

Isogenies and endomorphism rings of elliptic curves

ECC Summer School

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Notations

- We fix a perfect field k . Since our aim is cryptographic applications of elliptic curves, most of the time k will be a finite field.
- An elliptic curve E is a smooth complete curve of genus 1 with a base point O_E . This base point uniquely determine a structure of algebraic group on E .
- If k is a finite field, every smooth complete curve of genus 1 has a rational point, so is an elliptic curve.
- An elliptic curve E/\mathbb{F}_q over a finite field of characteristic p is said to be supersingular if $\#E[p] = \{0\}$. In this case $\#E[p^n] = \{0\}$ for all n . Otherwise, $\#E[p^n] = p^n$ for all n , and E is said to be ordinary.

Complex elliptic curve

- Over \mathbb{C} : an elliptic curve is a torus $E = \mathbb{C}/\Lambda$, where Λ is a lattice $\Lambda = \mathbb{Z} + \tau\mathbb{Z}$, ($\tau \in \mathfrak{H}$).
- Let $\wp(z, \Lambda) = \sum_{w \in \Lambda \setminus \{0_E\}} \frac{1}{(z-w)^2} - \frac{1}{w^2}$ be the Weierstrass \wp -function and $E_{2k}(\Lambda) = \sum_{w \in \Lambda \setminus \{0_E\}} \frac{1}{w^{2k}}$ be the Eisenstein series of weight $2k$.
- Then $\mathbb{C}/\Lambda \rightarrow E, z \mapsto (\wp'(z, \Lambda), \wp(z, \Lambda))$ is an analytic isomorphism to the elliptic curve

$$y^2 = 4x^3 - 60E_4(\Lambda)x - 140E_6(\Lambda).$$

Isogenies between elliptic curves

Definition

An isogeny is a (non trivial) algebraic map $f : E_1 \rightarrow E_2$ between two elliptic curves such that $f(P + Q) = f(P) + f(Q)$ for all geometric points $P, Q \in E_1$.

Example

- If E is an elliptic curve, the multiplication by $[m]$ is an isogeny.
- If $E : y^2 = x^3 + ax + b$ is an elliptic curve defined over a finite field \mathbb{F}_q of characteristic p , the Frobenius $E \rightarrow E^{(p)}, (x, y) \mapsto (x^p, y^p)$ is an isogeny.
- Let E be the elliptic curve $y^2 = x^3 + x$ over \mathbb{F}_{17} . Let f be the map $f(x, y) = (x, 4y)$. Is f an isogeny?

Remark

Isogenies are surjectives. In particular, if E is ordinary, any isogenous curve to E is also ordinary.

Isogenies and algebraic maps

Theorem

An algebraic map $f : E_1 \rightarrow E_2$ is an isogeny if and only if $f(0_{E_1}) = f(0_{E_2})$

Proof.

Over \mathbb{C} : a bit of work on analytic functions. □

Corollary

An algebraic map between two elliptic curves is either

- *trivial (i.e. constant)*
- *or the composition of a translation with an isogeny.*

Equivalent isogenies

- Two isogenies $f_1 : E_1 \rightarrow E_2$ and $f_2 : E'_1 \rightarrow E'_2$ are equivalent if the following diagram commutes:

$$\begin{array}{ccc} E_1 & \xrightarrow{f_1} & E_2 \\ \downarrow \wr & & \downarrow \wr \\ E'_1 & \xrightarrow{f_2} & E'_2 \end{array}$$

- Let $E_1 : y^2 = x^3 + 4x + 2$ and $E_2 : y^2 = x^3 + 8x + 7$ be two elliptic curves over \mathbb{F}_{17} .
- Let $f_1 : E_1 \rightarrow E_2$ be the isogeny given by

$$\left(\frac{x^9 - x^8 + 8x^7 - 2x^6 - 6x^5 + 5x^4 + x^3 - 4x^2 + 2}{x^8 - x^7 + 2x^6 - 5x^5 + 7x^4 + 4x^3 - 8x^2 + 3x - 2}, \frac{x^{12}y + 7x^{11}y + 8x^{10}y - 2x^9y + 6x^8y + 5x^7y + 8x^6y + 2x^5y + 7x^4y - 6x^3y - 7x^2y + 5xy + 4y}{x^{12} + 7x^{11} - 3x^{10} + 7x^9 - 2x^8 + 2x^7 - 4x^6 - 6x^5 - 8x^4 - 5x^3 + 3x^2 + 6x + 3} \right)$$

