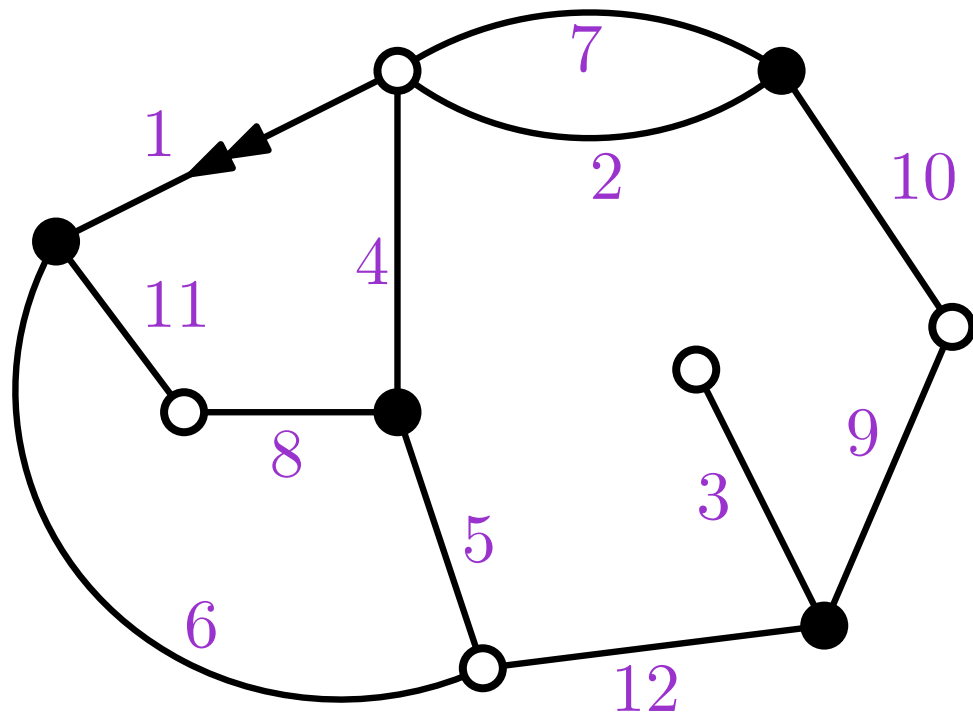
$\Rightarrow$ If a series is of that form then it satisfies the equations!

Algebraic combinatorics and KP/Toda

Maps = permutations factorisation!

Representation of **bipartite maps** (edge labeled, not planar nor connected) :
a cycle of a permutation = the edges around a vertex/face.

Degree = cycle length.



$$\sigma_{\circ} = (1, 4, 2, 7)(3)(5, 6, 12)(8, 11)(9, 10)$$

$$\sigma_{\bullet} = (1, 6, 11)(2, 10, 7)(3, 12, 9)(4, 8, 5)$$

$$\varphi = (1, 12, 10)(2, 9, 3, 5)(4, 11)(6, 8)(7)$$

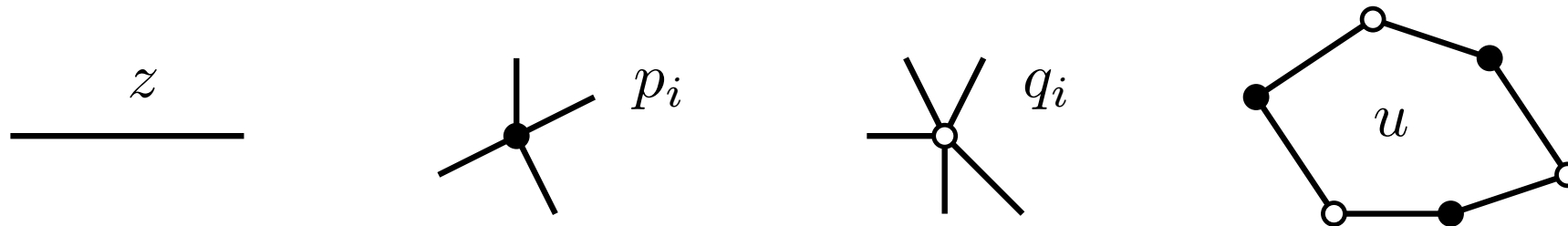
$$\sigma_{\bullet}\sigma_{\circ} = \varphi$$

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Using permutation encoding and representation theory, one can show that the series B of bipartite maps with weights



is an **hypergeometric tau** function, i.e. solution to the KP/Toda hierarchy.

↪ the PDEs translate as recurrence equations for given specialisations.

Combinatorial formulas from KP/Toda

Theorem. [Goulden-Jackson '08] Let $T_g(n)$ be the number of triangulations of genus g with $2n$ faces.

$$(n+1)T_g(n) = 4n(3n-2)(3n-4)T_{g-1}(n-2) + 4(3n-1)T_g(n-1) \\ + 4 \sum_{i+j=n-2} \sum_{g_1+g_2=g} (3i+2)(3j+2)T_{g_1}(i)T_{g_2}(j)$$

Theorem. [Louf '21] Let $B_g(\mathbf{f})$ be the number of bipartite maps of genus g with degree sequence $\mathbf{f}$.

$$\binom{n+1}{2} B_g(\mathbf{f}) = \sum_{\substack{\mathbf{s}+\mathbf{t}=\mathbf{f} \\ g_1+g_2+g^*=g}} (1+n_1) \binom{v_2}{2g^*+2} B_{g_1}(\mathbf{s}) B_{g_2}(\mathbf{t}) + \sum_{g^* \geq 0} \binom{v+2g^*}{2g^*+2} B_{g-g^*}(\mathbf{f})$$

- ▶ Very fast counting.
- ▶ Lots of parameters.
- ▶ New bijections explaining the structure of maps.

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- ▶ Single vertex case [Chapuy-Féray-Fusy '13]
- ▶ Planar case [Louf '19]

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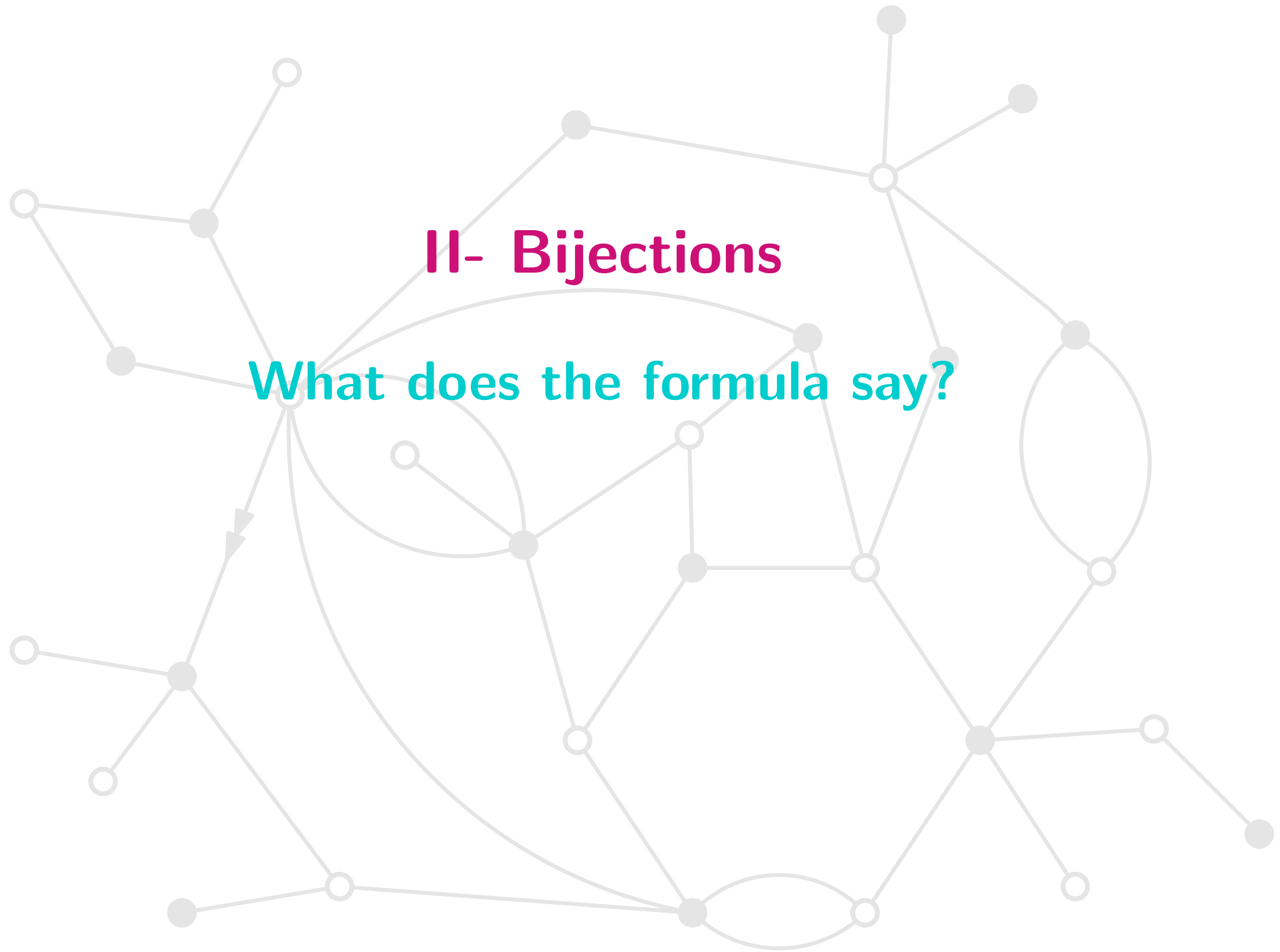
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- ▶ Planar case [S. '24]

A complex graph structure with nodes and edges, overlaid with text. The graph consists of numerous nodes, some of which are white with black outlines, and others are solid gray. These nodes are interconnected by a network of thin gray lines representing edges. Some edges form loops or cycles. Overlaid on this graph is the text "II- Bijections" in a bold, magenta font, and below it, "What does the formula say?" in a bold, teal font. A faint, light gray oval highlights a specific region of the graph, centered around the text. A small, dark gray arrow points towards one of the nodes within this highlighted area.

II- Bijections

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Bipartite maps with prescribed degrees

Let $\mathbf{d} = (d_1, d_2, \dots)$ be a sequence of integers. A bipartite map M has **degree sequence** $\mathbf{d}$ if it has exactly d_i faces with degree $2i$ for all i . Denote by $\mathcal{B}(\mathbf{d})$ their set and $B(\mathbf{d})$ their number.

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Euler formula: $v + f = n + 2$ for all planar maps.

A map $M \in \mathcal{B}(\mathbf{d})$ has :

- $f(\mathbf{d}) = \sum_{i \geq 1} d_i$ faces
- $n(\mathbf{d}) = \sum_{i \geq 1} i d_i$ edges
- $v(\mathbf{d}) = n(\mathbf{d}) + 2 - f(\mathbf{d})$ vertices

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Theorem. [Louf '21]

$$\left(\binom{n(\mathbf{d}) + 1}{2} - \binom{v(\mathbf{d})}{2} \right) B(\mathbf{d}) = \sum_{\mathbf{s} + \mathbf{t} = \mathbf{d}} (1 + n(\mathbf{s})) \binom{v(\mathbf{t})}{2} B(\mathbf{s}) B(\mathbf{t}).$$

$\swarrow$
 $s_i + t_i = d_i \forall i$

Louf's formula: planar case

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Euler formula: $n + 1 = v + f - 1$ for all planar maps.

$$\left((f(\mathbf{d}) - 1)v(\mathbf{d}) + \binom{f(\mathbf{d}) - 1}{2} \right) B(\mathbf{d}) = \sum_{\mathbf{s} + \mathbf{t} = \mathbf{d}} (v(\mathbf{s}) + f(\mathbf{s}) - 1) \binom{v(\mathbf{t})}{2} B(\mathbf{s}) B(\mathbf{t}).$$

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Theorem. [S. '24] The number of bipartite planar maps with profile $\mathbf{d}$ satisfy :

$$2v(\mathbf{d})(f(\mathbf{d}) - 1)B(\mathbf{d}) = \sum_{\mathbf{s} + \mathbf{t} = \mathbf{d}} v(\mathbf{s})v(\mathbf{t})(v(\mathbf{t}) - 1)B(\mathbf{s})B(\mathbf{t})$$

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this talk's
bijection

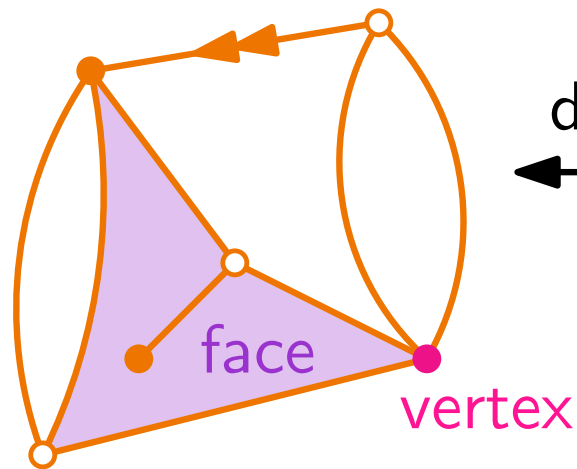
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Simplest proof: eulerian trees

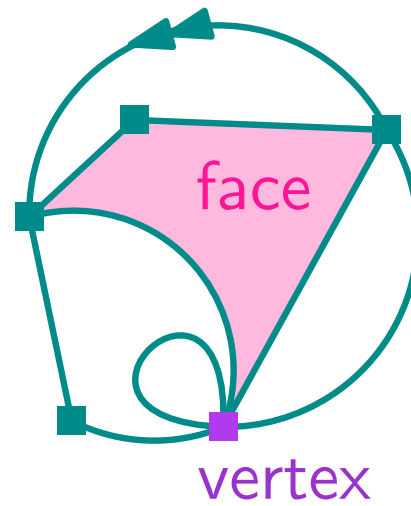
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Bipartite map



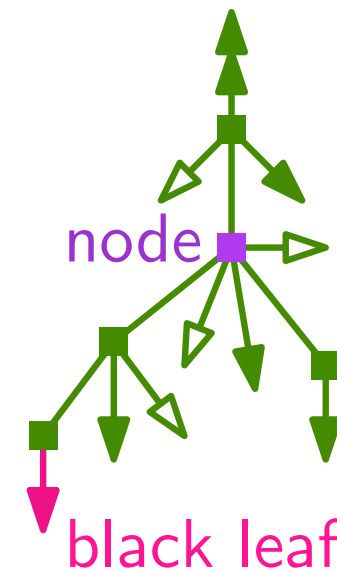
duality
↔

Eulerian map



Schaeffer
↔

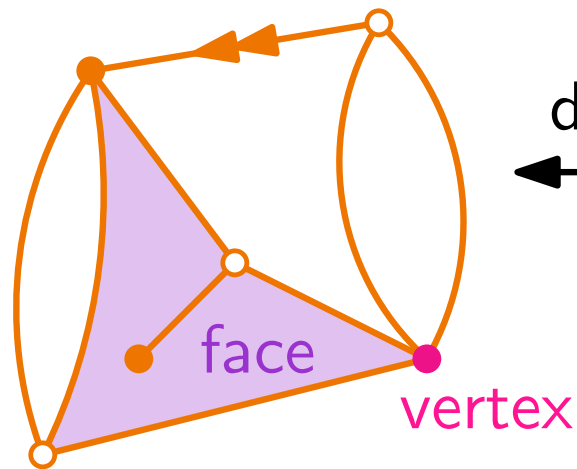
Eulerian Tree



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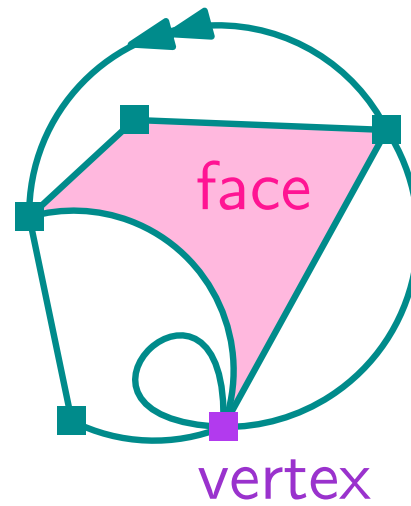
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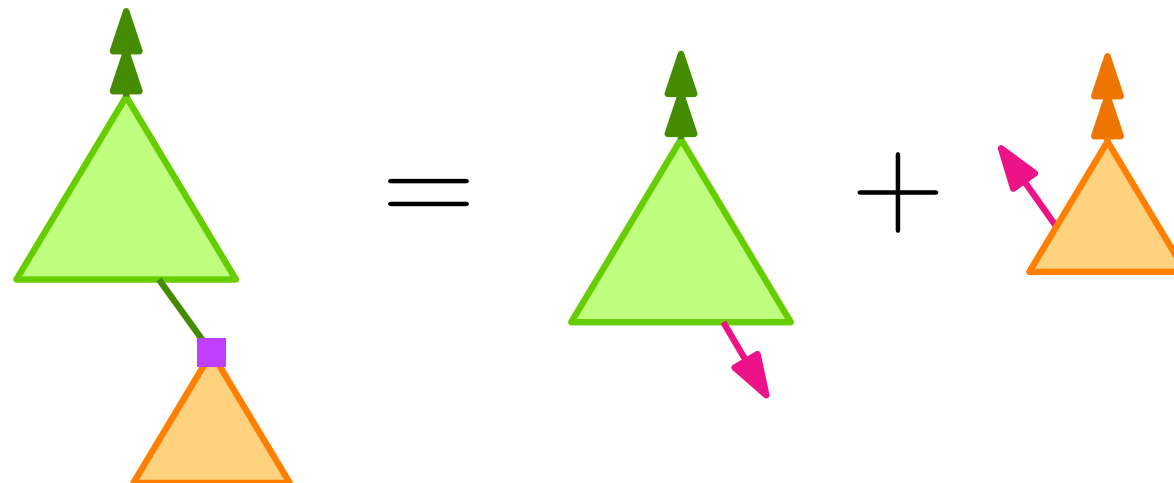
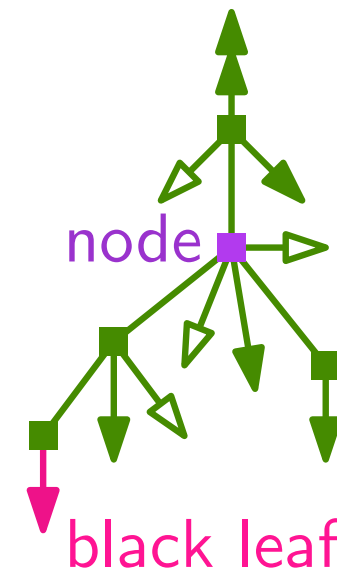
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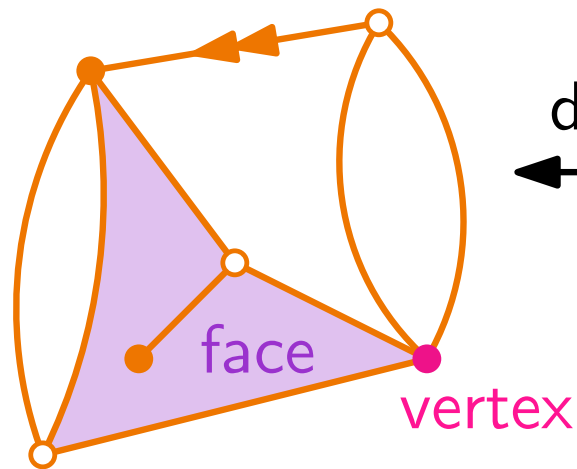
Eulerian Tree



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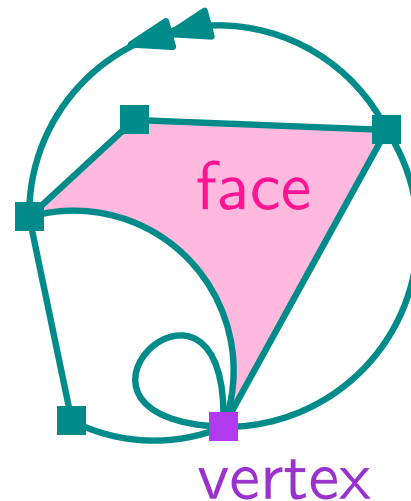
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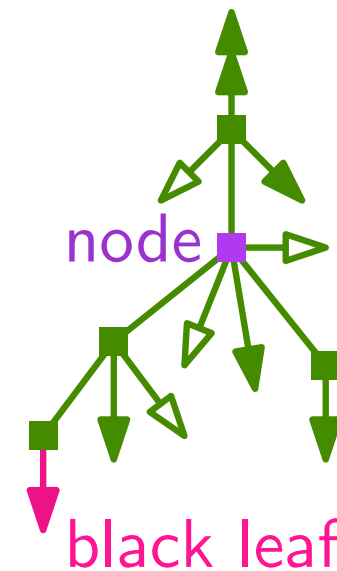
duality
↔

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Eulerian Tree



$$\begin{array}{ccccc}
 (M_1, v_1) + (M_2, v_2) & \xleftrightarrow{\text{duality}} & (M_1^*, v_1^*) + (M_2^*, v_2^*) & \xleftrightarrow{\text{Schaeffer}} & (T_1^*, l_1) + (T_2^*, l_2) \\
 \tilde{\varphi} \downarrow \uparrow \tilde{\psi} & & & & \varphi \downarrow \uparrow \psi \\
 (M, e) & \xleftrightarrow{\text{duality}} & (M^*, e^*) & \xleftrightarrow{\text{Schaeffer}} & (T^*, e^*)
 \end{array}$$

Not really enlightening ...

What happens here?



The 2 faces case

The two faces case of the original formula [Bouttier-Guitter-Miermont, '22].

$$4(\cancel{f(\mathbf{d})} - 1)B(\mathbf{d}) = \sum_{\mathbf{s} + \mathbf{t} = \mathbf{d}} v(\mathbf{s})v(\mathbf{t})B(\mathbf{s})B(\mathbf{t})$$

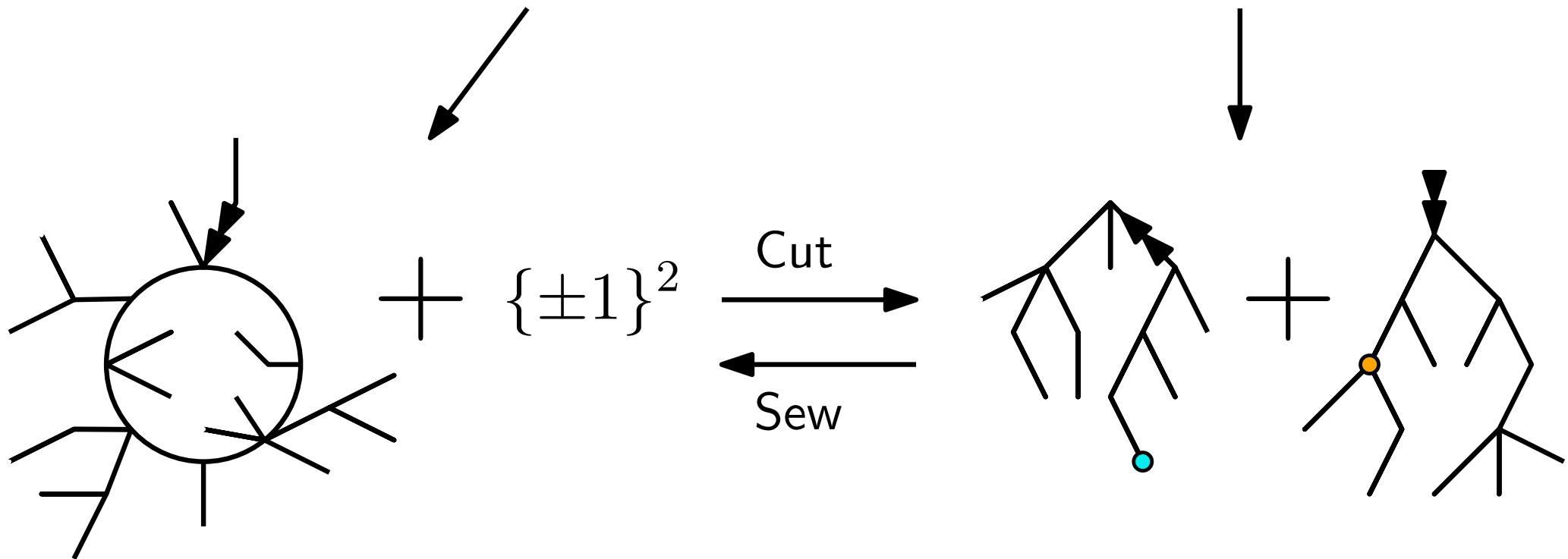
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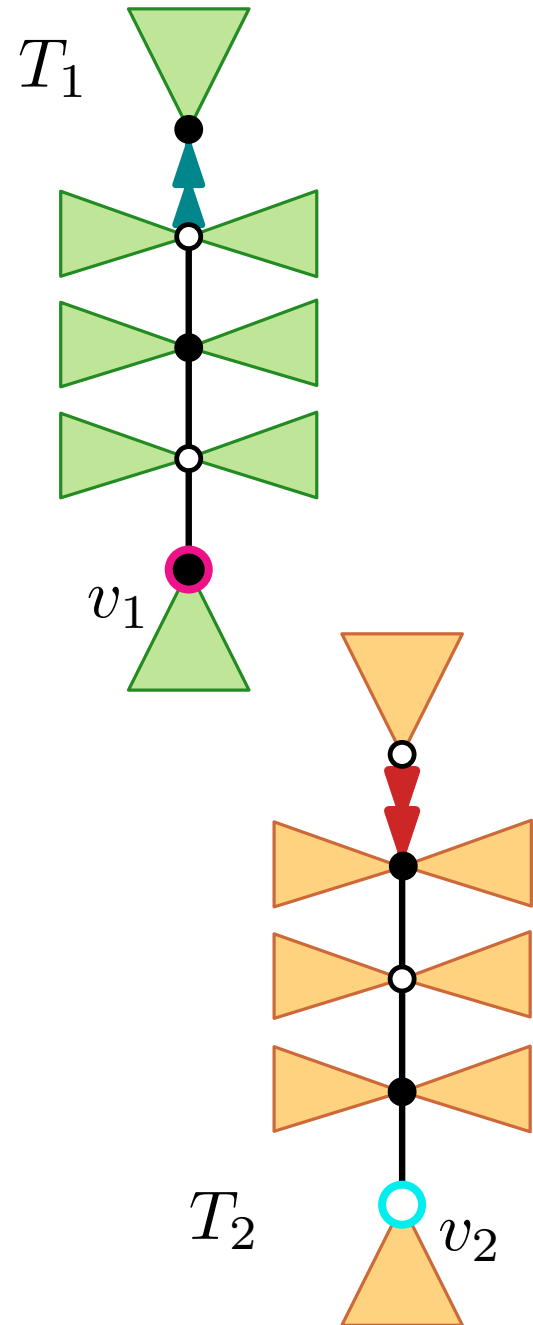
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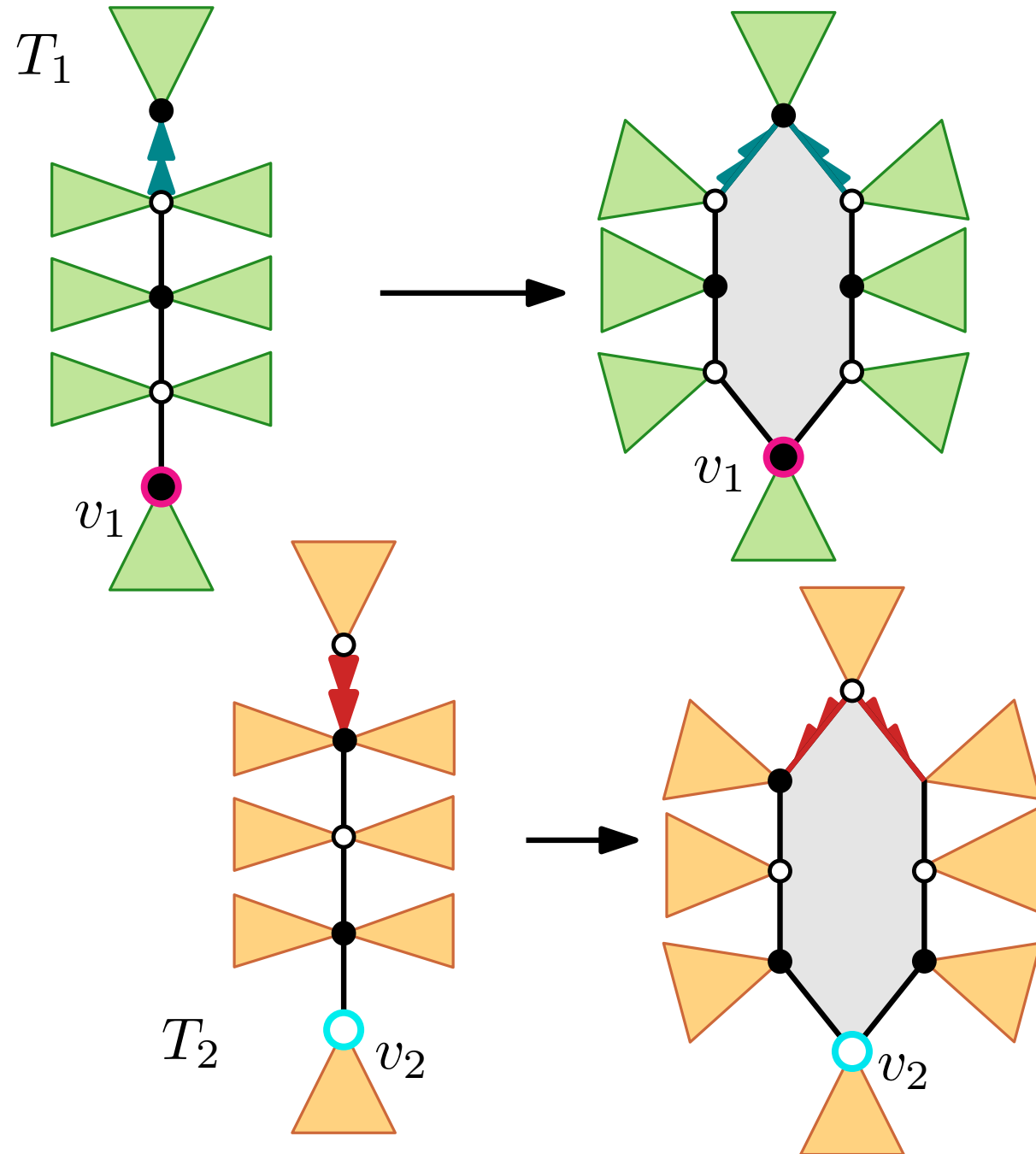
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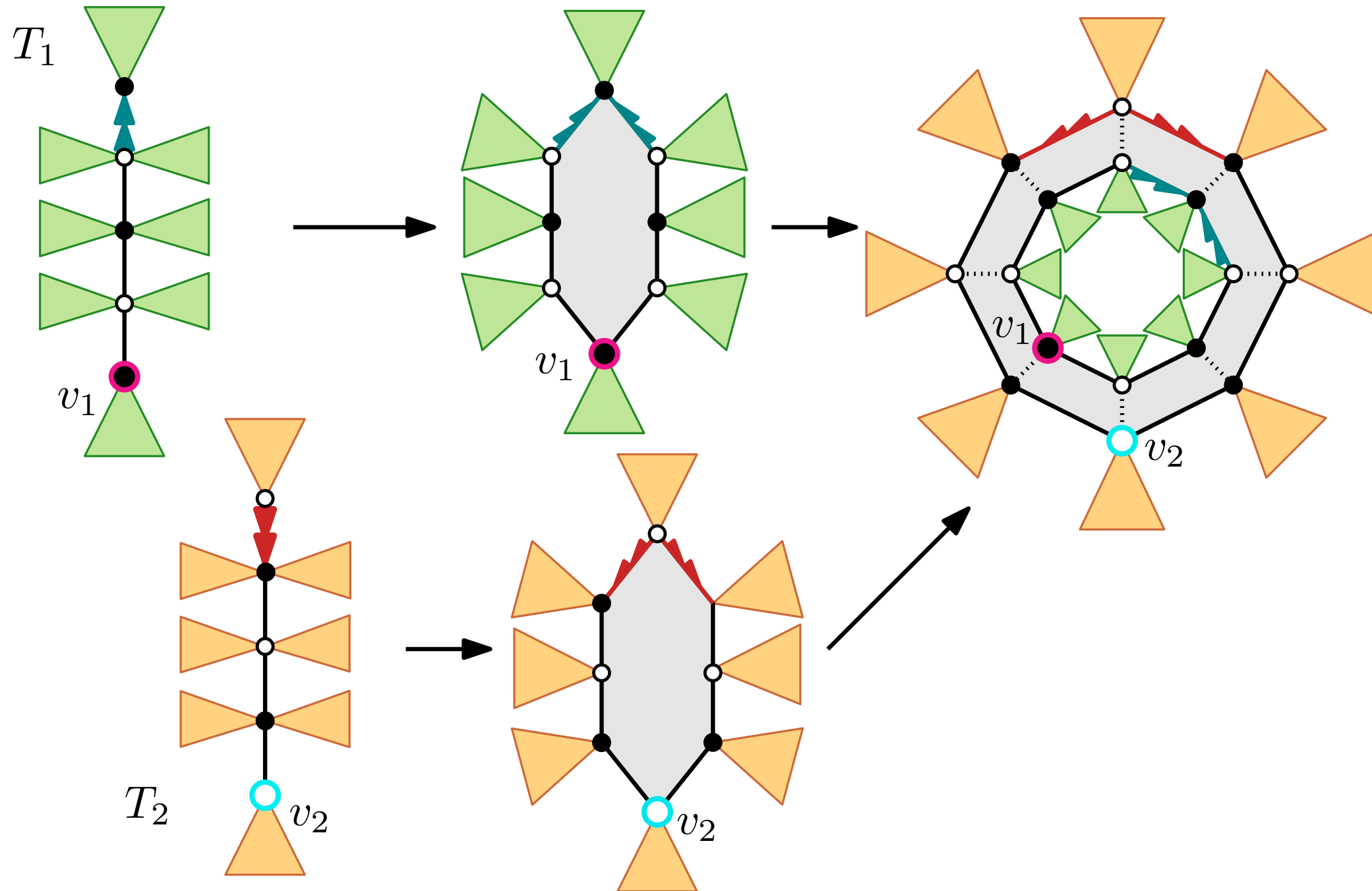
The 2 faces case: Slit and Sew