Let $f_2 : E_1 \rightarrow E_2$ be the isogeny given by

$$\left(\frac{x^9 + 3x^7 - 5x^6 + 4x^5 - 5x^4 - 3x^3 + 6x^2 - 2x + 6}{-8x^8 + 8x^6 + 8x^5 + 4x^4 - 4x^3 - 5x^2 - 3x + 1}, \frac{x^{12}y + 3x^{10}y - 2x^9y - 5x^8y - 8x^7y - 4x^6y - x^5y - 7x^4y + x^3y - 6x^2y - 2xy - 6y}{-7x^{12} + 2x^{10} + 2x^9 - 8x^8 - 2x^7 - 8x^6 - x^5 - 5x^4 + 8x^3 - 2x^2 + 4x + 1} \right)$$

- Is f_1 equivalent to f_2 ?

Equivalent isogenies

- f_1 and f_2 have the same degrees. But $E_1 \neq E_2$!
- But they have the same j -invariant ($j = 4$), so they are isomorphic.
- We could compose f_2 with an isomorphism $E_2 \xrightarrow{\sim} E_1$ and test if it is equal to f_1 . But even if the curves were equal, we could still compose with automorphisms.
- So we have to construct “canonical” isogenies from f_1 and f_2 .
- Easier way: compute the kernels!

$$\ker f_1 = x^4 + 8x^2 + 8x + 6$$

$$\ker f_2 = x^4 + 8x^3 + 3x^2 + 16x + 7$$

- The kernel are different, hence the isogenies are not the same. (Since $\text{Aut}(E_1) = \{\pm 1\}$).
- Exercise: prove that f_1 is equivalent to the multiplication by 3.

Isogenies and kernels

Definition (Kernel)

The kernel $\ker f$ of an isogeny $f : E_1 \rightarrow E_2$ is the set of geometric points $P \in E_1$ such that $f(P) = 0_{E_2}$.

Definition (Degree)

The degree of an isogeny f is the degree of the extension field $[k(E_1) : f^*k(E_2)]$. An isogeny is separable iff $\#\ker f = \deg f$.

- The Frobenius is an inseparable isogeny of degree p .
- Every isogeny is the composition of a separable isogeny with a power of the Frobenius \Rightarrow from now on we only focus on separable isogenies.

Theorem

There is a bijection between separable isogenies and finite subgroups of E :

$$\begin{aligned} (f : E_1 \rightarrow E_2) &\mapsto \ker f \\ (E_1 \rightarrow E_1/G) &\leftarrow G \end{aligned}$$

Isogenies and multiplications

- If $H \subset G$ are finite subgroups of E , then the isogeny $E \rightarrow E/G$ splits as $E \rightarrow E/H \rightarrow (E/H)/(G/H)$.
- In particular, for every (separable) isogeny $f : E \rightarrow E'$, there exists a contragredient isogeny $f' : E' \rightarrow E$ such that $f' \circ f = [m]$, where m is the exponent of $\ker f$.
- We can also identify f' as the dual isogeny \hat{f} of f (if $m = \deg f$):

$$\begin{array}{ccccccc}
 0 & \longrightarrow & K & \longrightarrow & E & \xrightarrow{f} & E' & \longrightarrow & 0 \\
 & & & & \downarrow \wr & & \downarrow \wr & & \\
 & & & & \hat{E} & \xleftarrow{\hat{f}} & E' & \xleftarrow{\quad} & \hat{K} & \xleftarrow{\quad} & 0
 \end{array}$$

Algorithms for manipulating isogenies

- 1 Given a finite subgroup $G \subset E$, construct the isogeny E/G .
- 2 Given E_1 and E_2 , test if they are isogenous. If so construct an (or all) isogenies $E_1 \rightarrow E_2$.
- 3 Given E and ℓ , find ℓ -isogenous curves to E (and iterate to construct the isogeny graph).
- 4 Find cyclic rational subgroups of E (by using the correspondance between isogenies and kernels).

Remark

Algorithm 4 can be obtained by combining algorithms 2 and 3: first compute all ℓ -isogenous curves E' , and from them compute the isogeny $E \rightarrow E'$ of degree ℓ , whose kernel give a cyclic subgroup of $E[\ell]$.

Destructive cryptographic applications

- An isogeny $f : E_1 \rightarrow E_2$ transports the DLP problem from E_1 to E_2 . This can be used to attack the DLP on E_1 if there is a weak curve on its isogeny class (and an efficient way to compute an isogeny to it).

Example

- extend attacks using Weil descent [[GHS02](#)] (remember Vanessa's talk!)
- Transfer the DLP from the Jacobian of a hyperelliptic curve of genus 3 to the Jacobian of a quartic curve [[Smi09](#)].

Constructive cryptographic applications

- One can recover informations on the elliptic curve E modulo ℓ by working over the ℓ -torsion.
- But by computing isogenies, one can work over a cyclic subgroup of cardinal ℓ instead.
- Since thus a subgroup is of degree ℓ , whereas the full ℓ -torsion is of degree ℓ^2 , we can work faster over it.

Example

- The SEA point counting algorithm [Sch95; Mor95; Elk97] (go to François' talk for more details).
- The CRT algorithms to compute class polynomials [Sut09; ES10].
- The CRT algorithms to compute modular polynomials [BLS09].

Further applications of isogenies

- Splitting the multiplication using isogenies can improve the arithmetic (remember Laurent's talk) [DIK06; Gau07].
- The isogeny graph of a supersingular elliptic curve can be used to construct secure hash functions [CLG09].
- Construct public key cryptosystems by hiding vulnerable curves by an isogeny (the trapdoor) [Tes06], or by encoding informations in the isogeny graph [RS06].
- Take isogenies to reduce the impact of side channel attacks [Sma03].
- Construct a normal basis of a finite field [CL09].
- Improve the discrete logarithm in \mathbb{F}_q^* by finding a smoothness basis invariant by automorphisms [CL08].

Class of isomorphisms of elliptic curves

- Every elliptic curve has a Weierstrass equation:

$$y^2 + a_1xy + a_3y = x^3 + a_2x^2 + a_4x + a_6 \quad (1)$$

with the discriminant $\Delta_E = -b_2b_8 - 8b_3 - 27b_2 + 9b_2b_4b_6 \neq 0$.

(Here $b_2 = a_1^2 + 4a_2$, $b_4 = 2a_4 + a_1a_3$, $b_6 = a_3^2 + 4a_6$,

$b_8 = a_1^2a_6 + 4a_2a_6 - a_1a_3a_4 + a_2a_3^2 - a_4^2$).

- The j -invariant of E is

$$j_E = \frac{(b_2^2 - 24b_4)^3}{\Delta_E}$$

Theorem

Two elliptic curves E and E' are isomorphic over \bar{k} if and only if $j_E = j_{E'}$.

The case of a finite field of characteristic $p > 3$

- We can always write the Weierstrass equation as

$$y^2 = x^3 + ax + b.$$

- The discriminant is $-16(4a^3 + 27b^2)$.
- The j -invariant is

$$j_E = 1728 \frac{4a^3}{4a^3 + 27b^2}.$$

Isomorphisms

- The isomorphisms (over \bar{k}) of isomorphisms of elliptic curves in Weierstrass form are given by the maps

$$(x, y) \mapsto (u^2x + r, u^3y + u^2sx + t)$$

for $u, r, s, t \in \bar{k}$, $u \neq 0$.

- If we restrict to elliptic curves of the form $y^2 = x^3 + ax + b$ then $s = t = 0$.

Proposition

Let E/\mathbb{F}_q and E'/\mathbb{F}_q be two *ordinary* elliptic curves such that $j_E = j_{E'}$. Then

$$E \simeq E' \text{ over } \mathbb{F}_q$$

$$\Leftrightarrow E \text{ and } E' \text{ are isogenous over } \mathbb{F}_q$$

$$\Leftrightarrow \#E = \#E'.$$

Twists

- A twist of an elliptic curve E/\mathbb{F}_q is an elliptic curve E'/\mathbb{F}_q isomorphic to E over $\overline{\mathbb{F}}_q$ but not over \mathbb{F}_q .
- Every elliptic curve $E : y^2 = x^3 + ax + b$ has a quadratic twist

$$E' : \delta y^2 = x^3 + ax + b$$

for any non square $\delta \in \mathbb{F}_q$. E and E' are isomorphic over \mathbb{F}_q^2 .