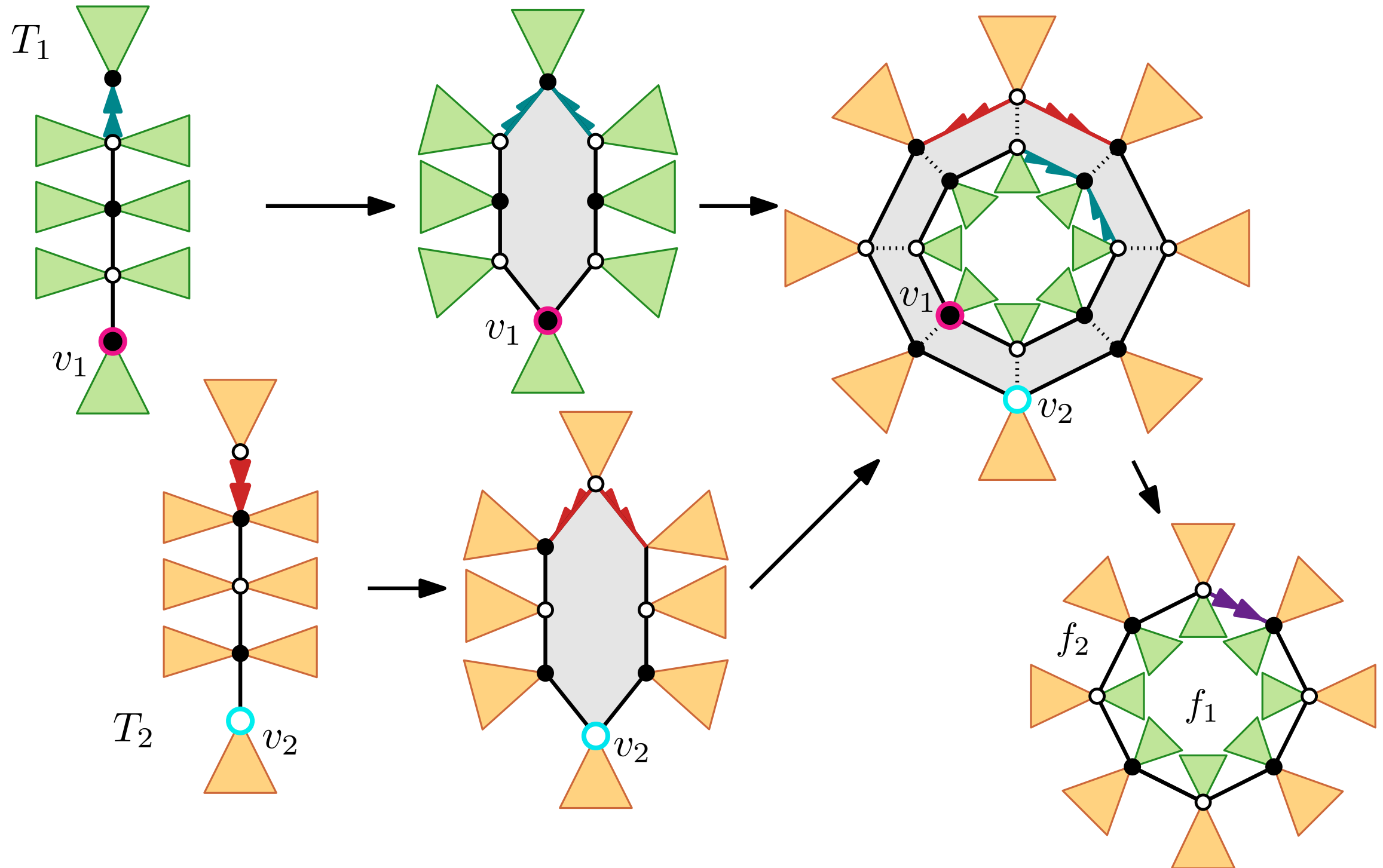
The 2 faces case: Slit and Sew



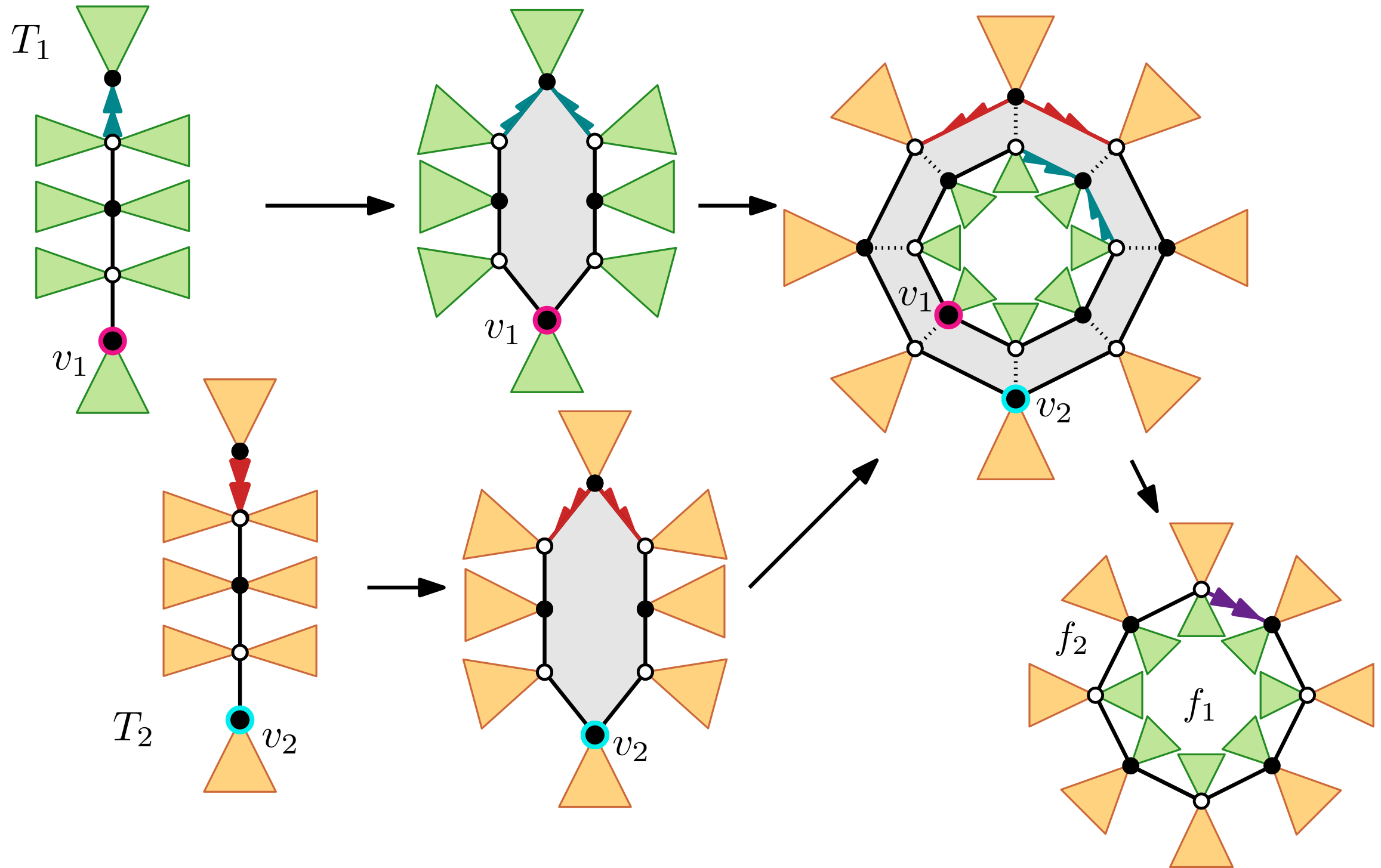
The 2 faces case: Slit and Sew



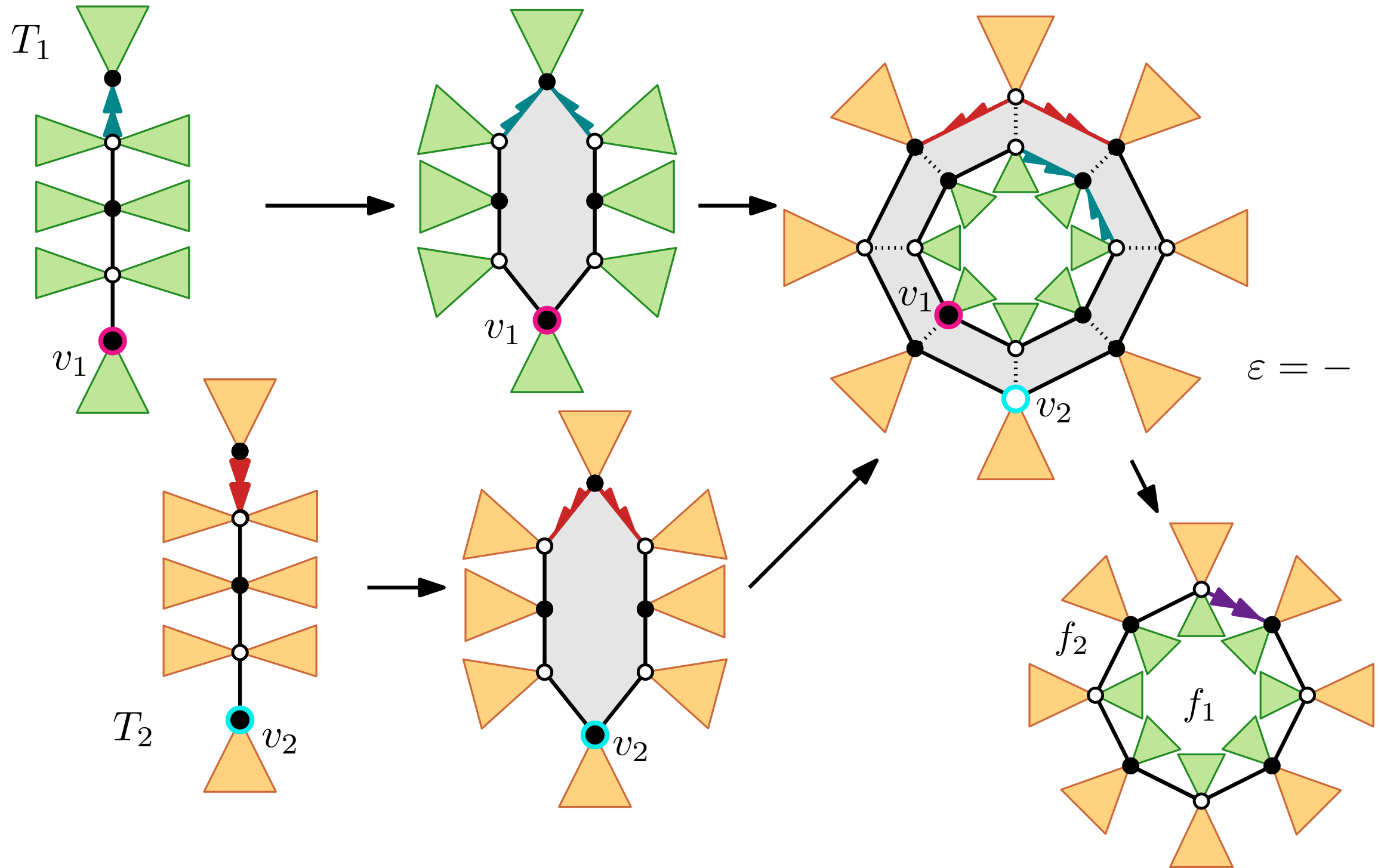
The 2 faces case: Slit and Sew



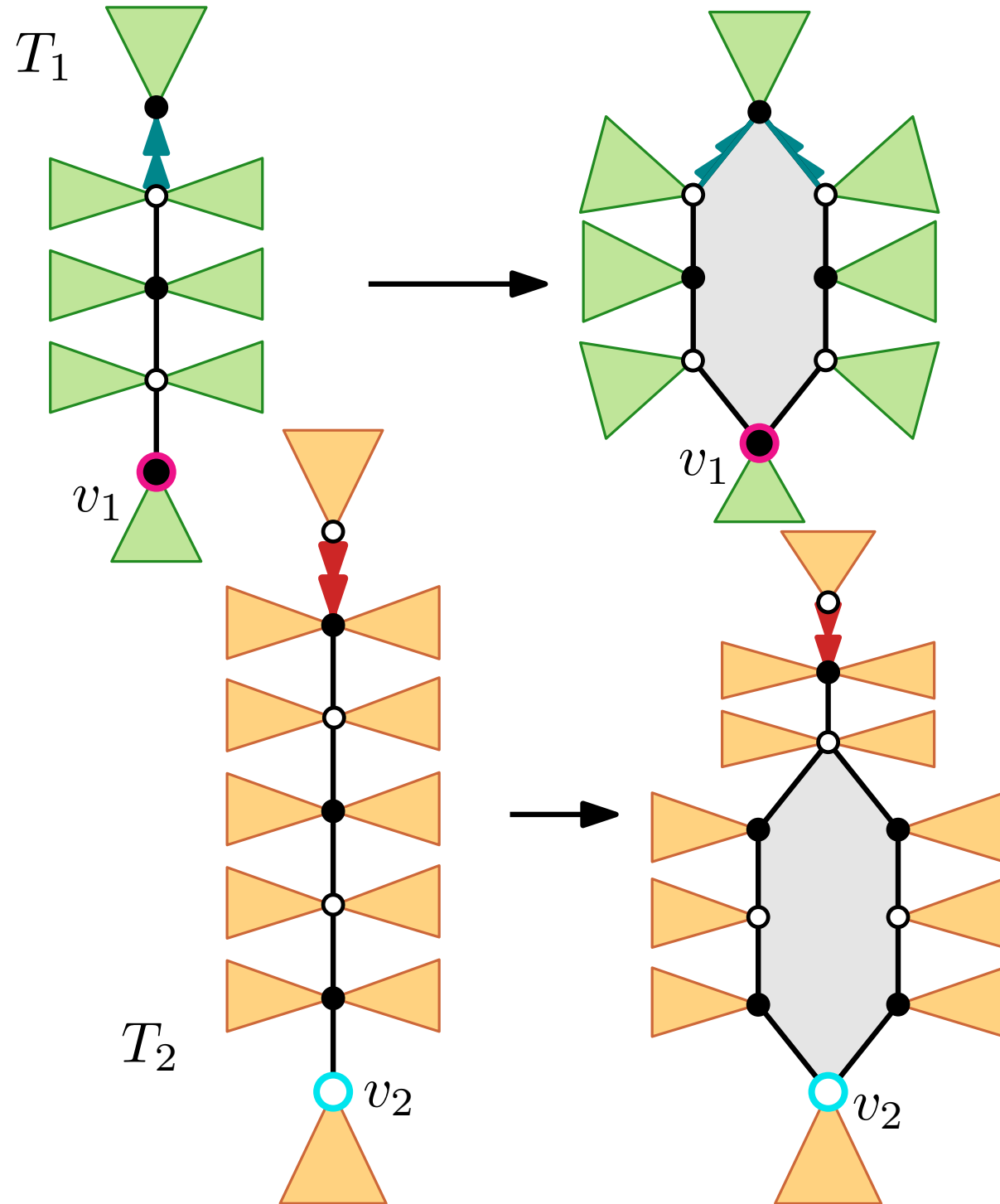
The 2 faces case: Slit and Sew



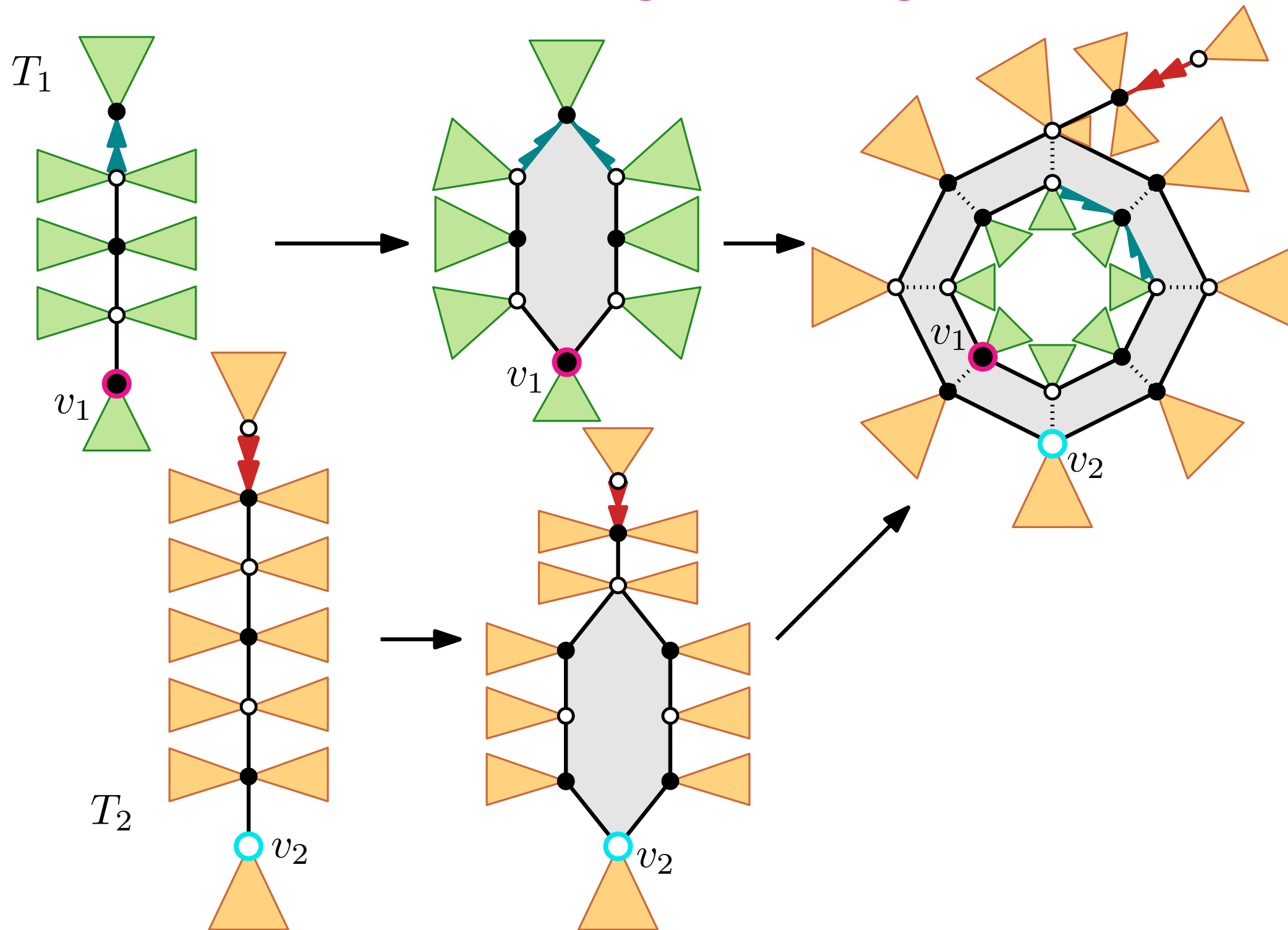
The 2 faces case: Slit and Sew



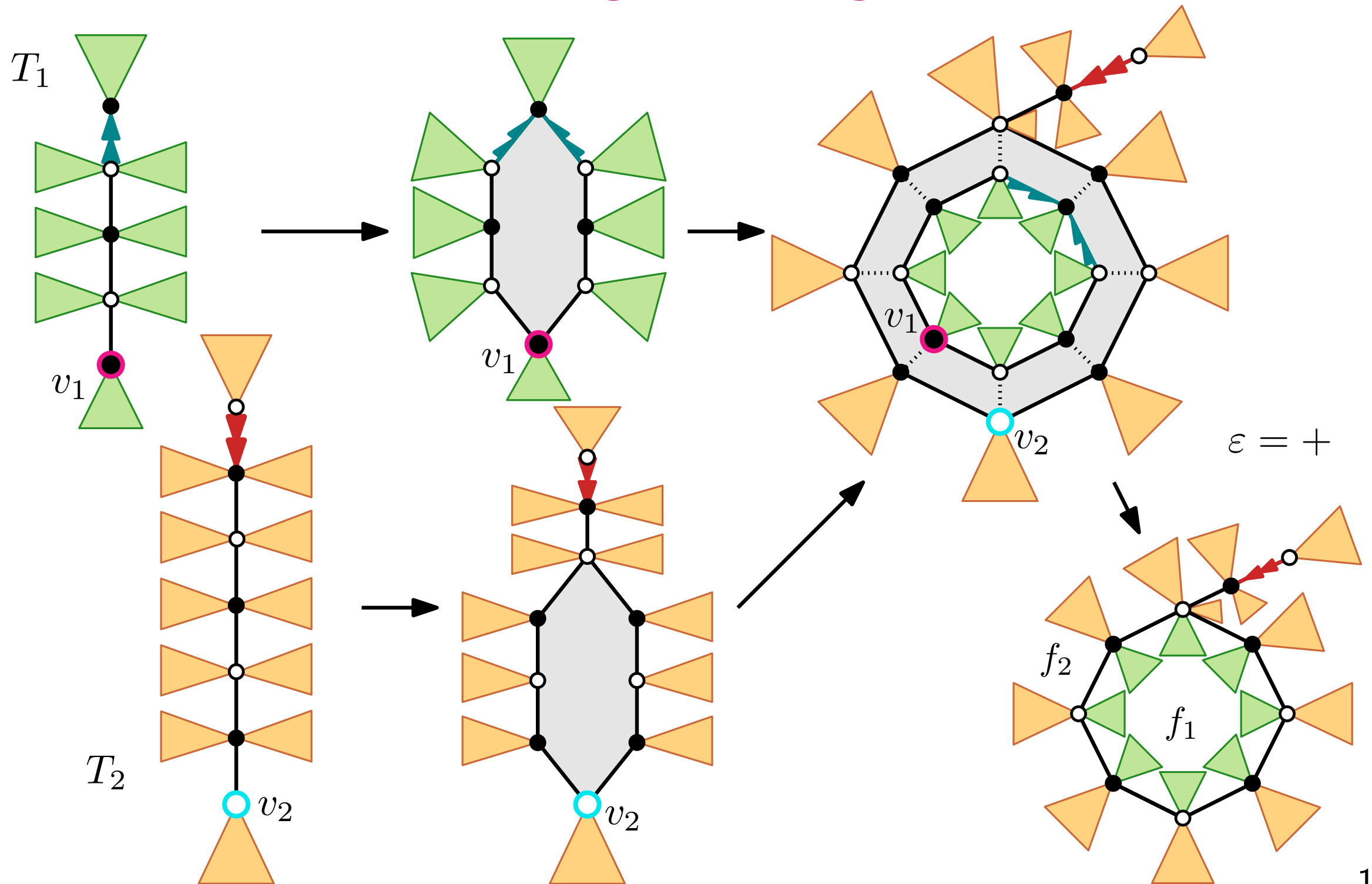
The 2 faces case: Slit and Sew



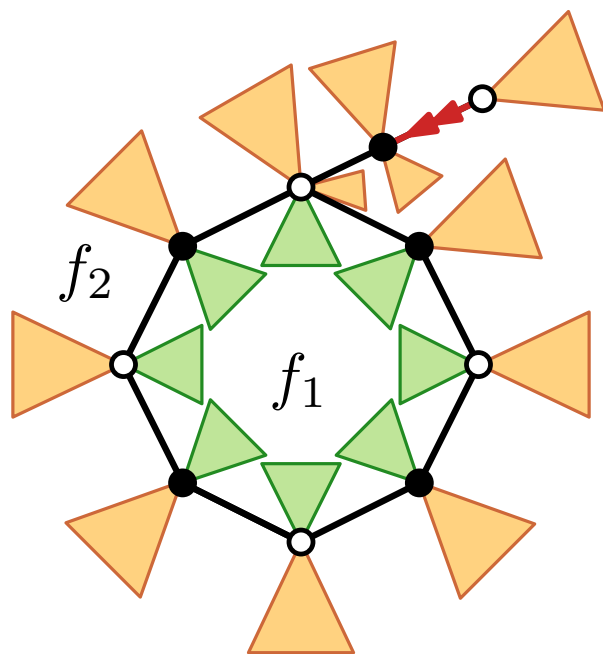
The 2 faces case: Slit and Sew



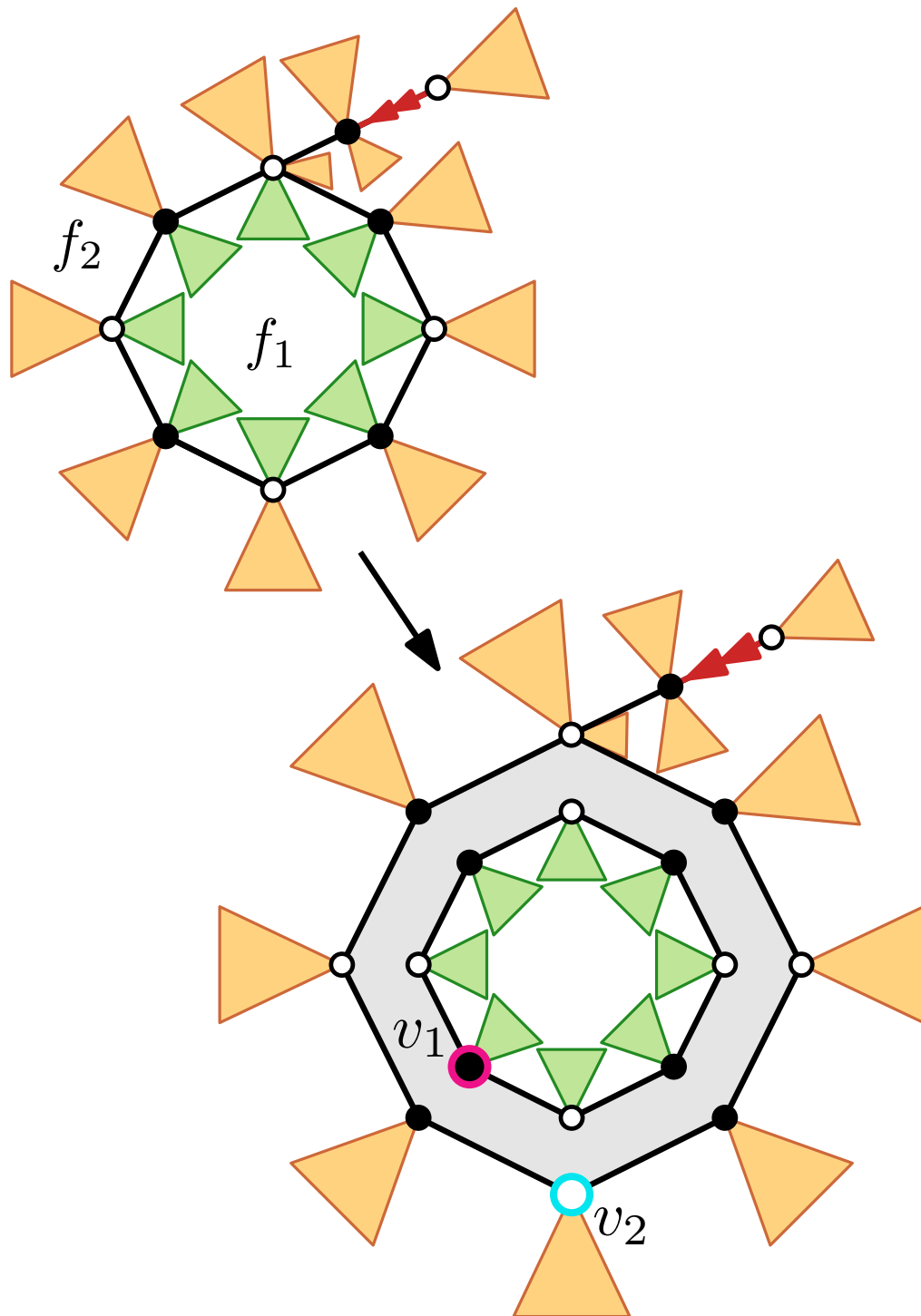
The 2 faces case: Slit and Sew



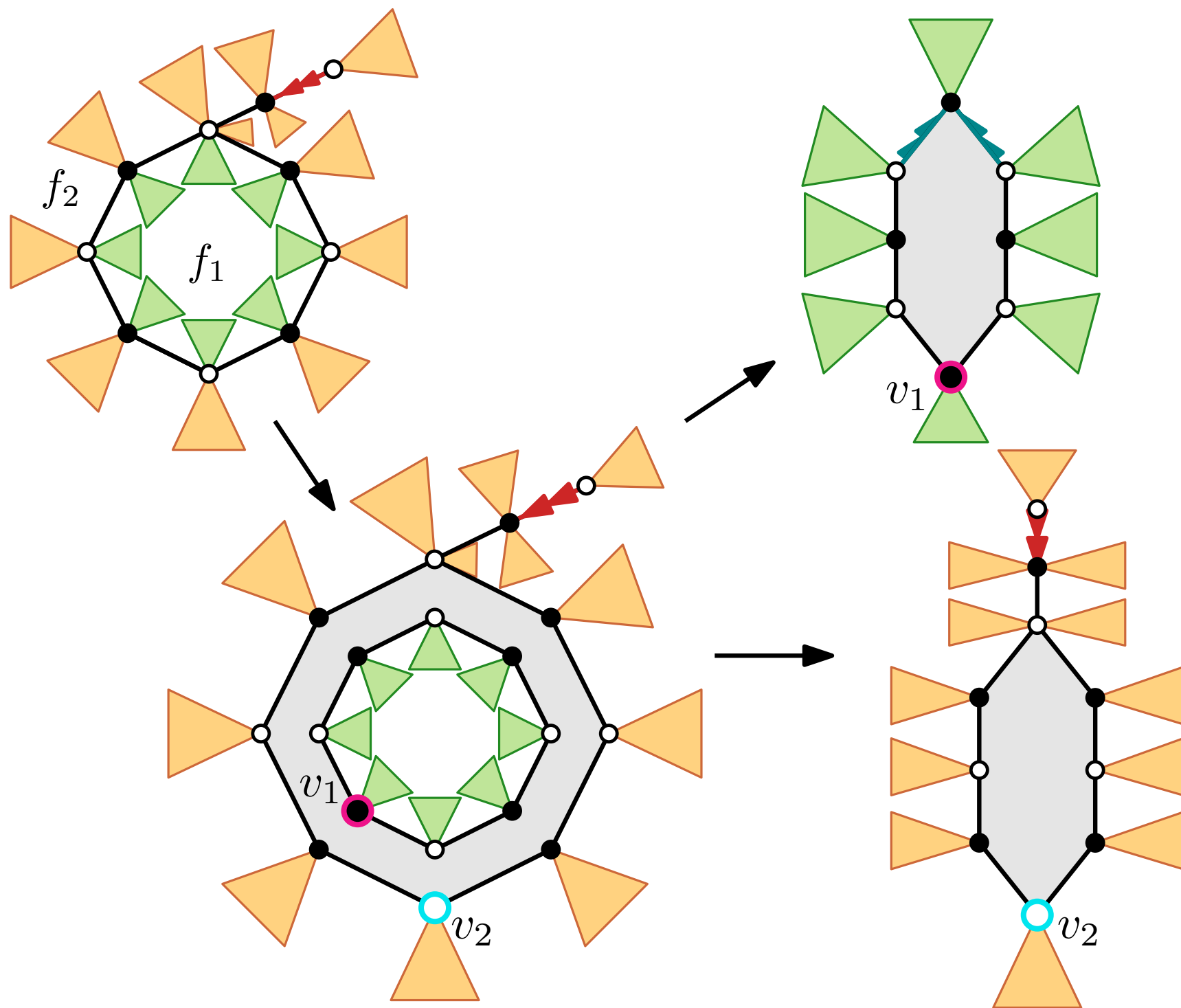
The 2 faces case: Cut and Close



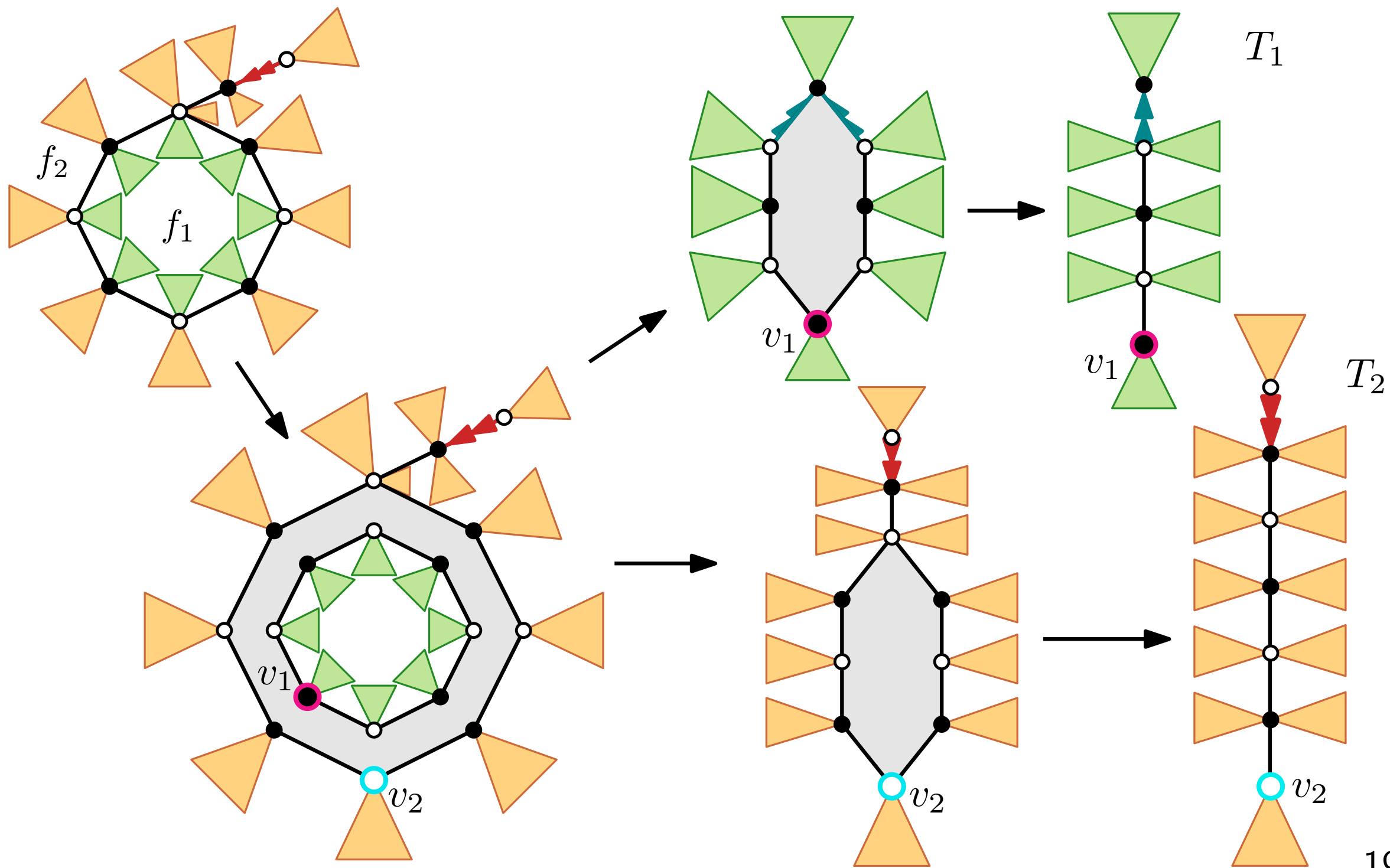
The 2 faces case: Cut and Close



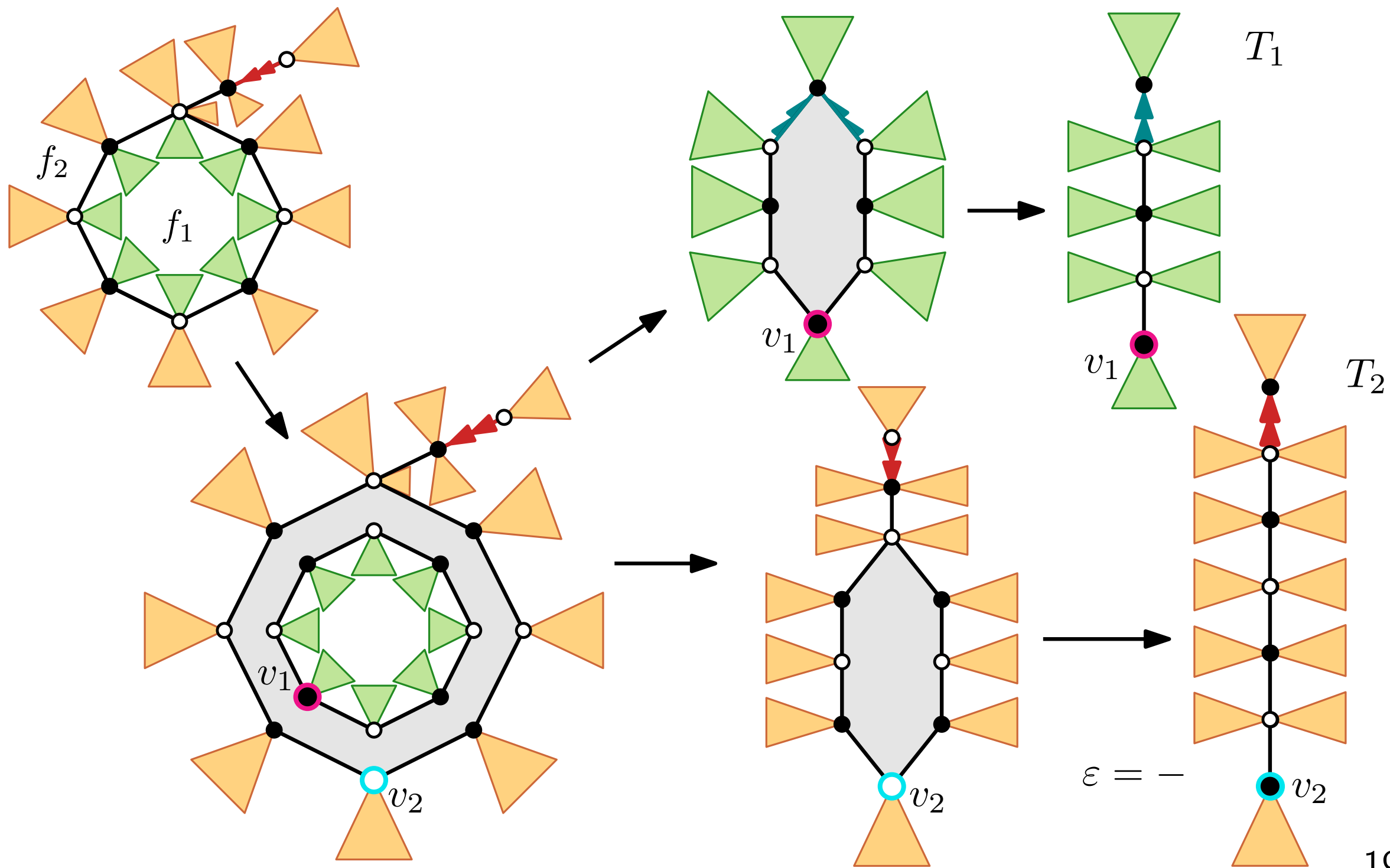
The 2 faces case: Cut and Close

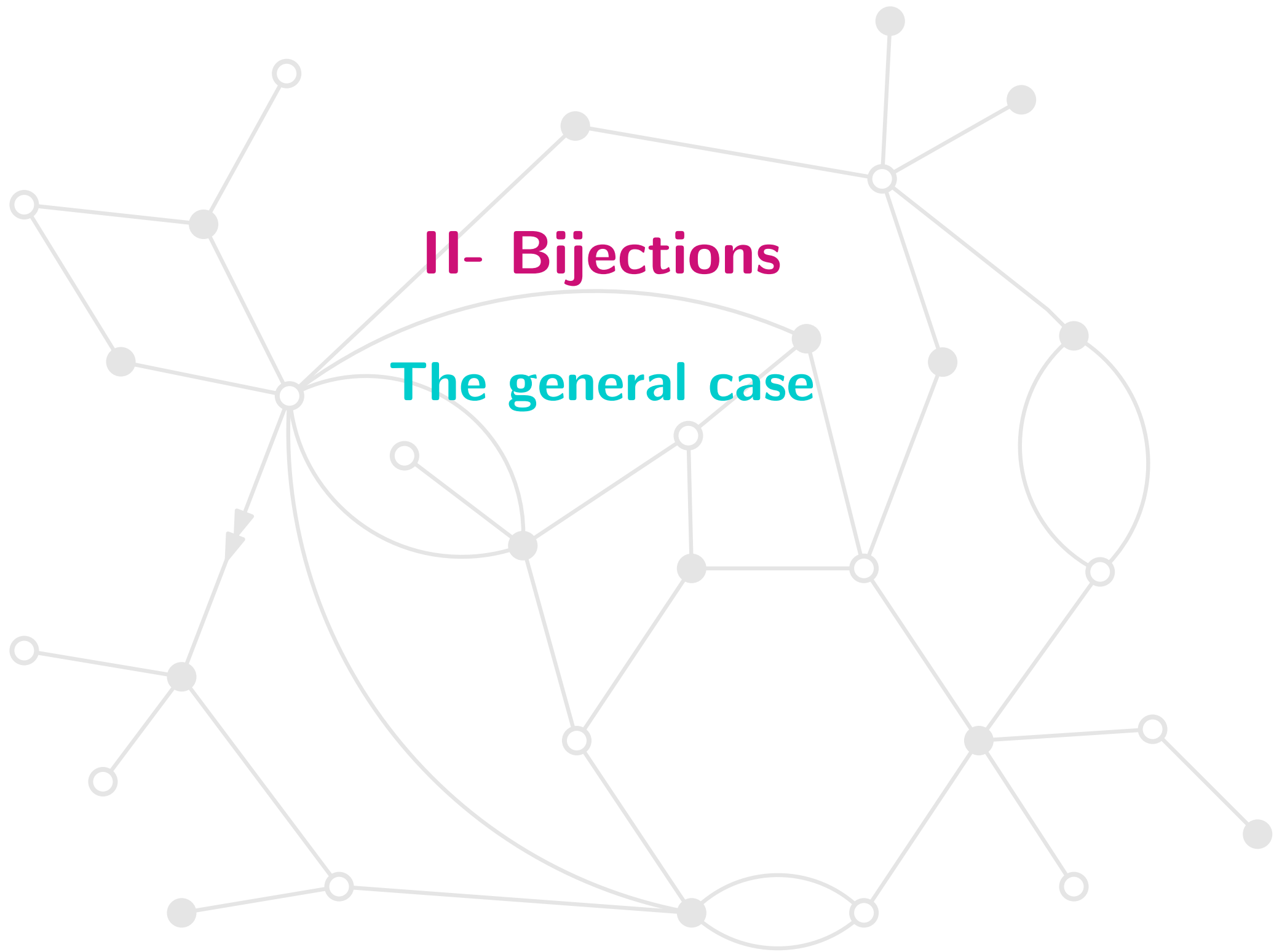


The 2 faces case: Cut and Close



The 2 faces case: Cut and Close





General case

$$4(f(\mathbf{d}) - 1)B(\mathbf{d}) = \sum_{\mathbf{s}+\mathbf{t}=\mathbf{d}} v(\mathbf{s})v(\mathbf{t})B(\mathbf{s})B(\mathbf{t})$$

Problems:

- ▶ Sew: Which path to slit?
- ▶ Cut: Which cycle to cut?