- If E/\mathbb{F}_q is an **ordinary** elliptic curve with $j_E \notin \{0, 1728\}$ then the only twist of E is the quadratic twist. If $j_E = 1728$, then E admits 4 twists. If $j_E = 0$, then E admits 6 twists.

When are two elliptic curves isogenous?

Theorem (Tate)

Two elliptic curves over \mathbb{F}_q are isogenous if and only if they have the same cardinal.

Proof.

- If E and E' are isogenous, they have the same cardinal: use the dual isogeny and look at the action of the Frobenius on $E[\ell]$ for ℓ not dividing the degree of the isogeny.
- The reciprocal is a theorem of Tate.



Isogenies between two elliptic curves

In this slide, E_1/\mathbb{F}_q and E_2/\mathbb{F}_q are **ordinary** elliptic curves over \mathbb{F}_q .

- If E_1 and E_2 are isogenous, then any isogeny over $\overline{\mathbb{F}}_q$ is in fact \mathbb{F}_q -rational.
- If $f : E_1 \rightarrow E_2$ is an isogeny over $\overline{\mathbb{F}}_q$ of prime degree, then there exist twists E'_1 and E'_2 of E_1 and E_2 such that f descends to an \mathbb{F}_q -rational isogeny $f : E'_1 \rightarrow E'_2$.
- Either $\text{Hom}_{\mathbb{F}_q}(E_1, E_2) = \{0\}$ or $\text{Hom}_{\mathbb{F}_q}(E_1, E_2)$ is a free \mathbb{Z} -module of rank 2.

Computing explicit isogenies

- If E_1 and E_2 are two elliptic curves given by Weierstrass equations, a morphism of curve $f : E_1 \rightarrow E_2$ is of the form

$$f(x, y) = (R_1(x, y), R_2(x, y))$$

where R_1 and R_2 are rational functions, whose degree in y is less than 2 (using the equation of the curve E_1).

- If f is an isogeny, $f(-P) = -f(P)$. If $\text{char } k > 3$ so we can assume that E_1 and E_2 are given by reduced Weierstrass forms, this mean that R_1 depends only on x , and R_2 is y time a rational function depending only on x .
- Let $w_E = dx/2y$ be the canonical differential. Then $f^*w_{E'} = cw_E$, with c in k .
- This show that f is of the form

$$f(x, y) = \left(\frac{g(x)}{h(x)}, cy \left(\frac{g(x)}{h(x)} \right)' \right).$$

$h(x)$ give (the x coordinates of the points in) the kernel of f (if we take it prime to g).

- If $c = 1$, we say that f is normalized.

Isogeny from the kernel

Remark

Every isogeny is a composition of a multiplication by $[m]$ and an isogeny with cyclic kernel (we could even further reduce to a composition with cyclic kernels of prime orders).

- Let E/k be an elliptic curve. Let $G = \langle P \rangle$ be a rational finite subgroup of E . We want to construct the isogeny $E \rightarrow E/G$.
- We need to find the Weierstrass coordinates X, Y on $k(E/G)$. But $k(E/G) = k(E)^G$ are the rational functions on E invariants under translation by a point of G .
- Moreover the Weierstrass coordinates x and y on E are characterized (up to isomorphism) by

$$\begin{aligned}
 v_{0_E}(x) &= -2 & v_P(x) &\geq 0 & \text{if } P \neq 0_E \\
 v_{0_E}(y) &= -3 & v_P(y) &\geq 0 & \text{if } P \neq 0_E \\
 y^2/x^3(0_E) &= 1
 \end{aligned}$$

Vélu's formula

- Vélu constructs the isogeny $E \rightarrow E/G$ as

$$X(P) = x(P) + \sum_{Q \in G \setminus \{0_E\}} (x(P+Q) - x(Q))$$

$$Y(P) = y(P) + \sum_{Q \in G \setminus \{0_E\}} (y(P+Q) - y(Q)).$$

The choices are made so that the formulas give a normalized isogeny.

- Moreover by looking at the expression of X and Y in the formal group of E , Vélu recovers the equations for E/G .
- For instance if $E : y^2 = x^3 + ax + b = f(x)$ then E/G is

$$y^2 = x^3 + (a - 5t)x + b - 7\omega$$

where $t = \sum_{Q \in G \setminus \{0_E\}} f'(Q)$, $u = 2 \sum_{Q \in G \setminus \{0_E\}} f(Q)$ and $\omega = \sum_{Q \in G \setminus \{0_E\}} x(Q)f'(Q)$.

Complexity of Vélu's formula

- Even if G is rational, the points in G may live to an extension of degree up to $\#G - 1$.
- Thus summing over the points in the kernel G can be expensive.
- Let $h(x) = \prod_{Q \in G \setminus \{O_E\}} (x - x(Q))$. The symmetry of X and Y allows us to express everything in term of h .
- For instance if E is given by a reduced Weierstrass equation $y^2 = f(x)$, we have

$$f(x, y) = \left(\frac{g(x)}{h(x)}, y \left(\frac{g(x)}{h(x)} \right)' \right), \text{ with}$$

$$\frac{g(x)}{h(x)} = \#G \cdot x - \sigma - f'(x) \frac{h'(x)}{h(x)} - 2f(x) \left(\frac{h'(x)}{h(x)} \right)',$$

where σ is the first power sum of h (i.e. the sum of the x -coordinates of the points in the kernel).

- When $\#G$ is odd, $h(x)$ is a square, so we can replace it by its square root.
- The complexity of computing the isogeny is then $O(M(\#G))$ operations in k .

Computing isogenous curves from E

- Let E be an elliptic curve and ℓ a prime number. We want to compute all ℓ -isogenous elliptic curves to E .
- Easy! Compute the rational cyclic subgroups of $E[\ell]$ and apply Vélu's formulas. These subgroups can be obtained as factors of the ℓ -division polynomial $\prod_{Q \in E[\ell] \setminus \{0_E\}} (x - x(Q))$.
- But the division polynomial has degree $(\ell^2 - 1)/2$ (if ℓ odd), and factorizing it will cost $O(\ell^{3.63})$. We only want to compute isogenies of degree ℓ . Can we do better?

Modular polynomials

Here $k = \bar{k}$.

Definition (Modular polynomial)

The modular polynomial $\varphi_\ell(x, y) \in \mathbb{Z}[x, y]$ is a bivariate polynomial such that $\varphi_\ell(x, y) = 0 \Leftrightarrow x = j(E)$ and $y = j(E')$ with E and E' ℓ -isogeneous.

- Roots of $\varphi_\ell(j(E), \cdot) \Leftrightarrow$ elliptic curves ℓ -isogeneous to E .
There are $\ell + 1 = \#\mathbb{P}^1(\mathbb{F}_\ell)$ such roots if ℓ is prime.
- φ_ℓ is symmetric.
- The height of φ_ℓ grows as $O(\ell)$.

Rational roots of the modular polynomials

Theorem

- Let E/\mathbb{F}_q be an *ordinary* elliptic curve with j -invariant not equal to 0 or 1728.
- Let ℓ be prime and j' be a root of $\varphi_\ell(j_E, \cdot)$ over \mathbb{F}_{q^n} .
- Then j' corresponds to a \mathbb{F}_{q^n} -rational ℓ -isogeny $E \rightarrow E'$.

Proof.

There exist a $\overline{\mathbb{F}_q}$ -isogeny between E and E' so a \mathbb{F}_{q^n} -isogeny on twists of E and E' . But with the hypothesis, the only twist of E is the quadratic one, so by applying a quadratic twist to the isogeny, we find a \mathbb{F}_{q^n} -rational isogeny starting from E . \square

Corollary

We can use the modular polynomial φ_ℓ to construct ℓ -isogeny graphs!