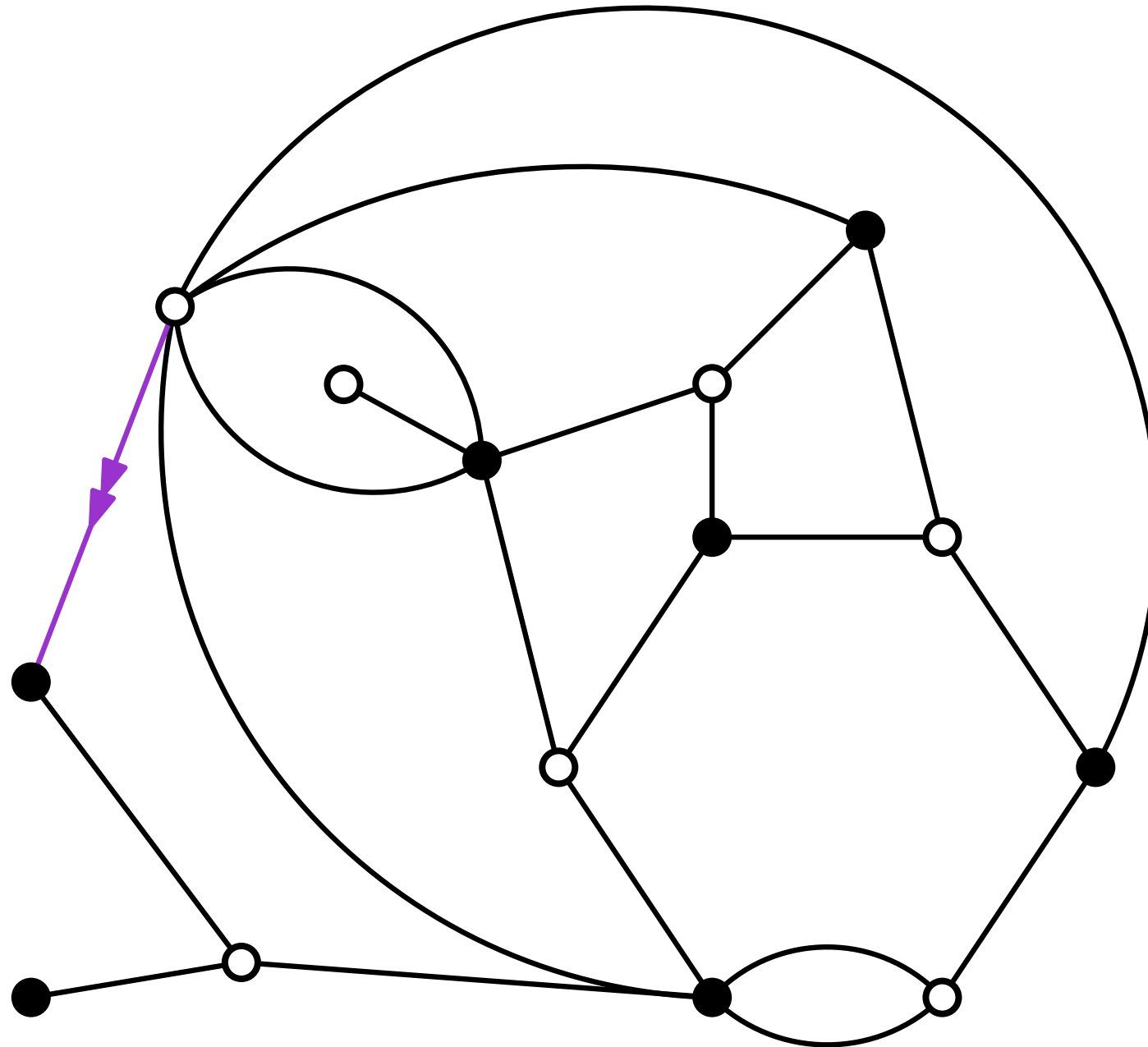
General case

$$4(f(\mathbf{d}) - 1)B(\mathbf{d}) = \sum_{\mathbf{s} + \mathbf{t} = \mathbf{d}} v(\mathbf{s})v(\mathbf{t})B(\mathbf{s})B(\mathbf{t})$$

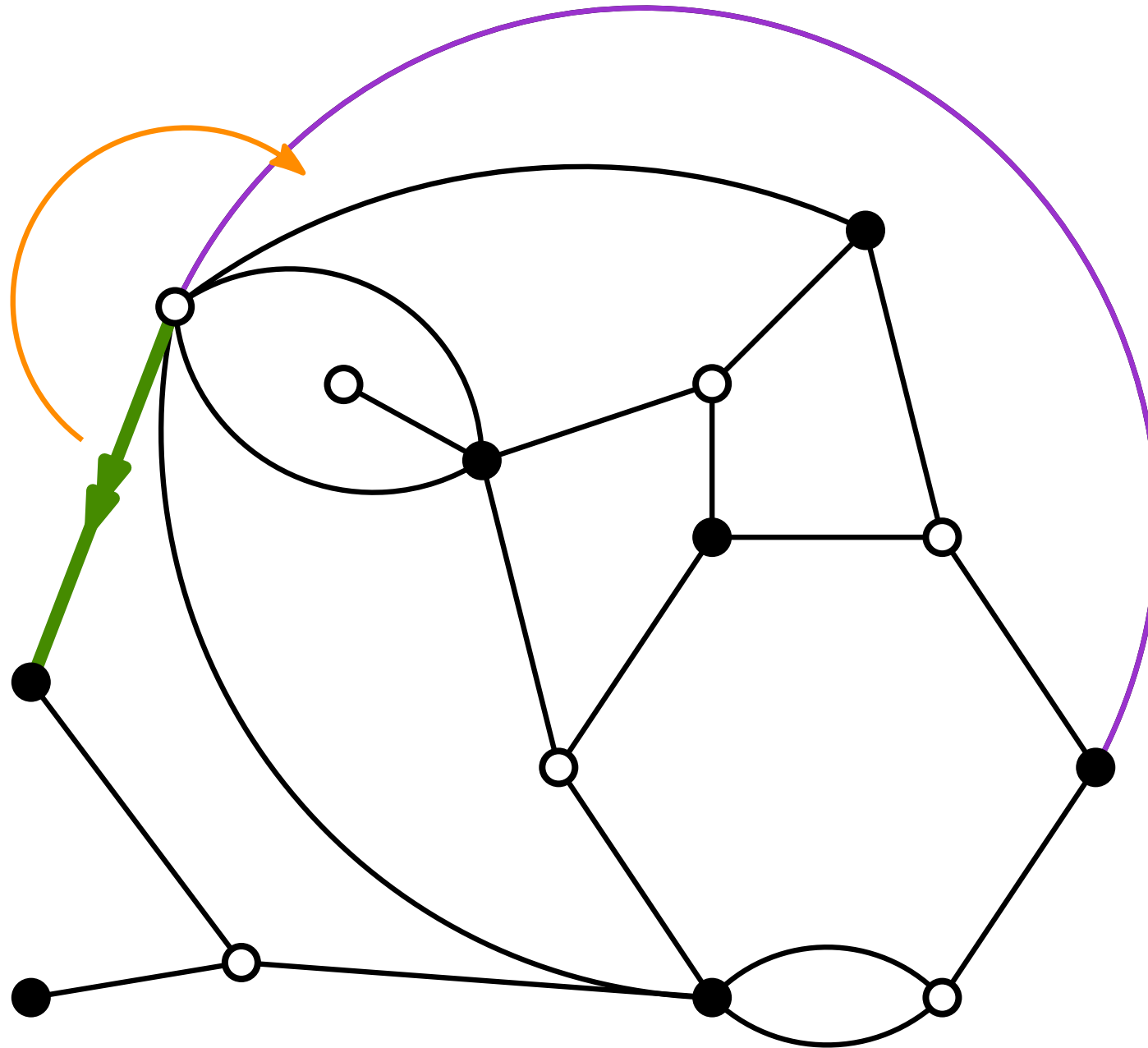
Problems:

- ▶ Sew: Which path to slit?
- ▶ Cut: Which cycle to cut?
- ⇒ Take a spanning tree!
 - ↪ Canonical path to slit.
 - ↪ $2(f(\mathbf{d}) - 1) =$ an oriented edge outside of the spanning tree.