Computing the modular polynomial

- 1 The complex analytic method: if we see $\tau \mapsto j(\tau)$ and $\tau \mapsto j(\tau/\ell)$ as a modular functions on \mathfrak{H} ; then $\varphi_\ell(\cdot, j)$ is the minimal polynomial of $j(\cdot/\ell)$ in $\mathbb{C}(j)$. One can then recover the polynomial by computing the Fourier coefficients of j and $j(\cdot/\ell)$ with high precision.
- 2 The CRT method: use Vélu's formulas to compute $\varphi_\ell \bmod p$ for small p and the CRT to recover the full modular polynomial.

Remark

- *Using asymptotically fast algorithms, both algorithms are quasilinear in the size ℓ^3 of φ_ℓ , so the computations are memory bounded. But the CRT algorithm allow to compute the specialization $\varphi_\ell(j, \cdot) \in \mathbb{F}_p[x]$ directly and is the faster in practice.*
- *To reduce the size of the coefficients, one use a different modular function in $X_0^*(\ell)$ than $j(\tau/\ell)$.*

Finding an isogeny between two isogenous elliptic curves

- Let E and E' be ℓ -isogenous abelian varieties (we can check that $\varphi_\ell(j_E, j_{E'}) = 0$). We want to compute the isogeny $f : E \rightarrow E'$.
- The explicit forms of isogenies are given by Vélu's formula, which give normalized isogenies. We first need to normalize E' .
- Over \mathbb{C} , the equation of the normalized curve E' is given by the Eisenstein series $E_4(\ell\tau)$ and $E_6(\ell\tau)$. We have $j'(\ell\tau)/j(\ell\tau) = -E_6(\tau)/E_4(\tau)$. By differencing the modular polynomial, we recover the differential logarithms.
- We obtain that from $E : y^2 = x^3 + ax + b$, a normalized model of $j_{E'}$ is given by the Weierstrass equation

$$y^2 = x^3 + Ax + B$$

$$\text{where } A = -\frac{1}{48} \frac{J^2}{j_{E'}(j_{E'} - 1728)}, B = -\frac{1}{864} \frac{J^3}{j_{E'}^2(j_{E'} - 1728)} \text{ and } J = -\frac{18}{\ell} \frac{b}{a} \frac{\varphi'_\ell(x)(j_E, j_{E'})}{\varphi'_\ell(y)(j_E, j_{E'})} j_E.$$

Remark

$E_2(\tau)$ is the differential logarithm of the discriminant. Similar methods allow to recover $E_2(\ell\tau)$, and from it $\sigma = \sum_{P \in K \setminus \{0_E\}} x(K)$.

Finding the isogeny between the normalized models (I: Stark's method)

- We need to find the rational function $I(x) = g(x)/h(x)$ giving the isogeny $f : (x, y) \mapsto (I(x), yI'(x))$ between E and E' .
- Over \mathbb{C} the coordinates of the elliptic curve are given by the elliptic functions: $x = \wp(z)$ and $y = \wp'(z)$.
- We have to find I such that $\wp_{E'}(z) = I \circ \wp_E(z)$.
- Stark's idea is to develop $\wp_{E'}$ as a continuous fraction in \wp_E , and approximate I as p_n/q_n .
- This algorithm is quasi-quadratic ($\tilde{O}(\ell^2)$).

Finding the isogeny between the normalized models (II: Elkies' method)

- We need to find the rational function $I(x) = g(x)/h(x)$ giving the isogeny $f : (x, y) \mapsto (I(x), yI'(x))$ between E and E' .
- Plugging f into the equation of E' shows that I satisfy the differential equation

$$(x^3 + ax + b)I'(x)^2 = I(x)^3 + AI(x) + B.$$

- Using an asymptotically fast algorithm to solve this equation yields $I(x)$ in time quasi-linear ($\tilde{O}(\ell)$).
- Knowing σ gains a logarithmic factor.

Finding an isogeny between two isogenous elliptic curves (the case of small characteristic)

- The preceding algorithm needs $p > 8\ell - 5$ to solve the differential equation.
- Idea in small characteristic: lift the curves to \mathbb{Q}_q by taking lifts \tilde{j}_E and $\tilde{j}_{E'}$ such that $\varphi_\ell(\tilde{j}_E, \tilde{j}_{E'}) = 0$ and apply the preceding algorithm.
- Even if E' is normalized, we need the modular polynomial to lift E' and normalize the lift.