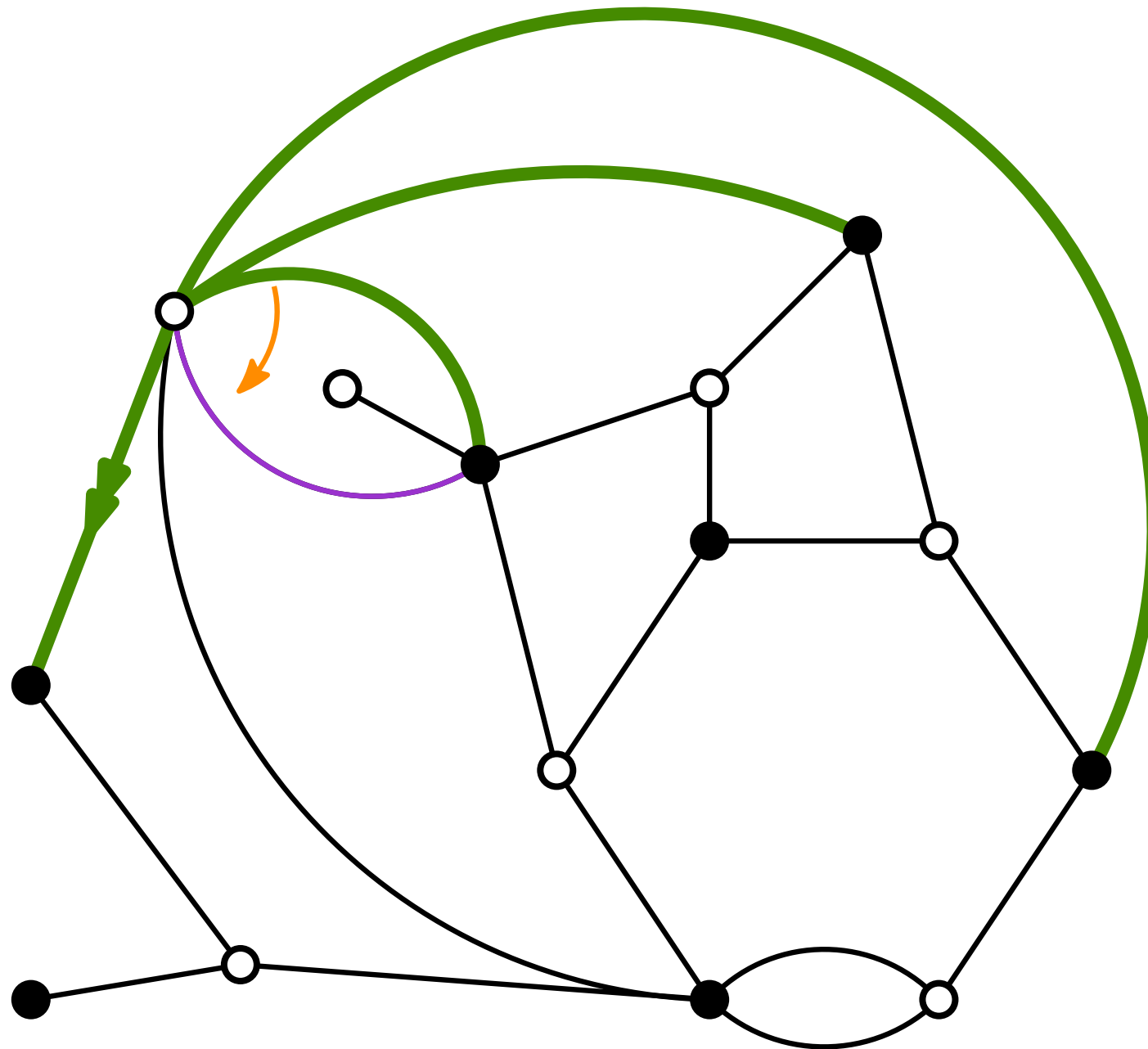
Breadth-first search tree



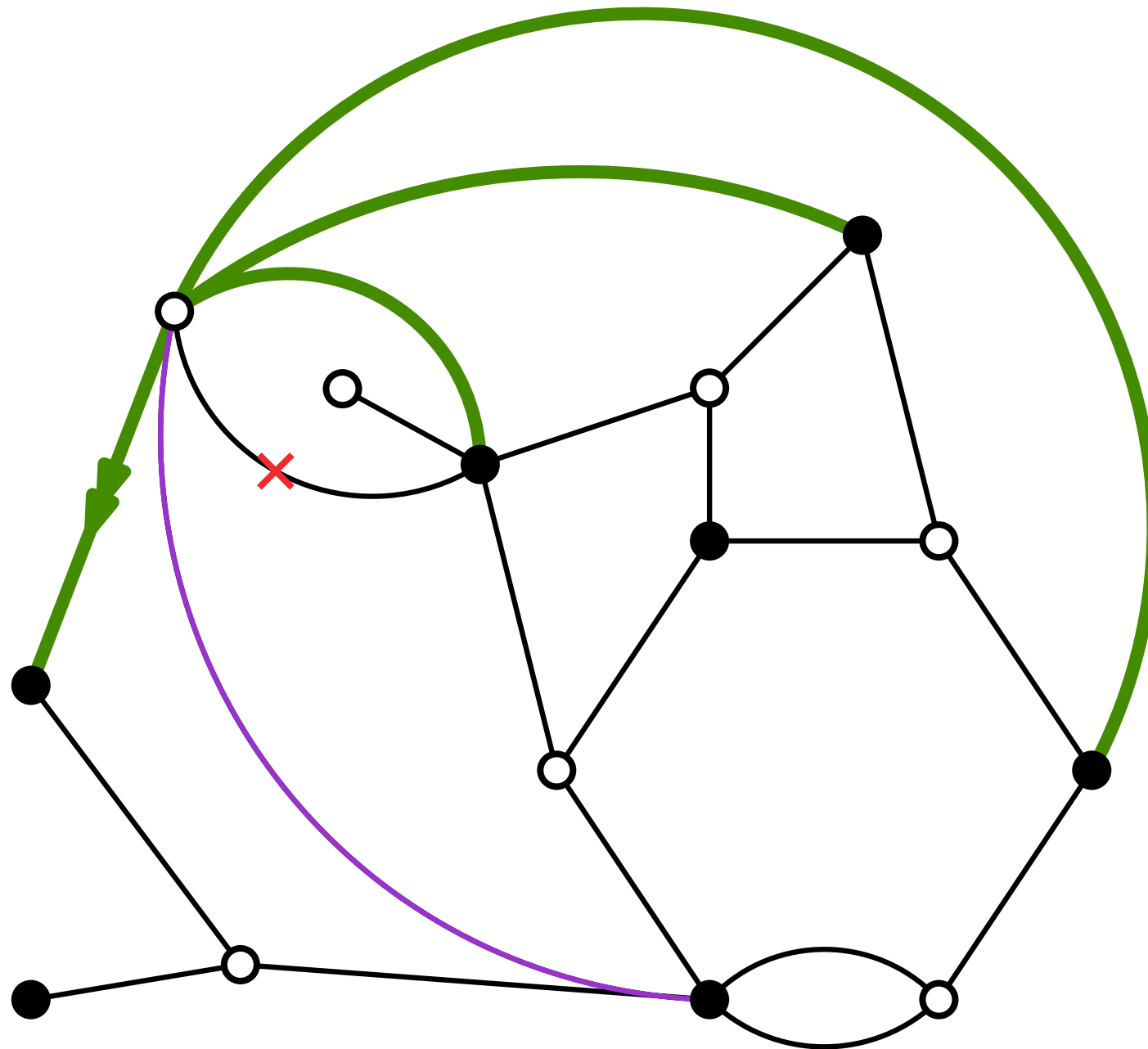
Breadth-first search tree



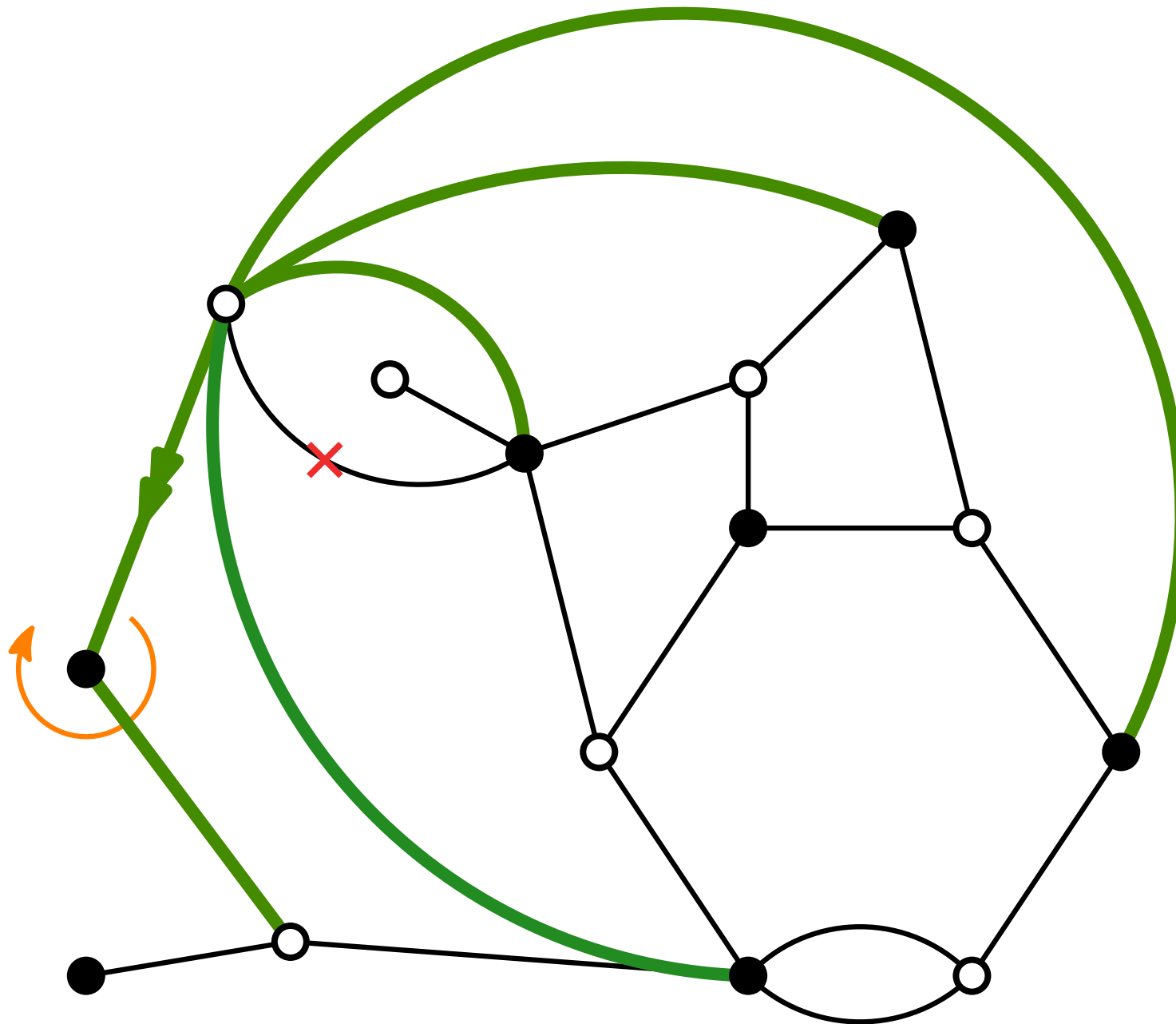
Breadth-first search tree



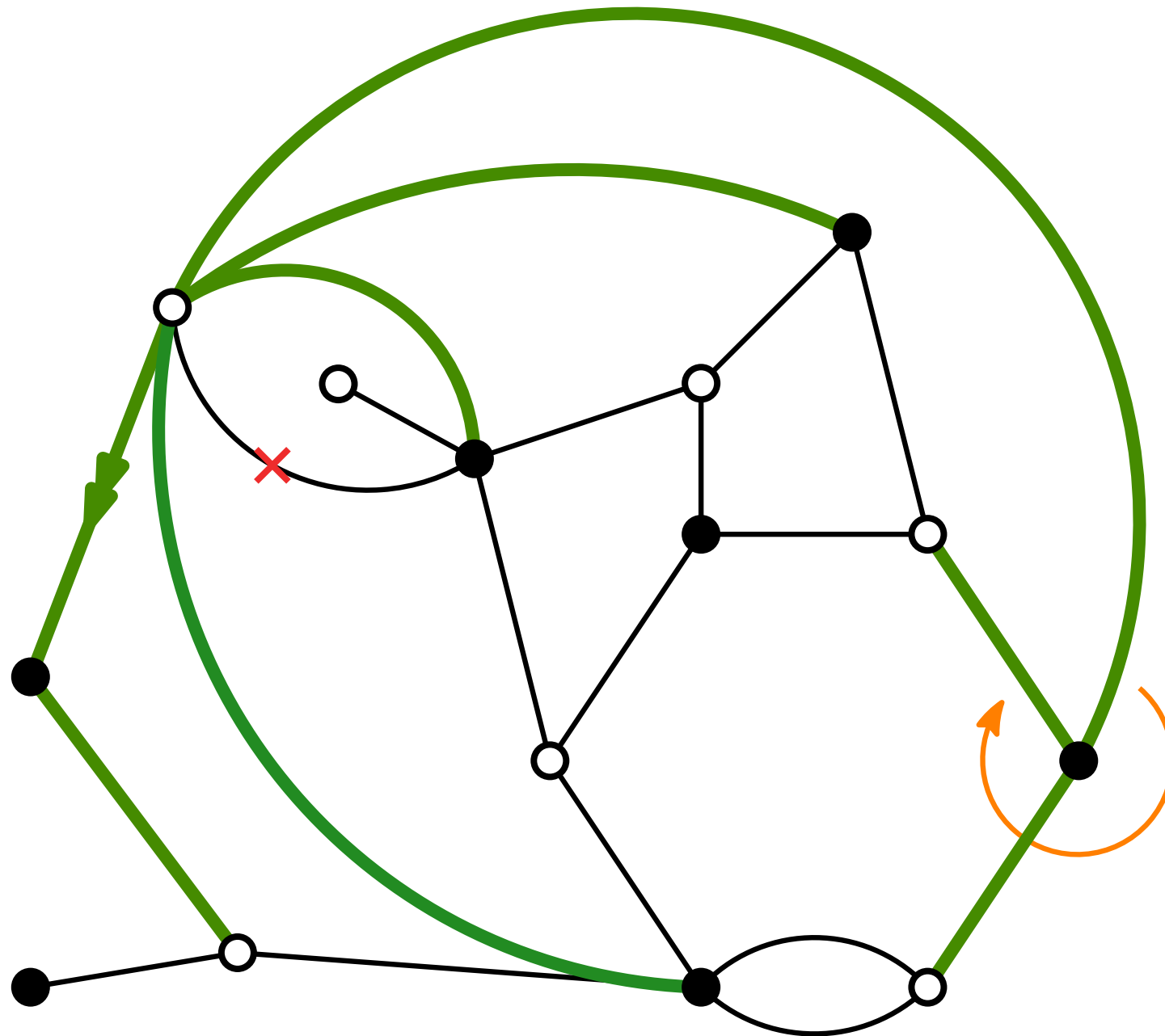
Breadth-first search tree



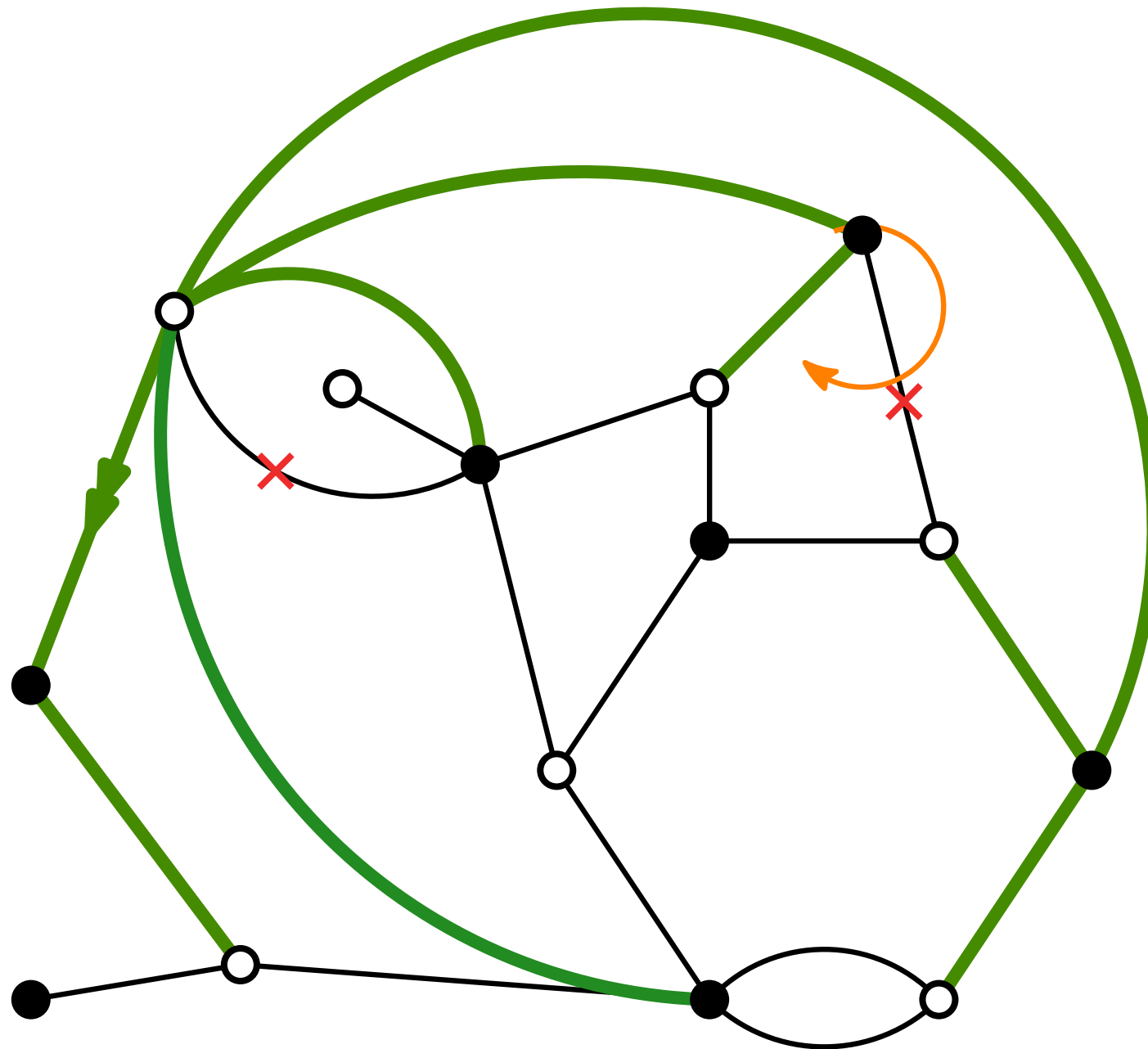
Breadth-first search tree



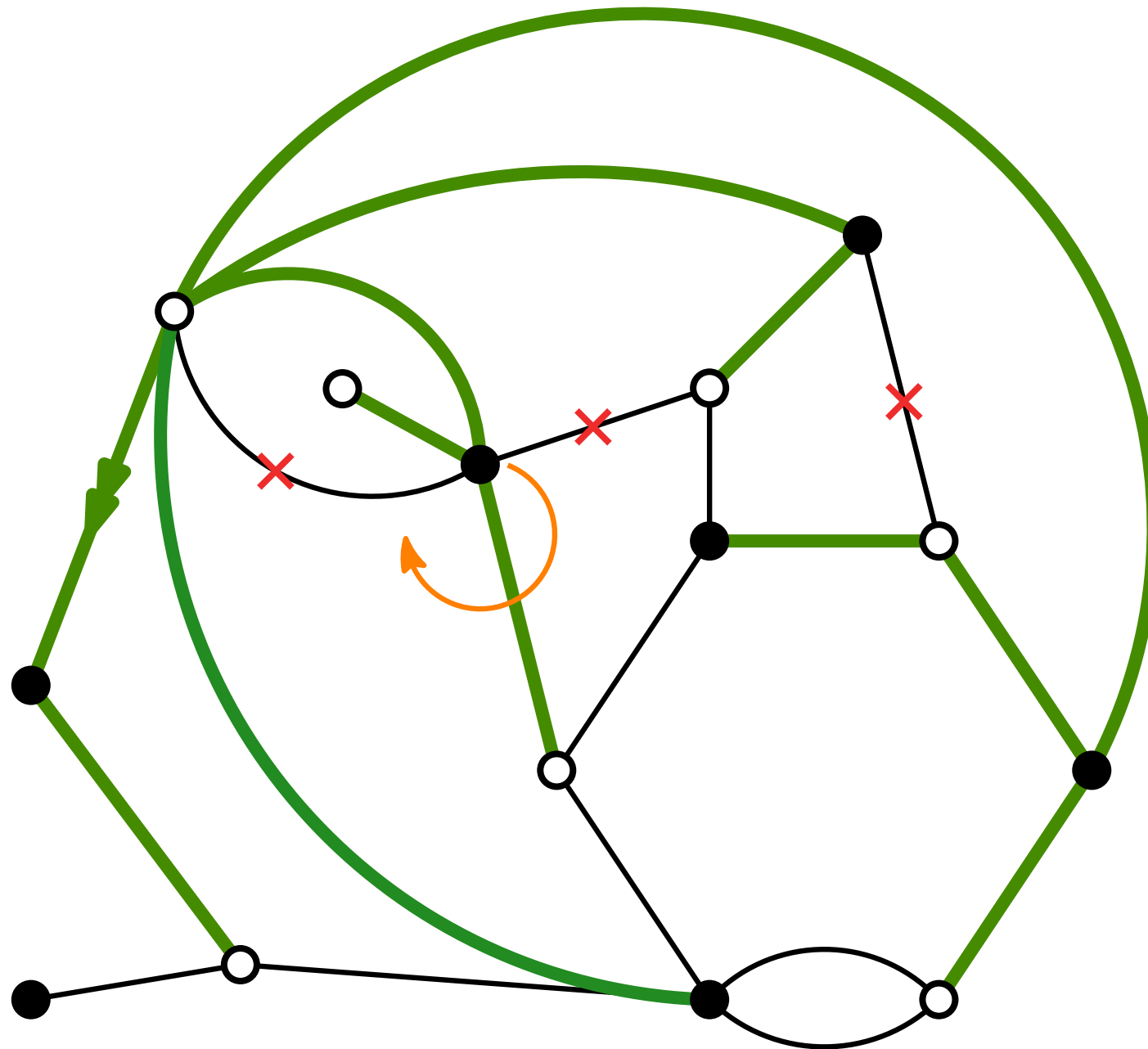
Breadth-first search tree



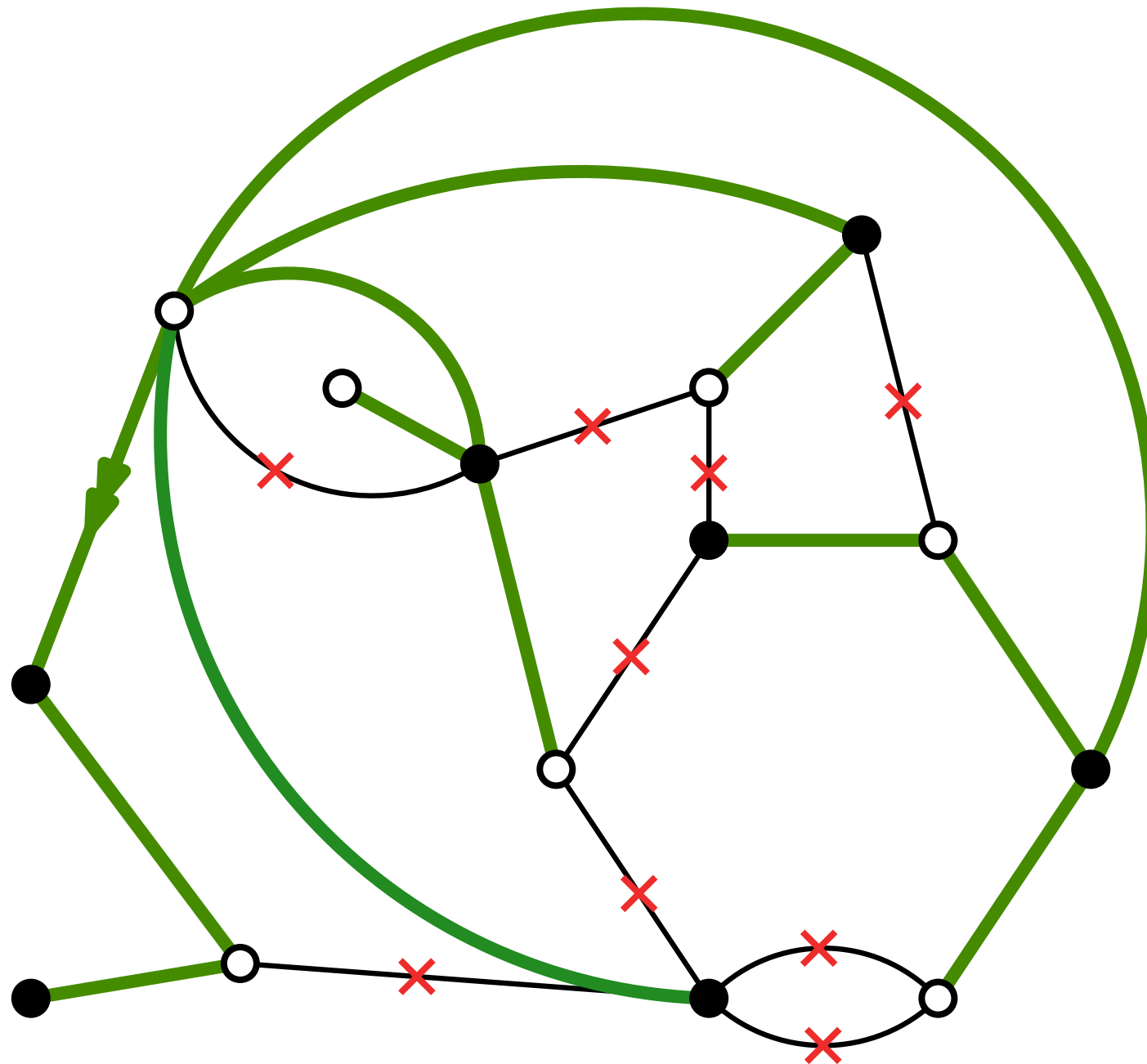
Breadth-first search tree



Breadth-first search tree

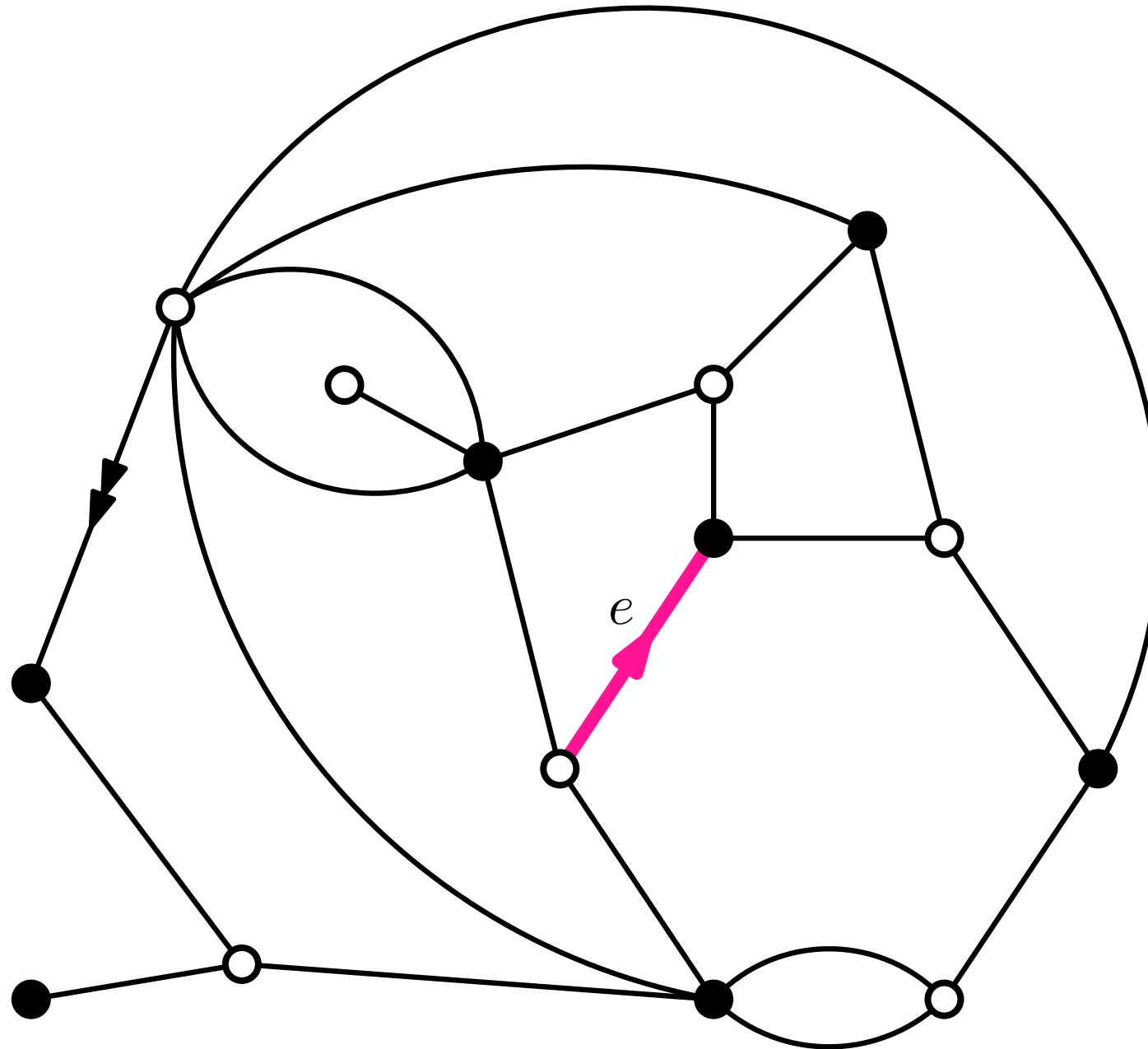


Breadth-first search tree



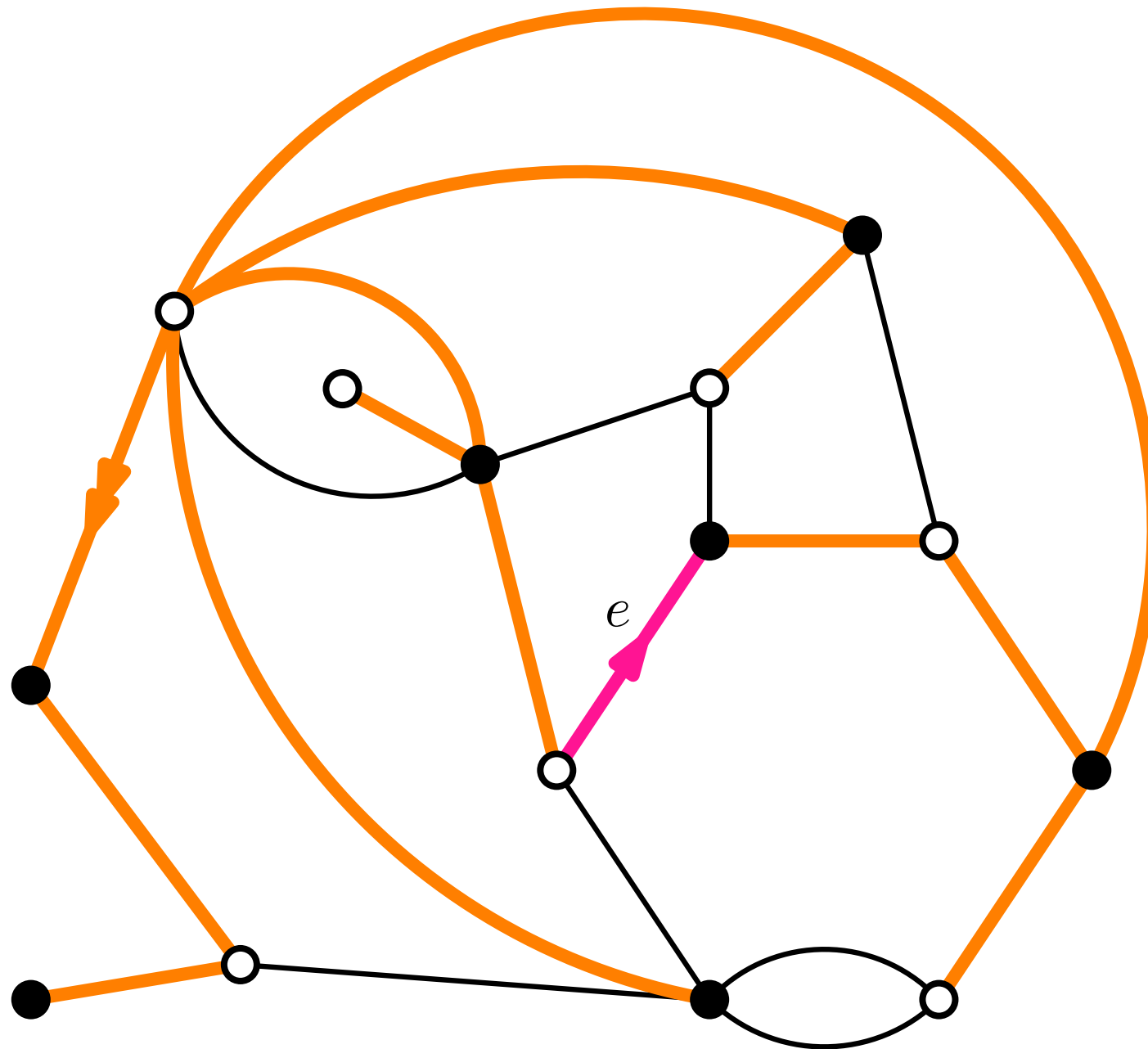
General case : Cut and Close

$(M, -)$



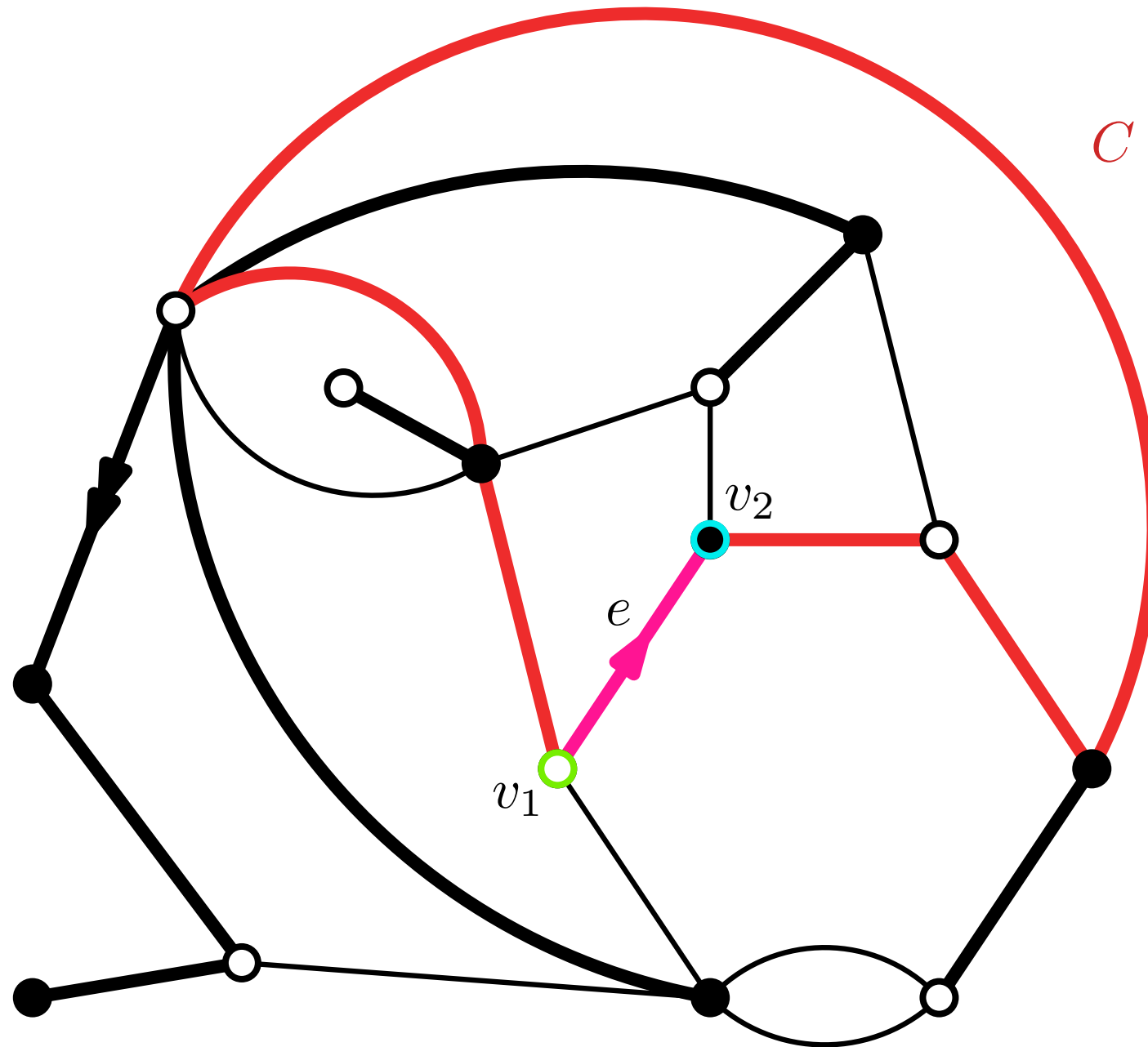
General case : Cut and Close

$(M, -)$

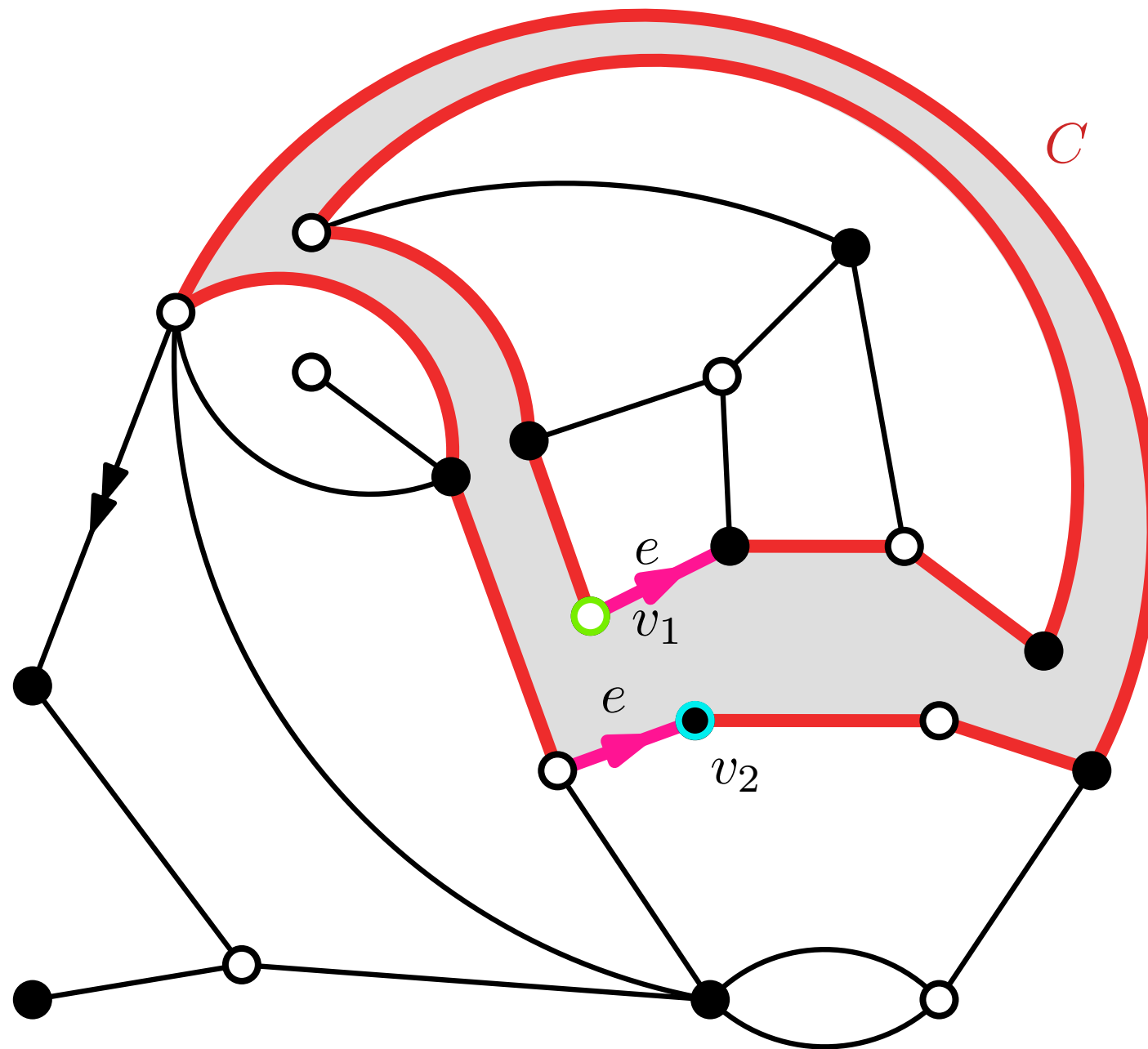


General case : Cut and Close

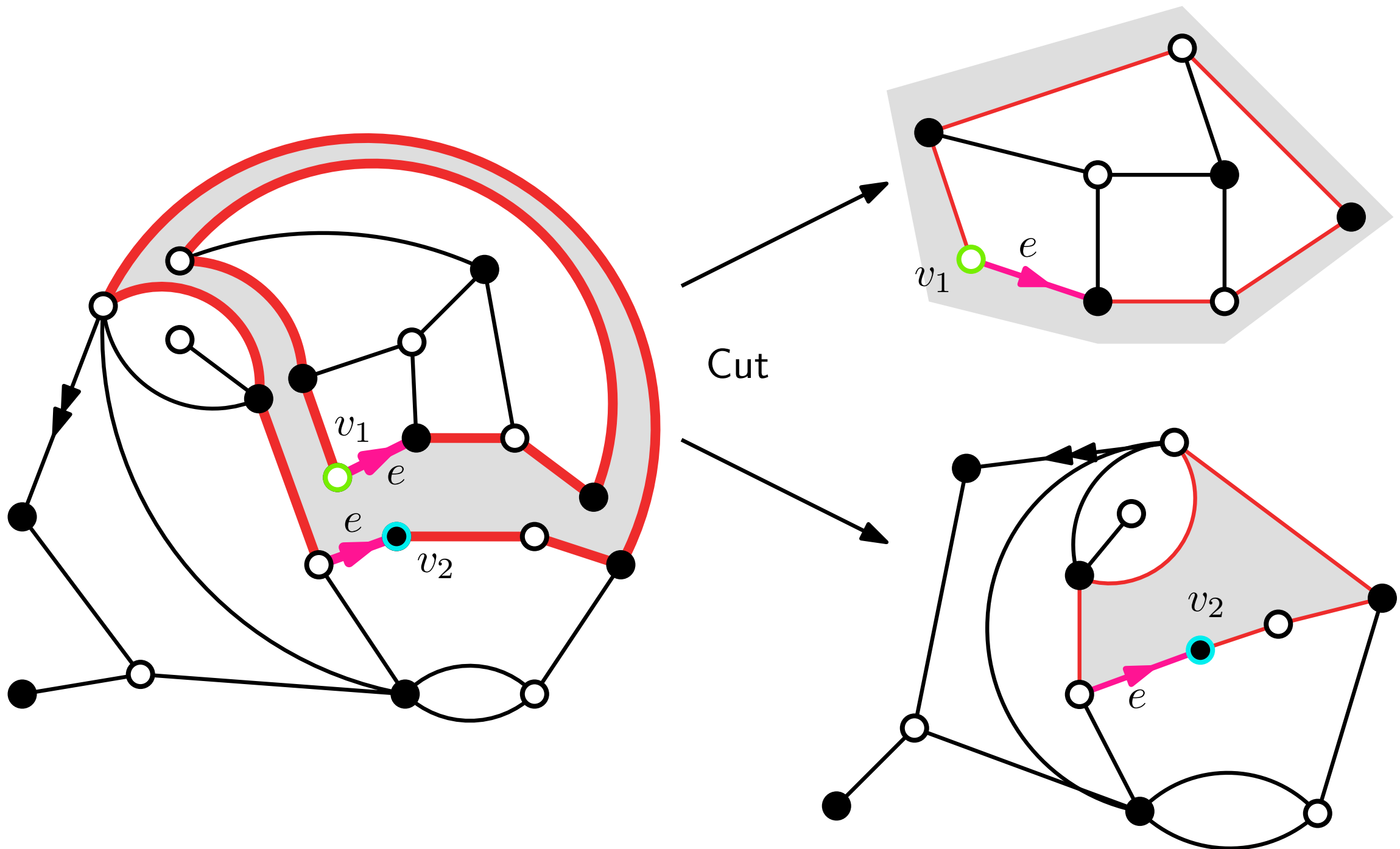
$(M, -)$



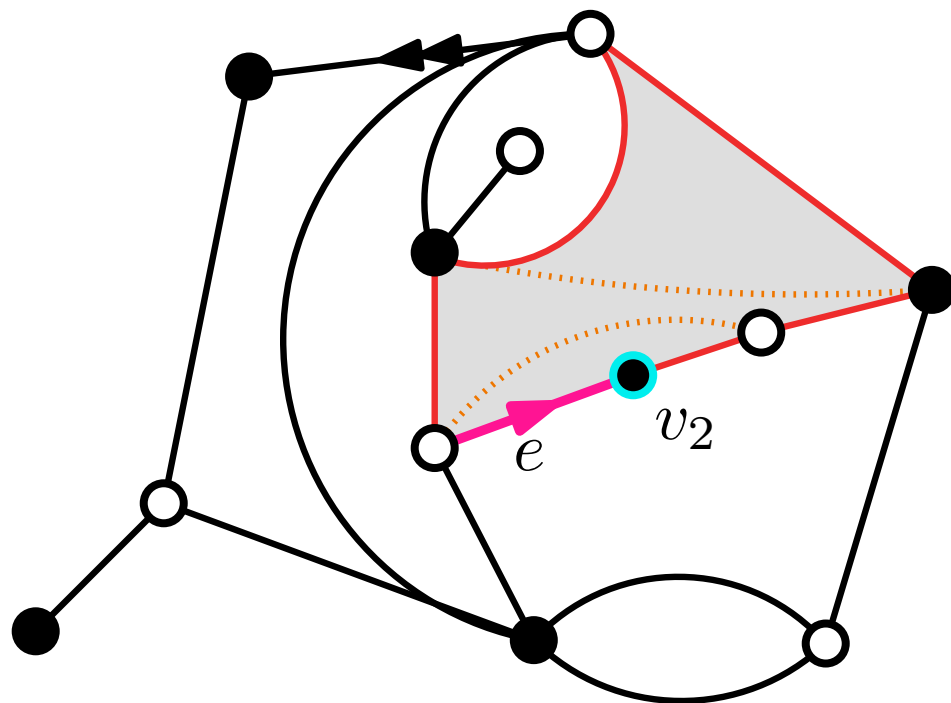
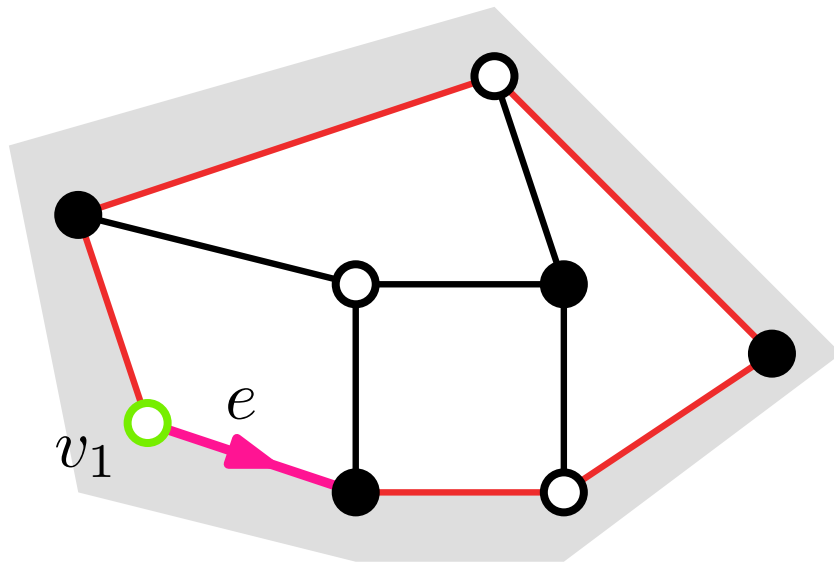
General case : Cut and Close



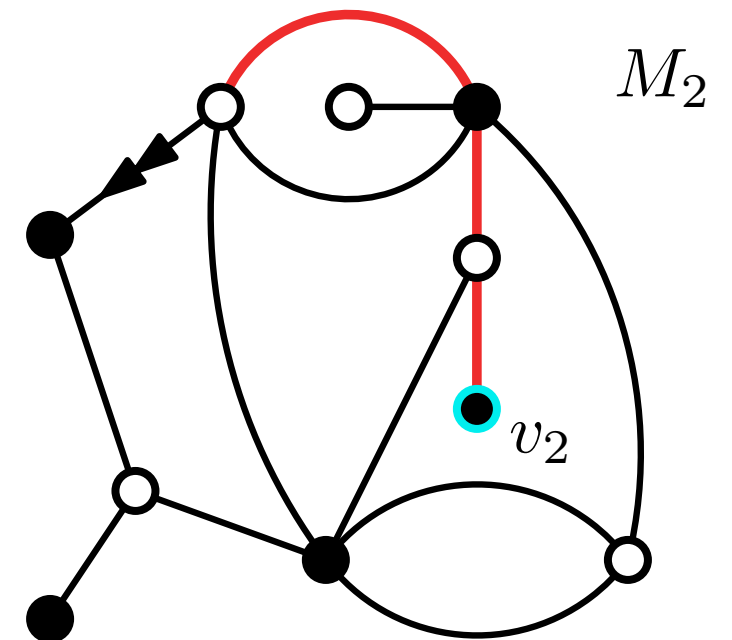
General case : Cut and Close



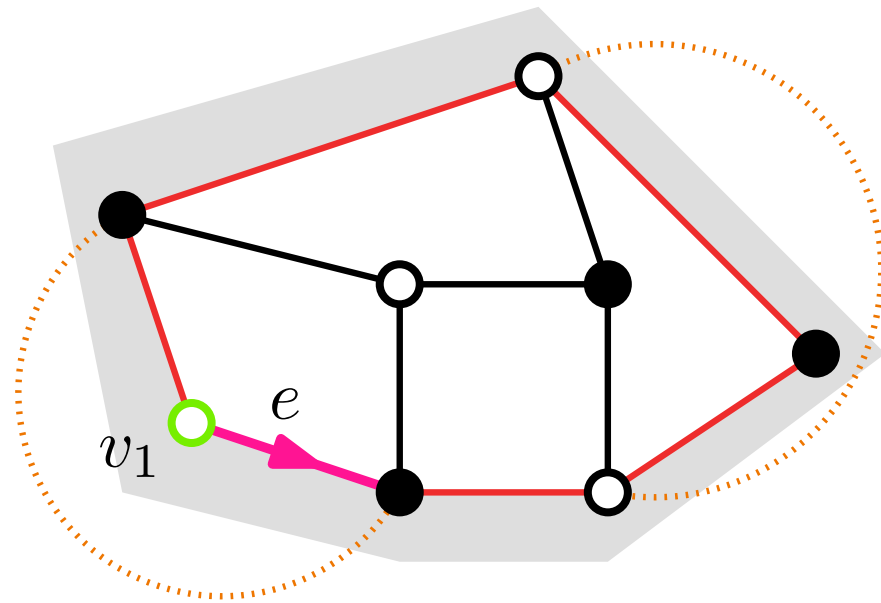
General case : Cut and Close



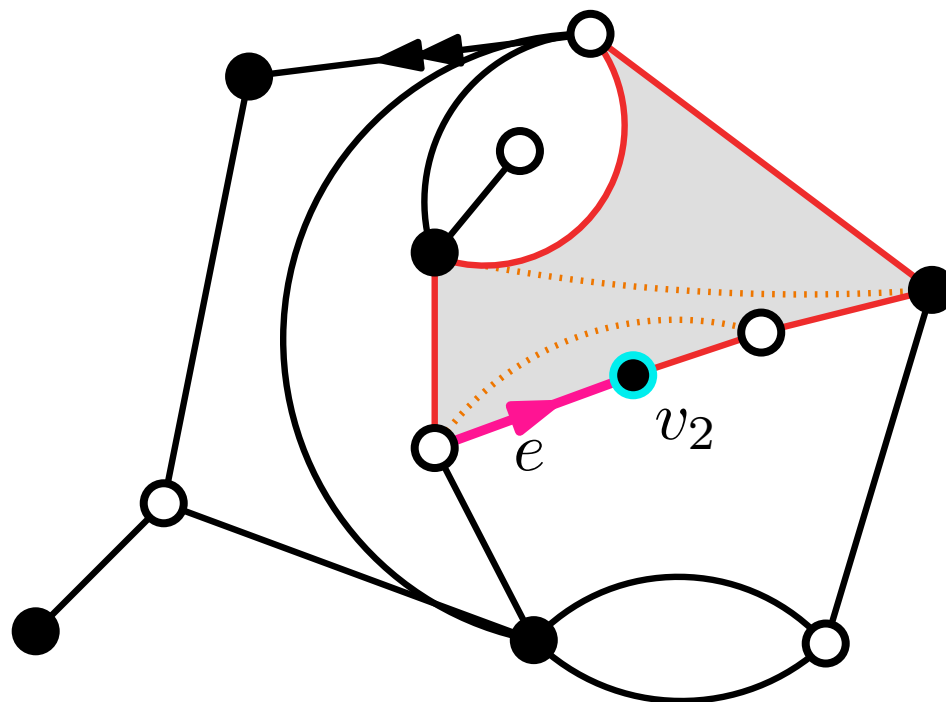
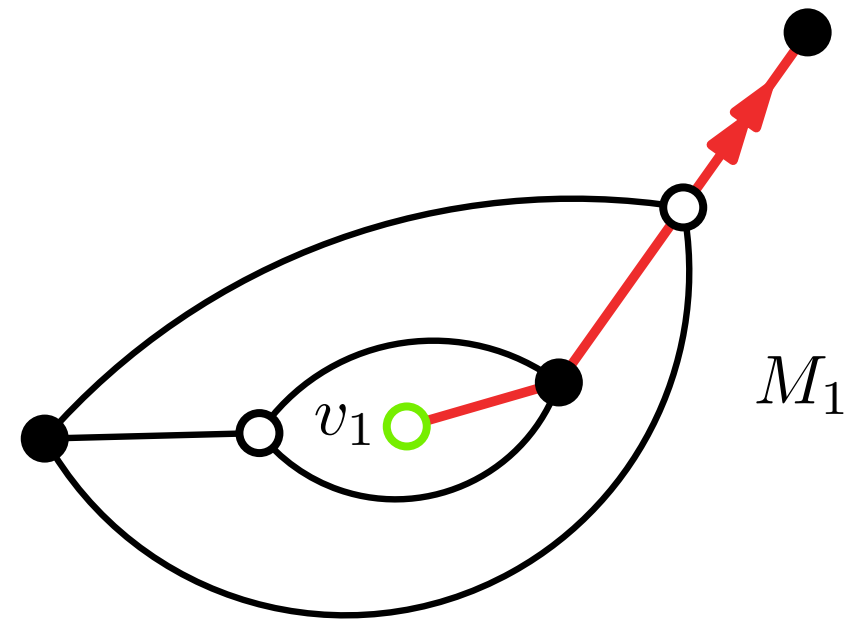
Close



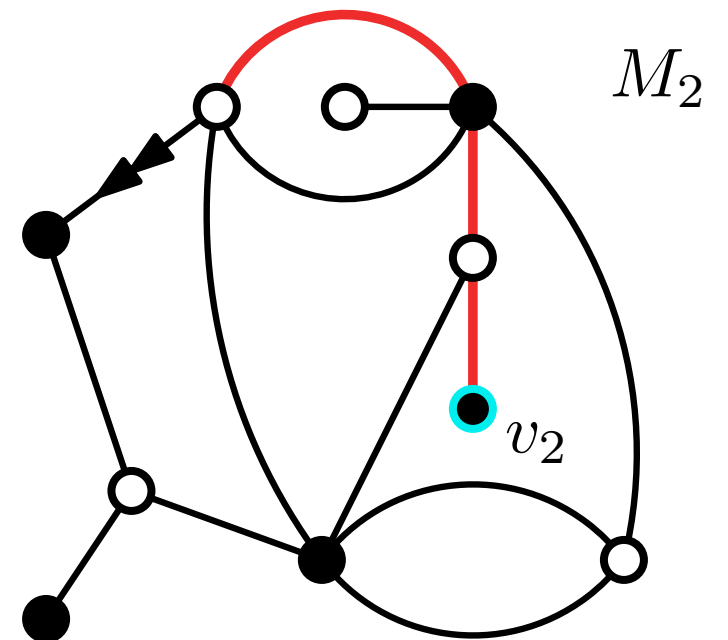
General case : Cut and Close



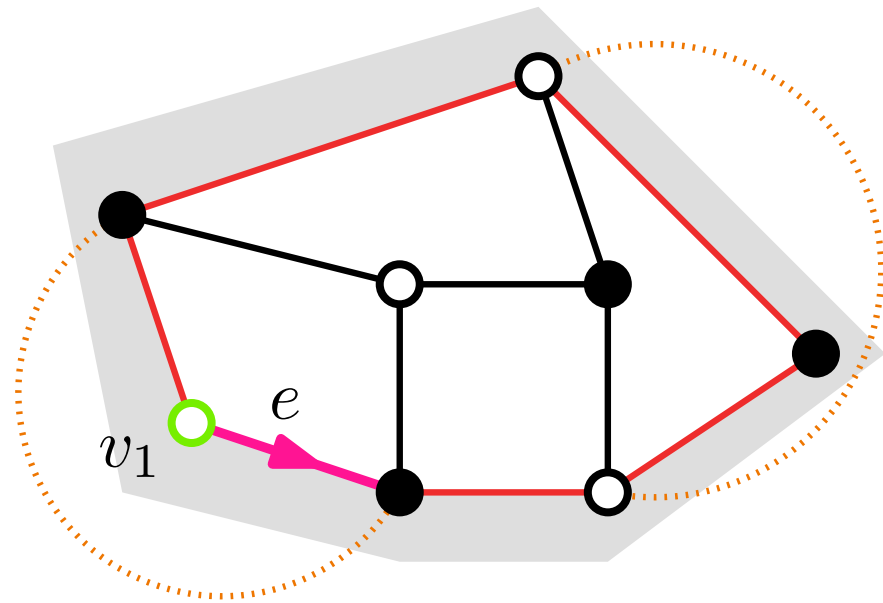
Close
→



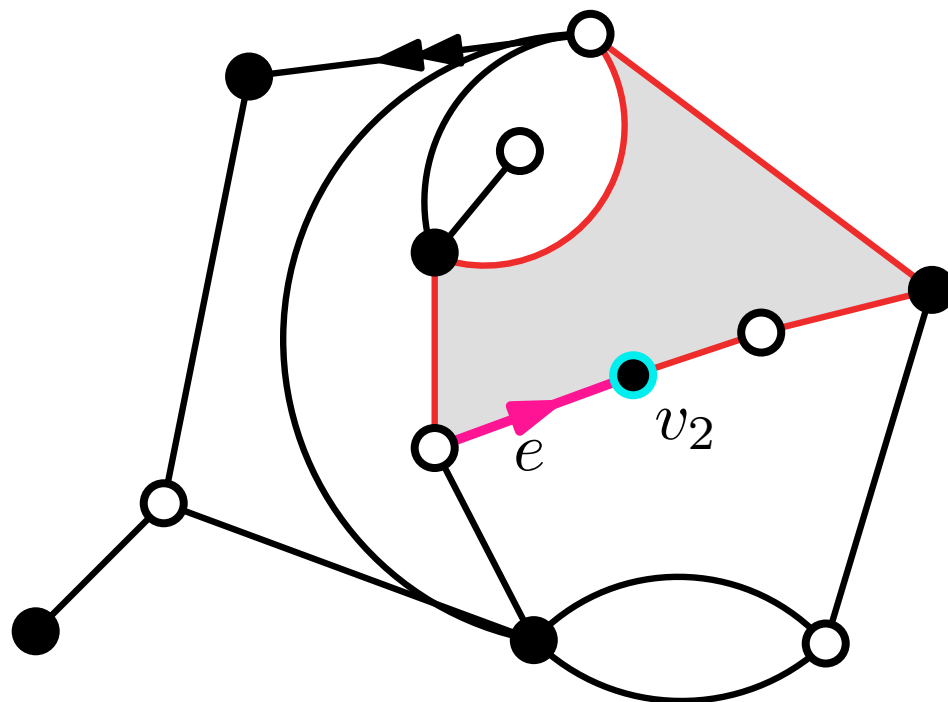
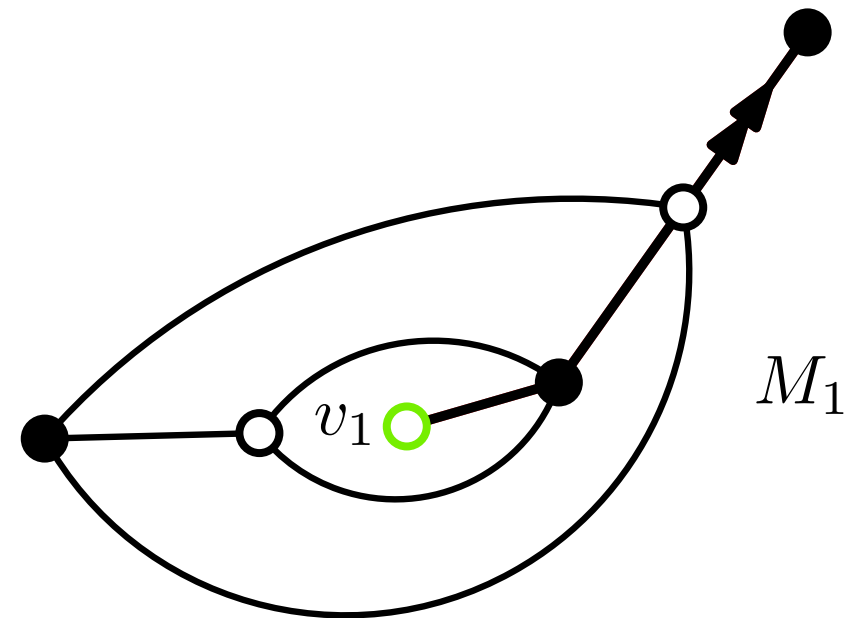
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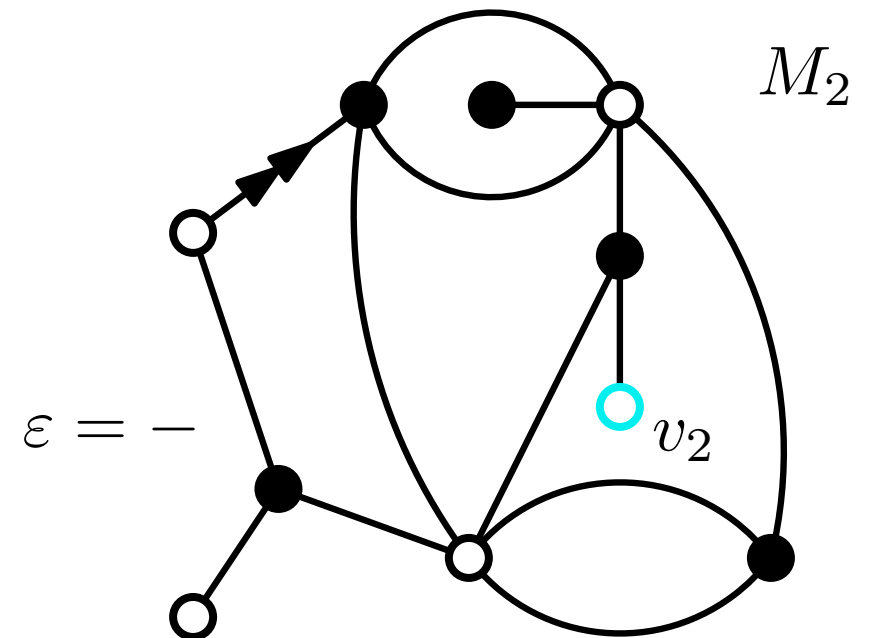
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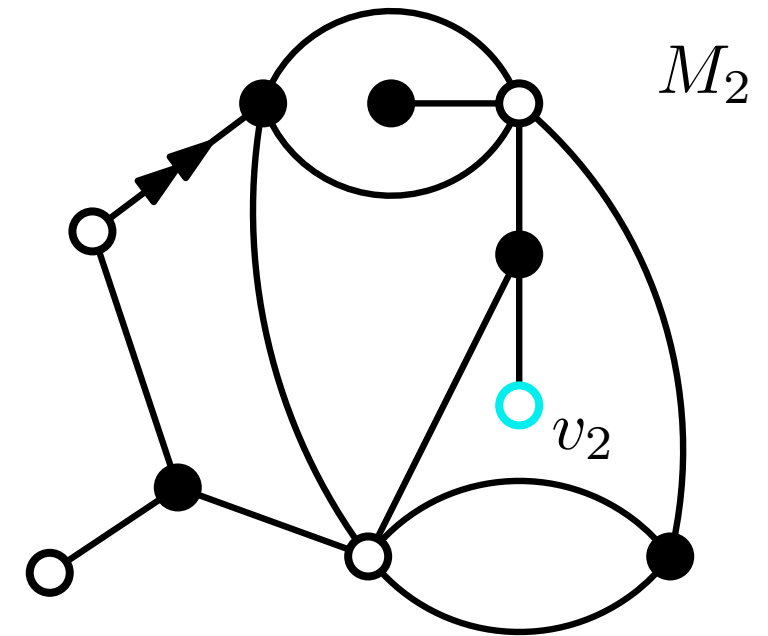
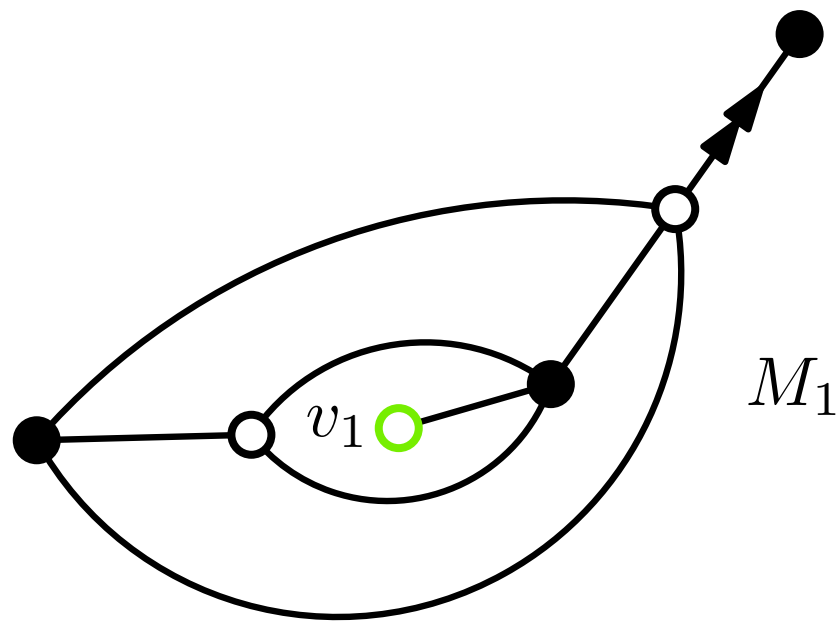
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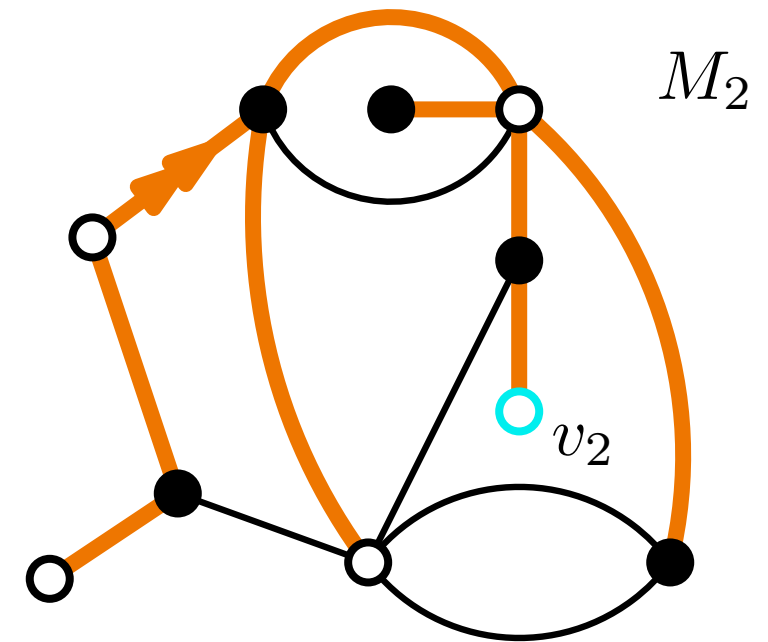
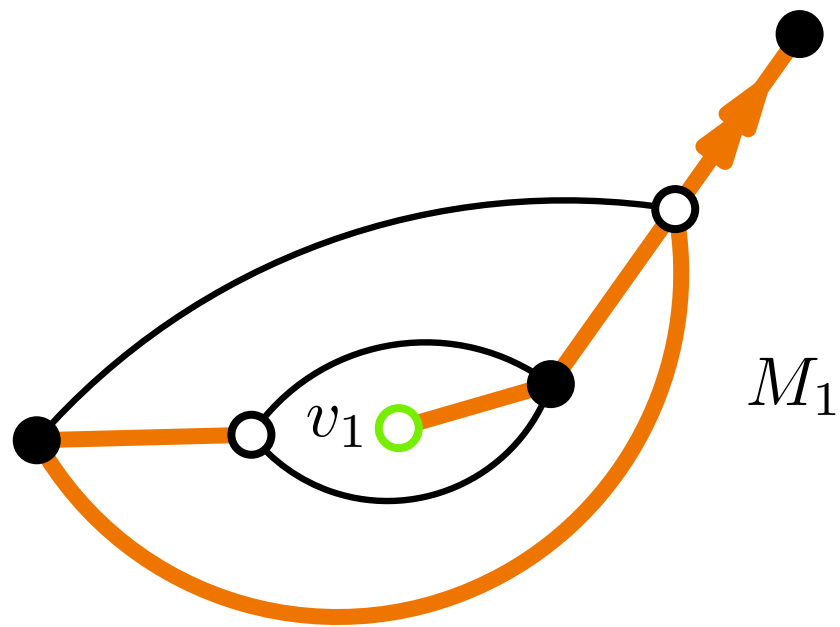
Close
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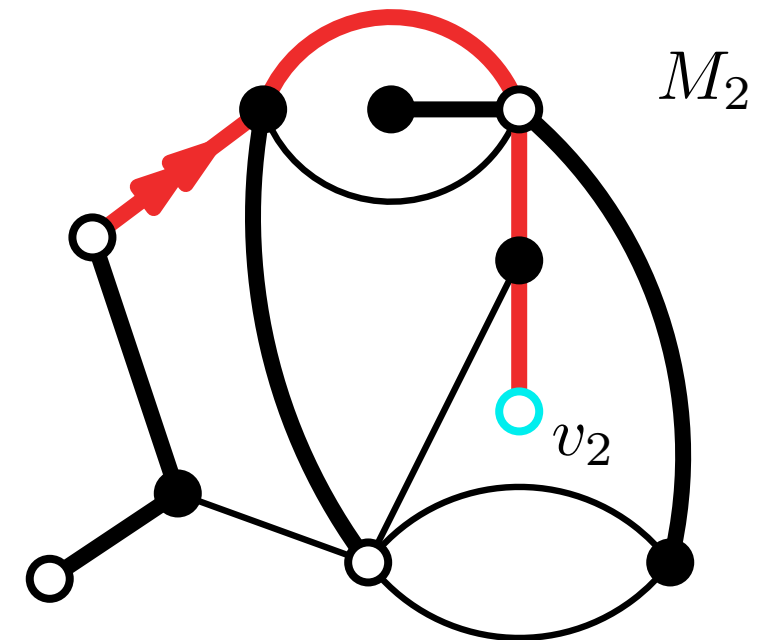
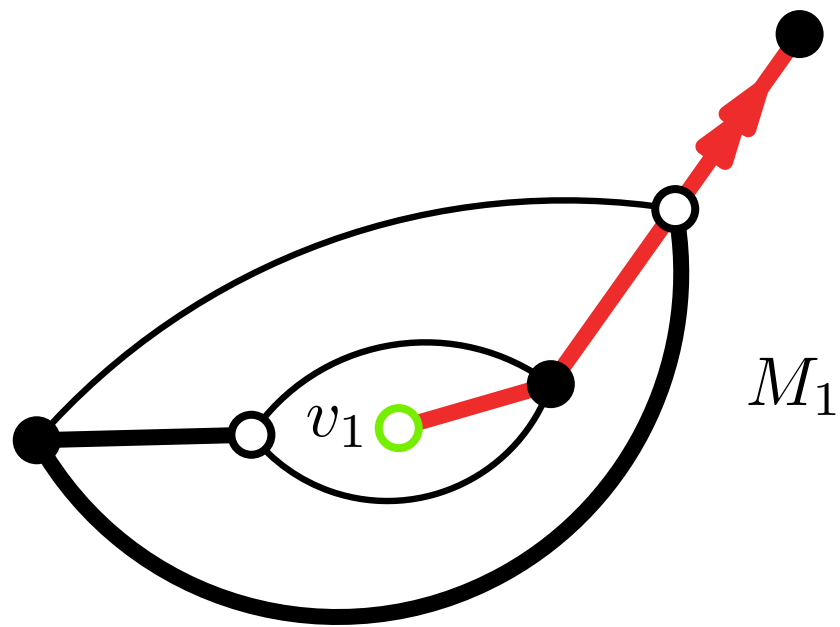
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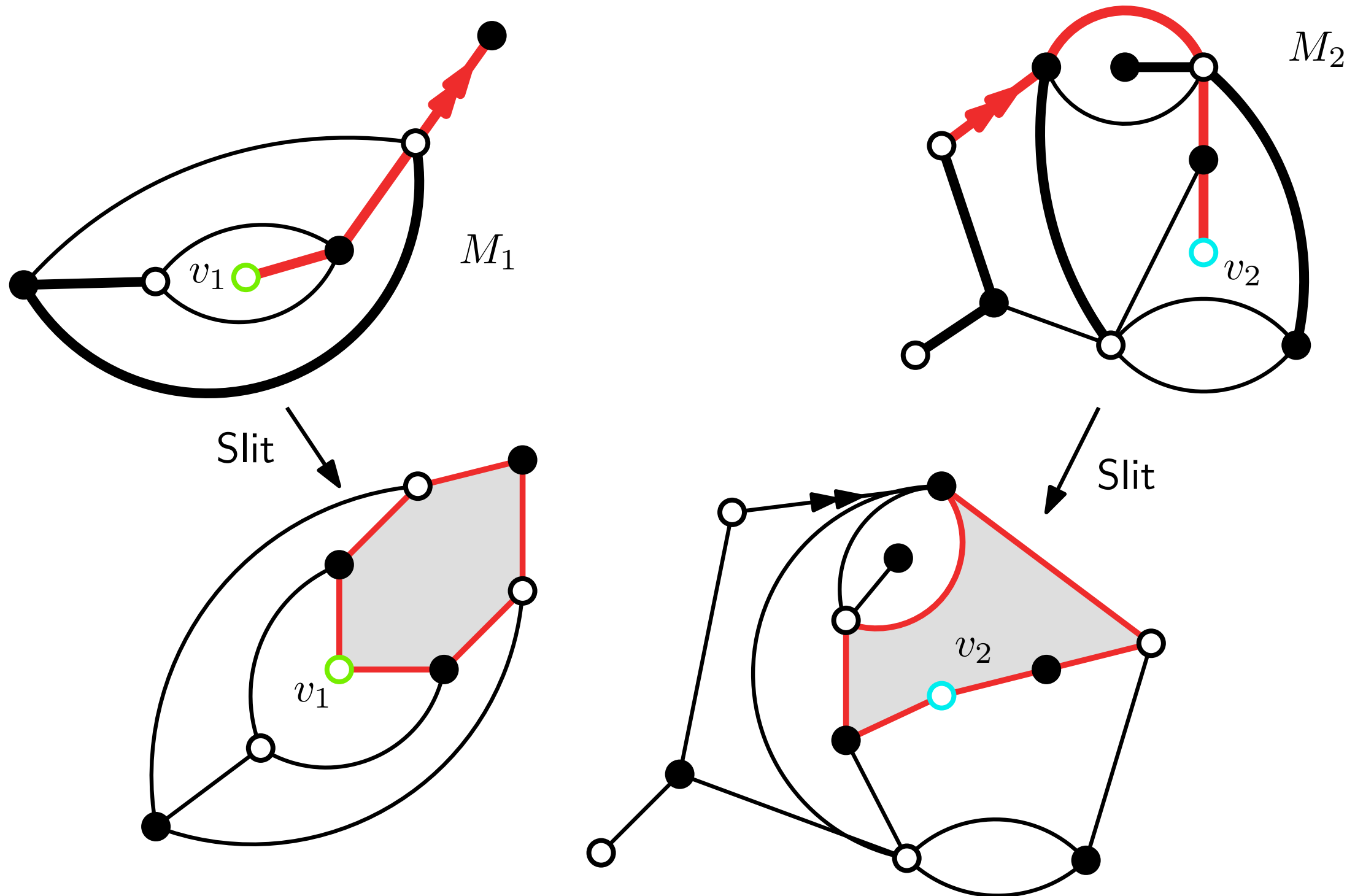
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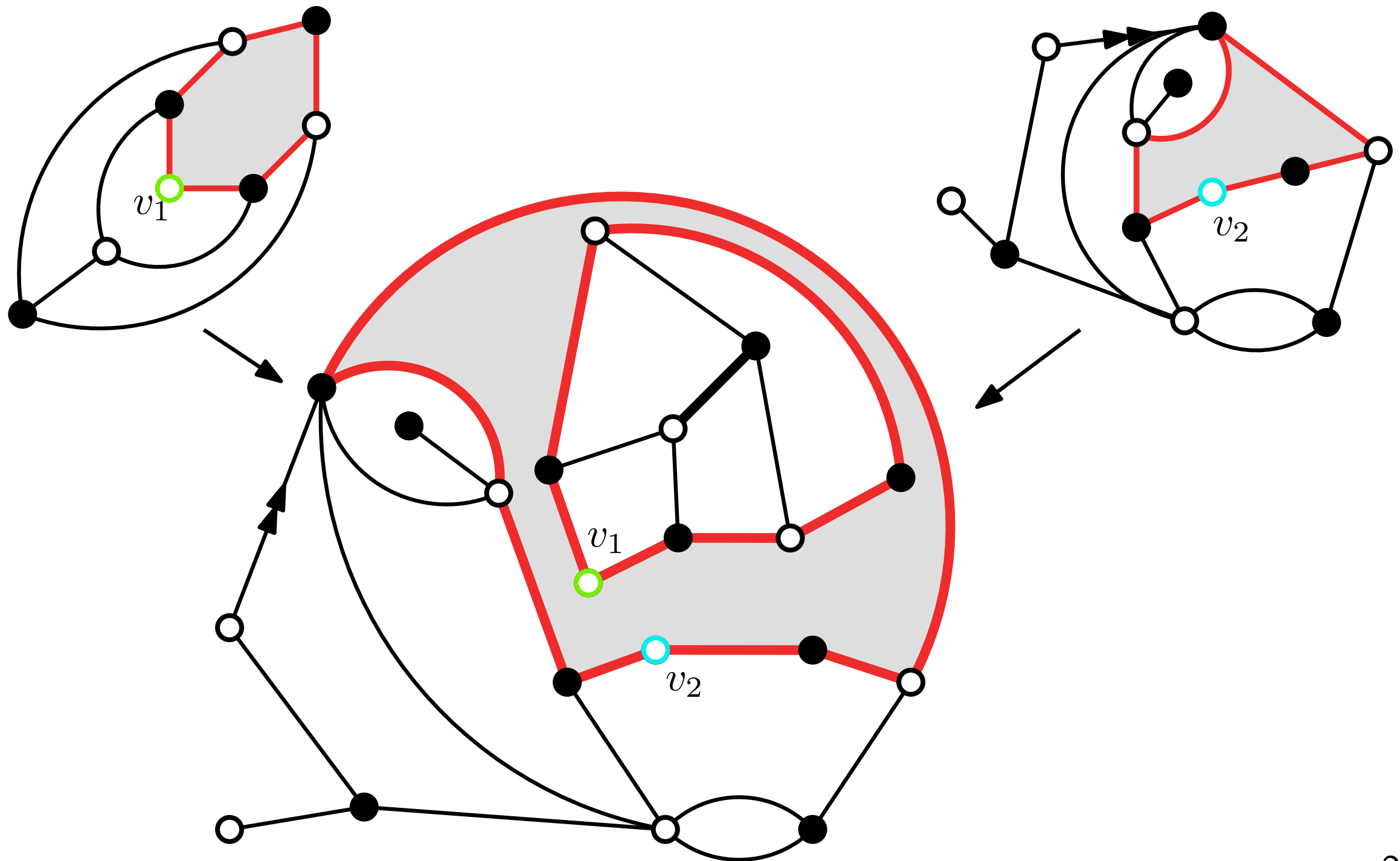
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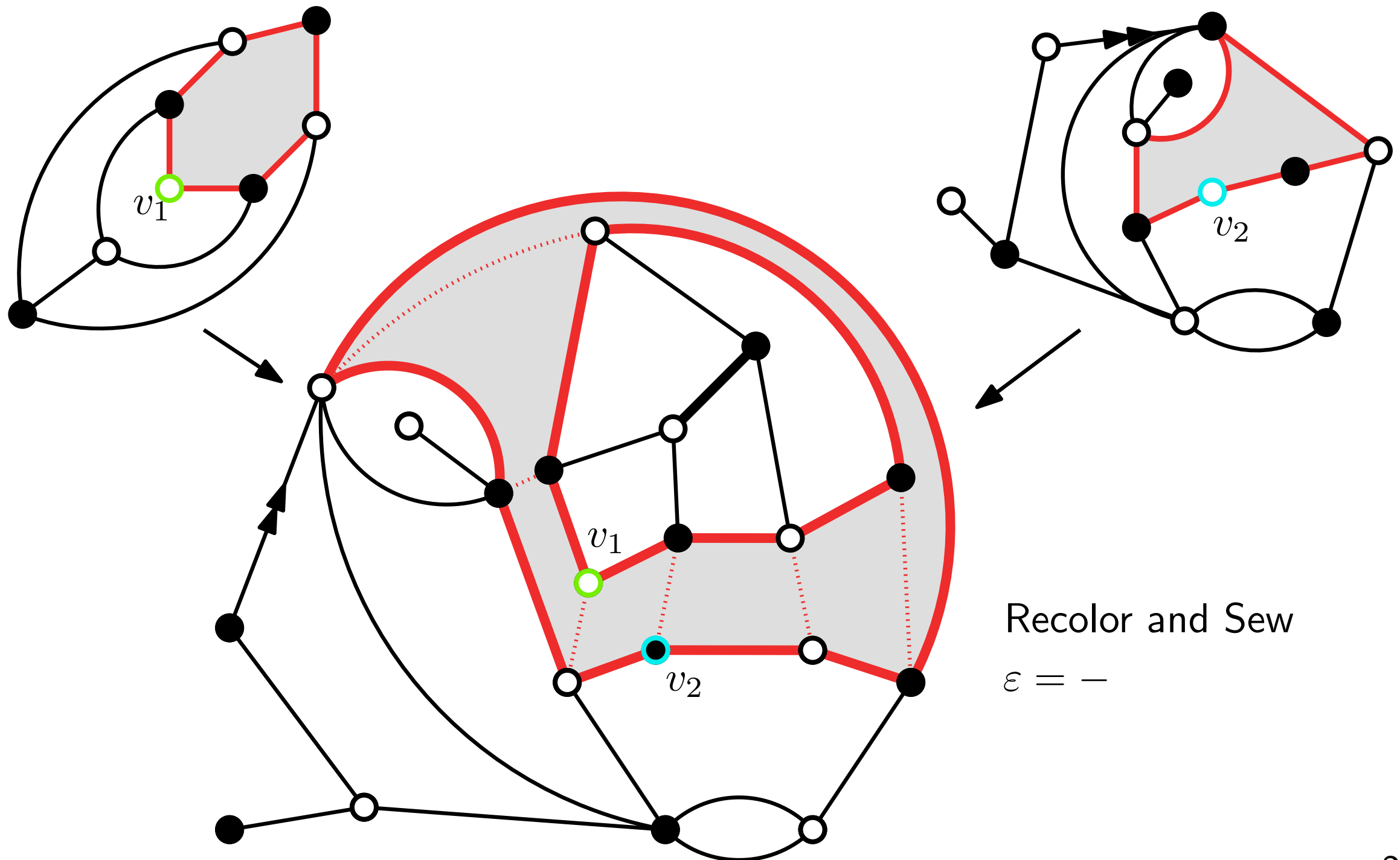
General case : Slit and Sew



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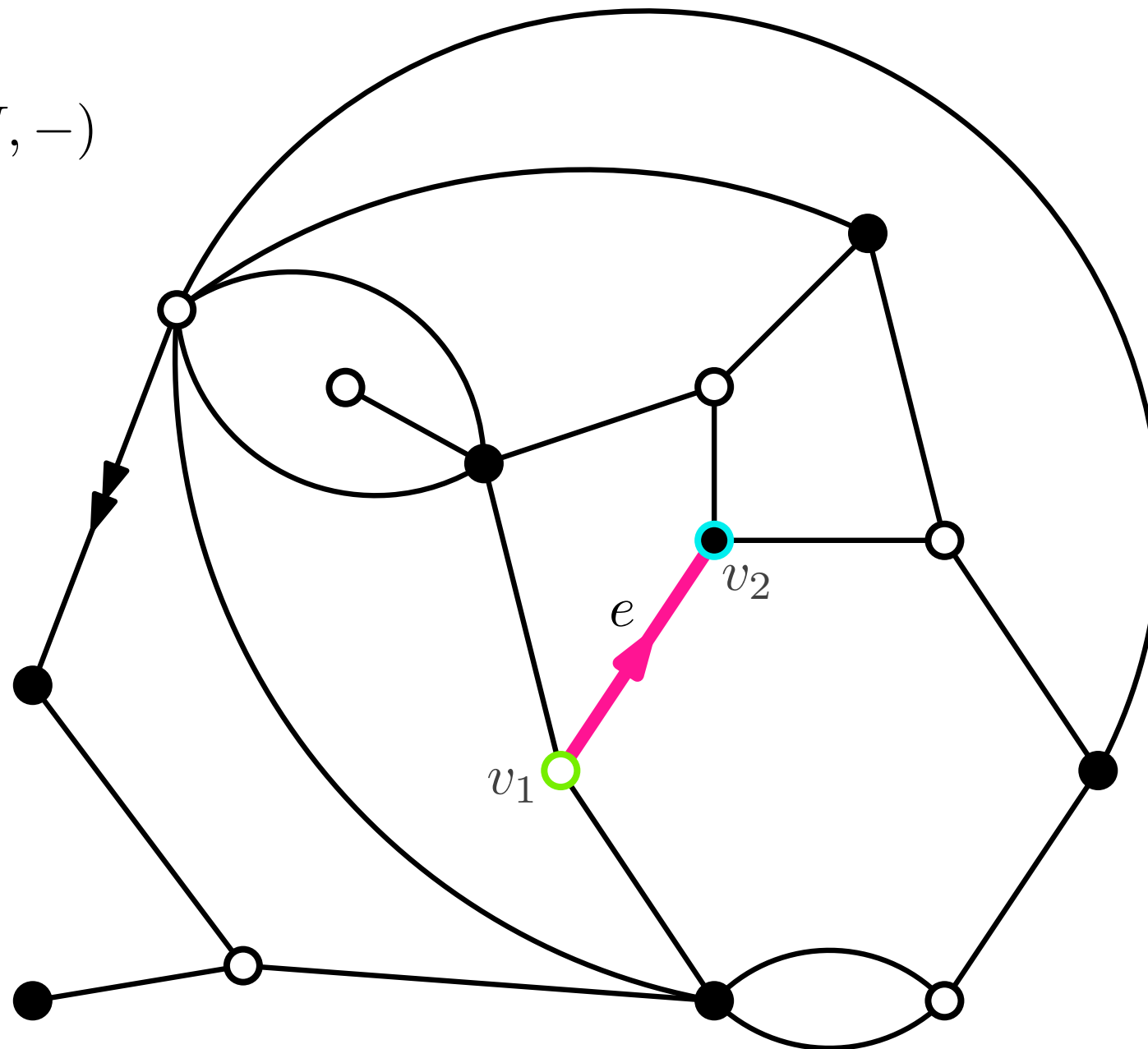


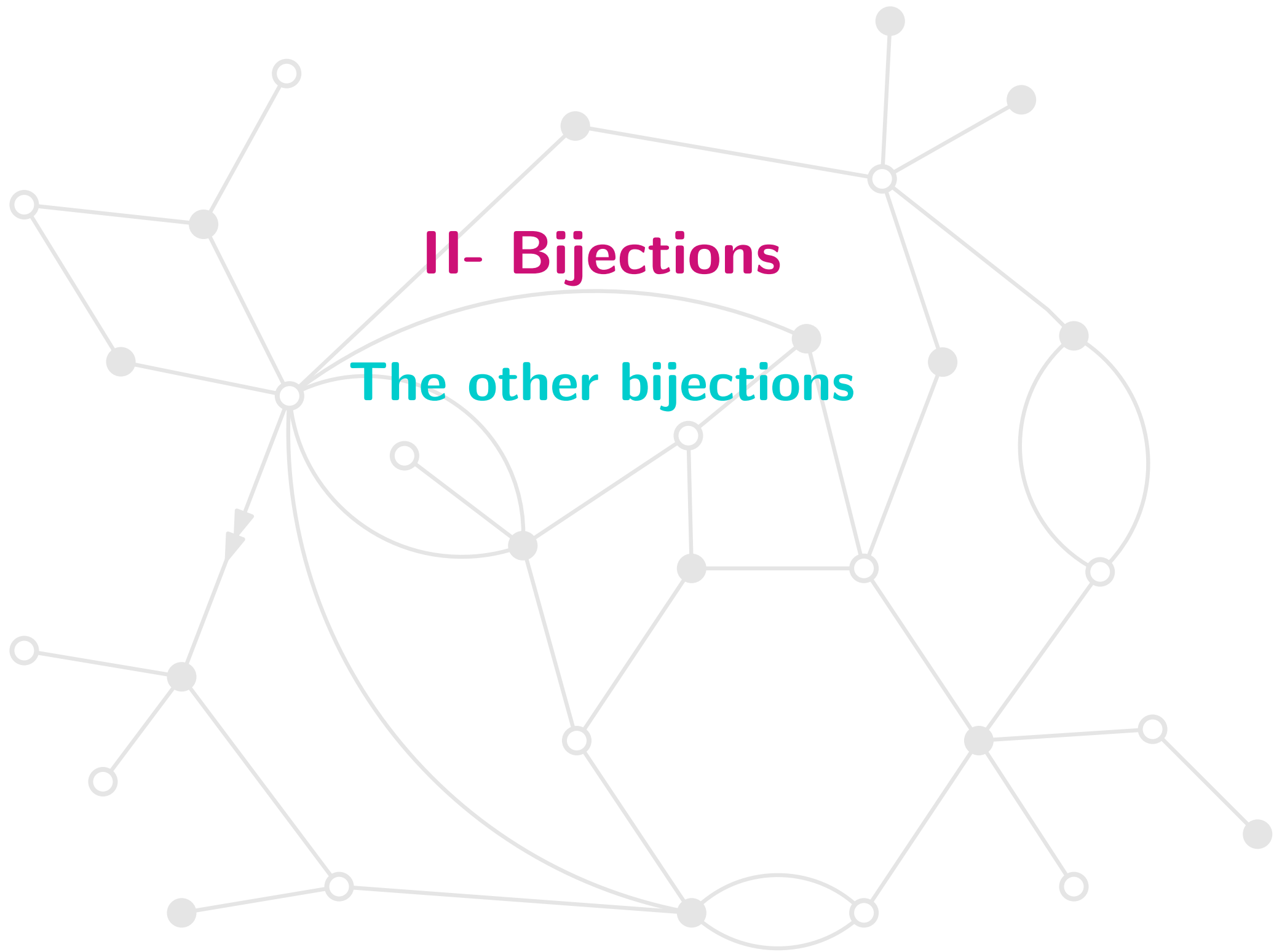
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$(M, -)$



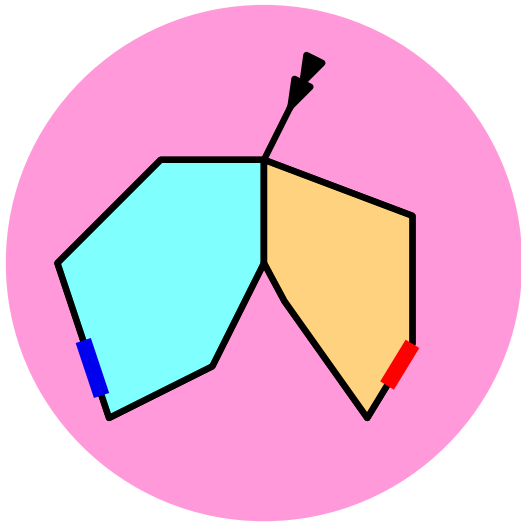


The other bijection

$$(f(\mathbf{d}) - 1)(f(\mathbf{d}) - 2)B(\mathbf{d}) = \sum_{\mathbf{s} + \mathbf{t} = \mathbf{d}} (f(\mathbf{s}) - 1)v(\mathbf{t})(v(\mathbf{t}) - 1)B(\mathbf{s})B(\mathbf{t})$$

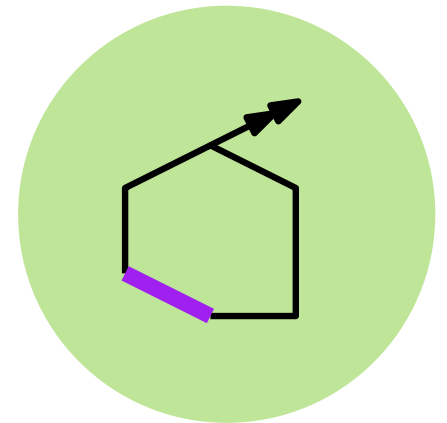
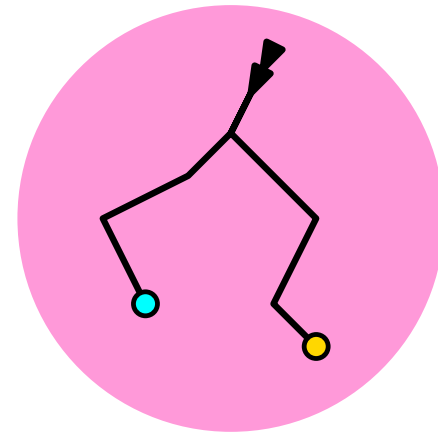
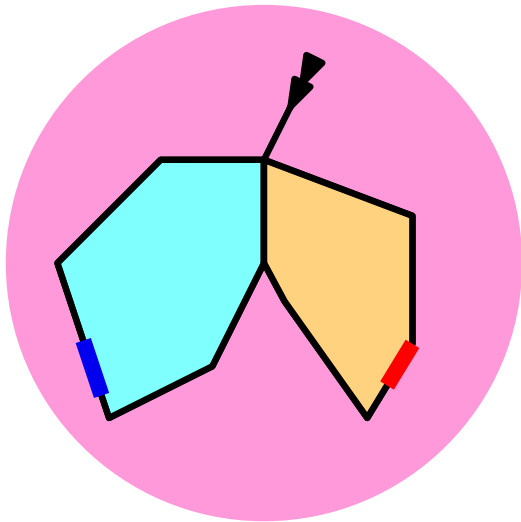
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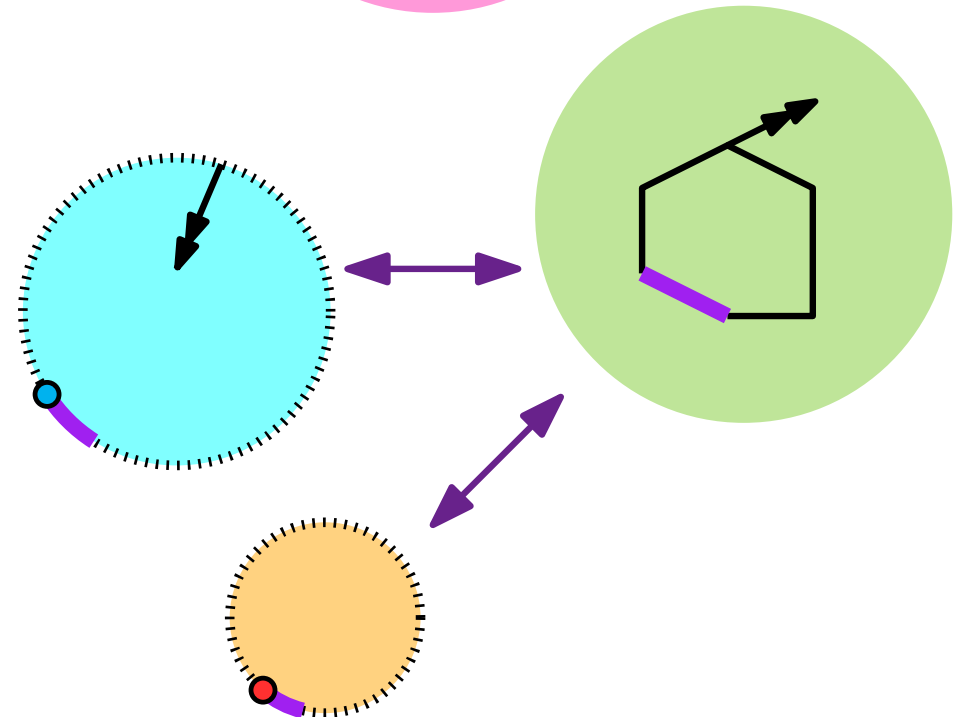
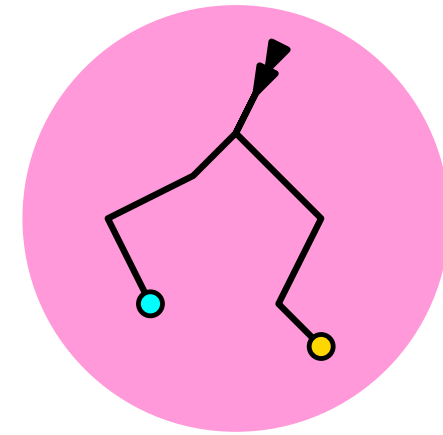
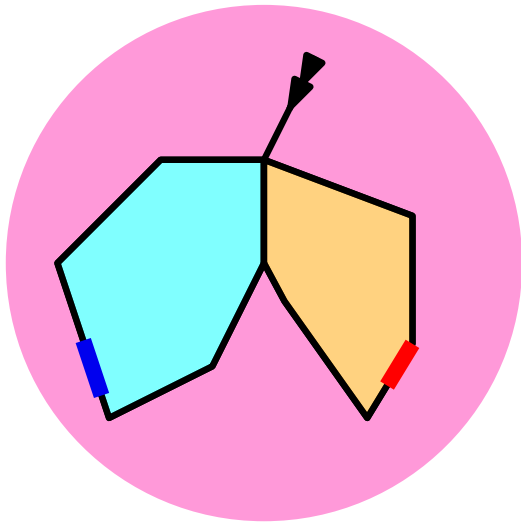
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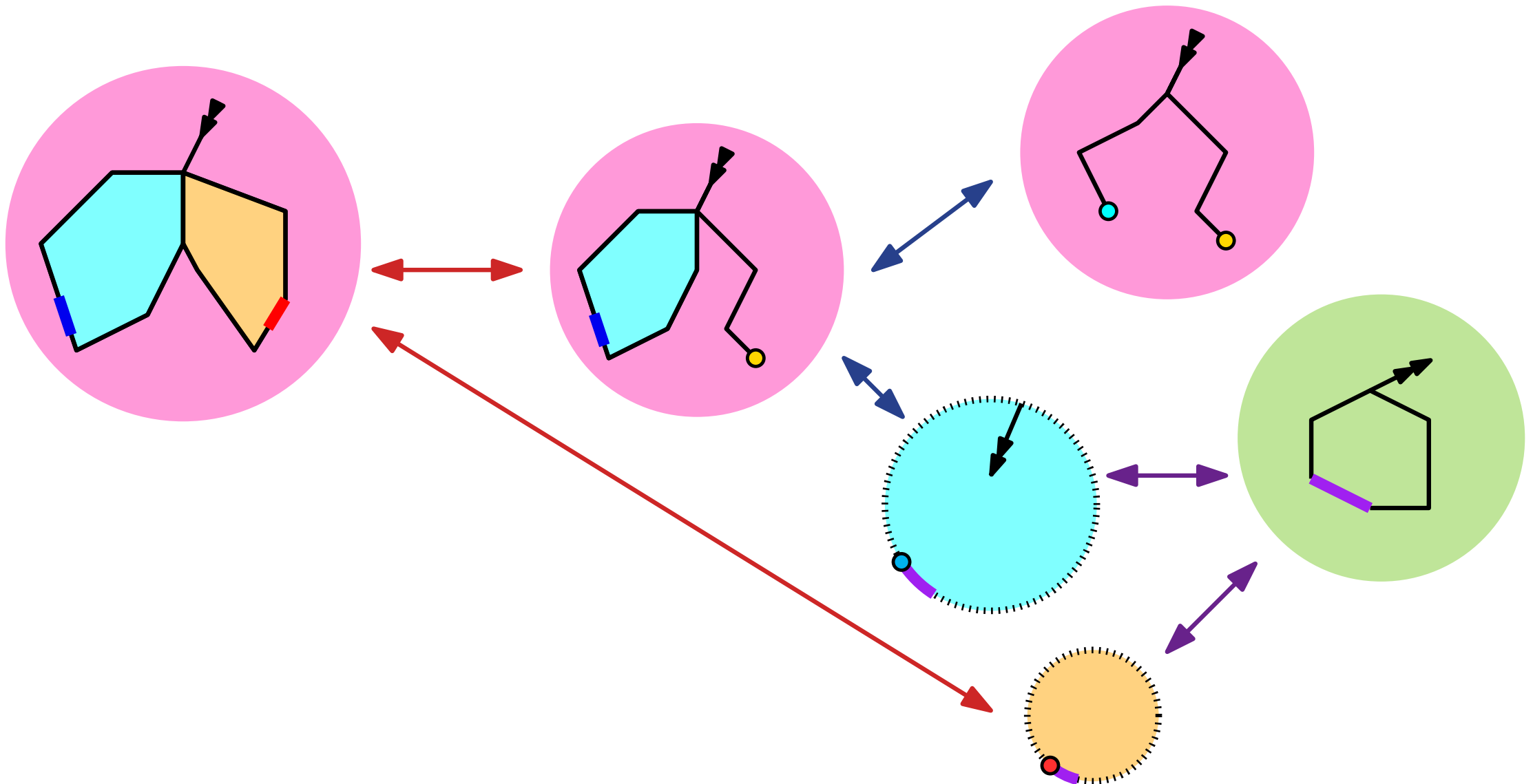
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Perspectives

- Unicellular case of Louf's formula:

$$\binom{n+1}{2} B_g(n) = \sum_{g^* \geq 0} \binom{v+2g^*}{2g^*+2} B_{g-g^*}(n).$$

Refinement:

$$\binom{n+1}{2} B_{v_\circ, v_\bullet}(n) = \sum_{\substack{w^*, b^* \geq 0 \\ 0 < w^* + b^* \leq 2g+2 \\ w^* \equiv b^* \pmod{2}}} \binom{v_\circ - 1 + w^*}{w^*} \binom{v_\bullet - 1 + b^*}{b^*} B_{v_\circ - 1 + w^*, v_\bullet - 1 + b^*}(n).$$

- Tutte's formula for q -colored triangulations:

$$\begin{aligned} (n+3)(n+2)T(n+1) &= 2(q-4)(3n+1)T(n) \\ &\quad + 2 \sum_{k=0}^n (3k+1)T(k)(n-k+2)(n-k+1)T(n-k). \end{aligned}$$

- General case of formulas from the KP hierarchy (Goulden-Jackson, Carrell-Chapuy).

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Thank you!

A third equation

By refining Louf's formula to count black and white vertices, we get:

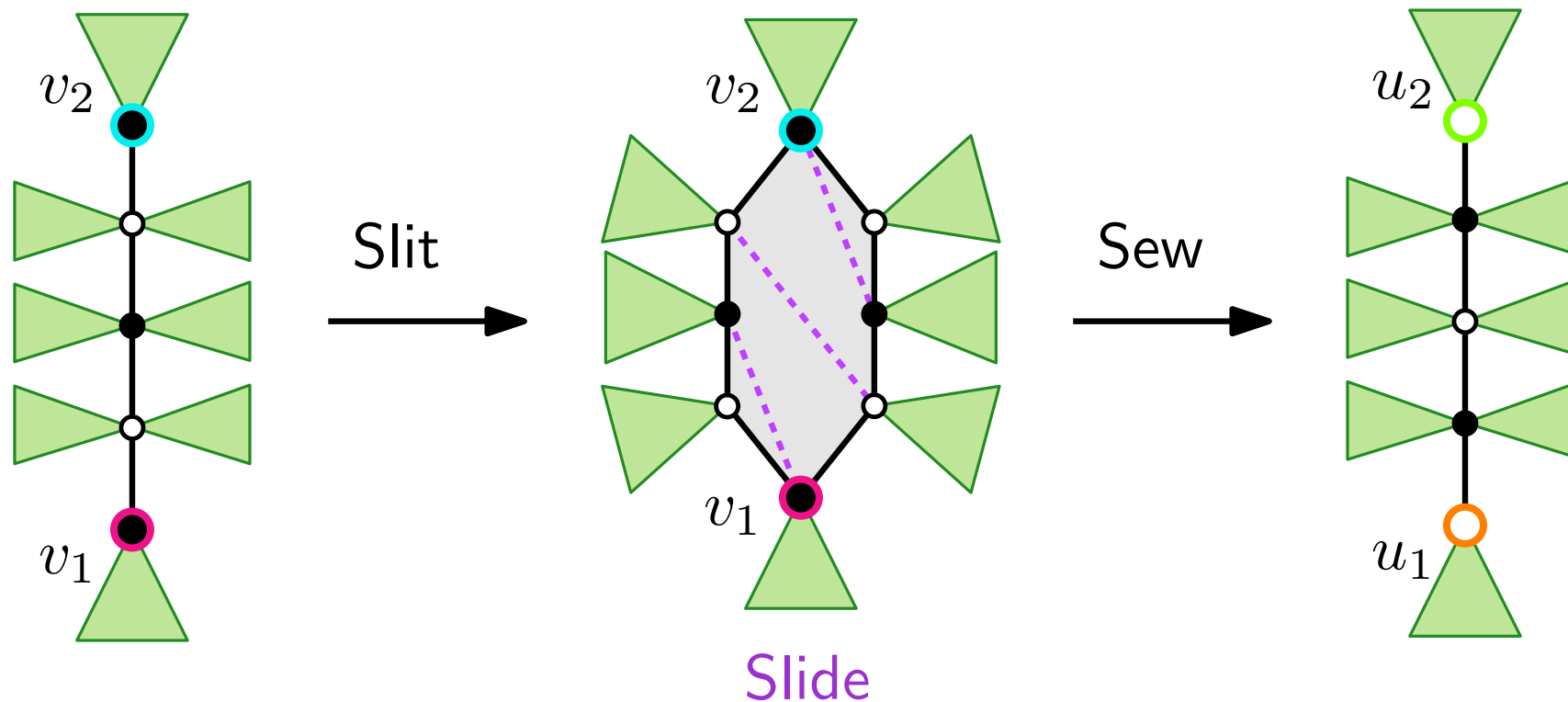
$$\begin{aligned} \binom{n+1}{2} B_{v_{\circ}, v_{\bullet}}(\mathbf{f}) &= \sum_{\substack{\mathbf{s}+\mathbf{t}=\mathbf{f} \\ w_1+w_2-w^*=v_{\circ} \\ b_1+b_2-b^*=v_{\bullet} \\ w^*+b^*=2}} (1+n_1) \binom{w_2}{w^*} \binom{b_2}{b^*} B_{w_1, b_1}(\mathbf{s}) B_{w_2, b_2}(\mathbf{t}) \\ &+ \sum_{\substack{w^*, b^* \geq 0 \\ w^*+b^*=2}} \binom{v_{\circ}-1+w^*}{w^*} \binom{v_{\bullet}-1+b^*}{b^*} B_{v_{\circ}-1+w^*, v_{\bullet}-1+b^*}(\mathbf{f}) \end{aligned}$$

So the $\binom{v}{2} B(\mathbf{f}) = \binom{v}{2} B(\mathbf{f})$ becomes

$$\begin{aligned} \binom{v_{\circ} + v_{\bullet}}{2} B_{v_{\circ}, v_{\bullet}}(\mathbf{f}) &= v_{\circ} v_{\bullet} B_{v_{\circ}, v_{\bullet}}(\mathbf{f}) + \binom{v_{\circ} + 1}{2} B_{v_{\circ}+1, v_{\bullet}-1}(\mathbf{f}) \\ &+ \binom{v_{\bullet} + 1}{2} B_{v_{\circ}-1, v_{\bullet}+1}(\mathbf{f}). \end{aligned}$$

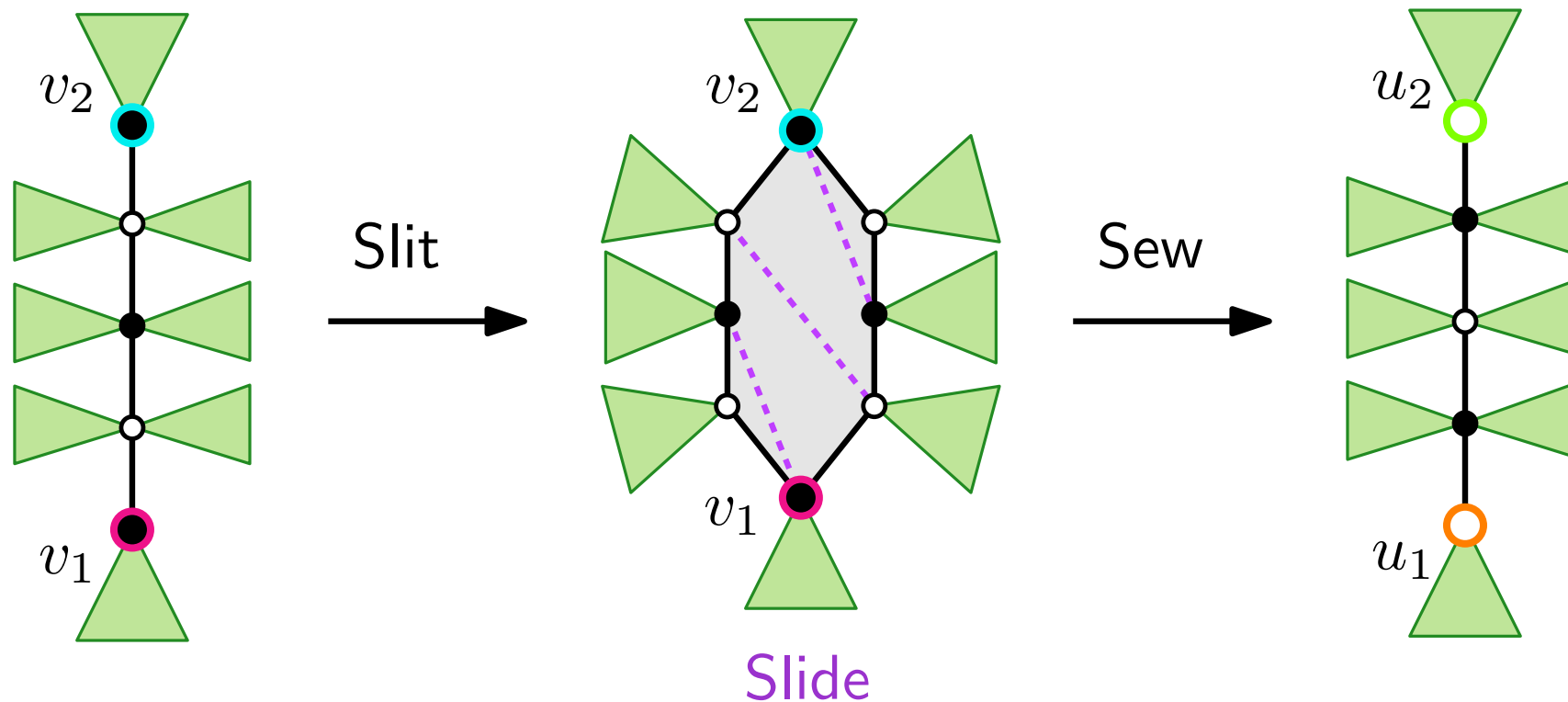
Bijection on trees

$$\binom{v_{\circ} + v_{\bullet}}{2} T_{v_{\circ}, v_{\bullet}} = v_{\circ} v_{\bullet} T_{v_{\circ}, v_{\bullet}}(\mathbf{f}) + \binom{v_{\circ}}{2} T_{v_{\circ}-1, v_{\bullet}+1} + \binom{v_{\bullet}}{2} T_{v_{\circ}+1, v_{\bullet}-1}.$$



Bijection on trees

$$\binom{v_{\circ} + v_{\bullet}}{2} T_{v_{\circ}, v_{\bullet}} = v_{\circ} v_{\bullet} T_{v_{\circ}, v_{\bullet}}(\mathbf{f}) + \binom{v_{\circ}}{2} T_{v_{\circ}-1, v_{\bullet}+1} + \binom{v_{\bullet}}{2} T_{v_{\circ}+1, v_{\bullet}-1}.$$



↪ Can this be generalized to bipartite maps ?