Finding an isogeny: total complexity

To summarize, we have the following algorithm to find an isogeny from E in large characteristic:

Algorithm ([BMS+08])

- 1 Compute φ_ℓ (cost $\tilde{O}(\ell^3)$)
- 2 Specialize on j_E to obtain $\varphi_\ell(X, j_E)$ (cost $\tilde{O}(\ell^2 \log q)$)
- 3 Find a root $j_{E'}$ of $\varphi_\ell(X, j_E)$ to obtain the j -invariant of a ℓ -isogenous curve E' (cost $\tilde{O}(\ell \log^2 q)$).
- 4 Compute the normalized model for E' (cost $\tilde{O}(\ell^2 \log q)$).
- 5 Solve the differential equation (cost $\tilde{O}(\ell \log q)$).

Finding an isogeny: total complexity

With the adaptation in small characteristic still of total cost $\tilde{O}(\ell^3 + \ell \log^2 q)$:

Algorithm ([LS08])

- 1 Compute $\varphi_\ell(X, j_E)$ (cost $\tilde{O}(\ell^3 + \ell^2 \log q)$).
- 2 Lift j_E and find a root $\tilde{j}_{E'}$ in precision $O(1 + \log^2 \ell / \log q)$ (cost $\tilde{O}(\ell \log^2 q)$).
- 3 Compute the normalized model for \tilde{E}' (cost $\tilde{O}(\ell^2 \log q)$).
- 4 Solve the differential equation in \mathbb{Q}_q (cost $\tilde{O}(\ell \log q)$).
- 5 Reduce in \mathbb{F}_q (cost $\tilde{O}(\ell \log q)$).

Finding an isogeny between two isogenous elliptic curves (the case of small characteristic): Couveigne's algorithm

Another idea to compute the isogeny in the **ordinary** case comes from Couveigne:

Algorithm

- ① Find generators P and P' of the cyclic groups $E[p^\alpha]$ and $E'[p^\alpha]$ for $p^\alpha \ll \ell$.
 - ② Interpolate the algebraic map $f : E[p^\alpha] \rightarrow E'[p^\alpha], iP \mapsto iP'$.
 - ③ Test if f is an isogeny.
- [Cou94] works with formal groups.
 - [Cou96] use p -descent and towers of Artin-Schreier extensions. The best implementation [Feo10a] has complexity $\tilde{O}(\ell^2)$.
 - But the complexity is exponential in $\log(p)$.

Other algorithms to compute the isogeny

- Lercier for $p = 2$: solve the differential equation using linear algebra. Cost $\tilde{O}(\ell^3 \log q)$ operations, in practice the fastest for $p = 2$.
- Joux and Lercier: lift in \mathbb{Q}_q with precision $O(\ell)$. Cost $\tilde{O}(\ell^2(1 + \ell/p) \log q)$; useful for the intermediate case $p \approx \log q$.
- When the degree ℓ is not known but only bounded by L . The naive method is to apply one of the above algorithm for all $\ell \leq L$. This increase the cost by a degree 1 in L . However, Couveigne's algorithm can be adapted to stay in $\tilde{O}(L^2)$ [Feo10b].
- Subexponential algorithms for computing isogenies of large degree [JS10; CJS10].

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The characteristic polynomial of the Frobenius

From now on k will represent a finite field: $k = \mathbb{F}_q$.

- There exist a unique polynomial χ_π such that for every n prime to the characteristic p , $\chi_\pi \pmod n$ is the characteristic polynomial of the action of the Frobenius π on $E[n]$ (here $\pi = \text{Fr}_{\mathbb{F}_q}$).
- We have $\chi_\pi(\pi) = 0$, and $\#E = \chi_\pi(1)$.
- We have $\chi_\pi = X^2 - tX + q$ where the trace t is such that $|t| \leq 2\sqrt{q}$ (Hasse).

The endomorphism ring

Definition

- If E_1 and E_2 are elliptic curves, we note $\text{Hom}_k(E_1, E_2)$ the \mathbb{Z} -module of all k -morphisms from E_1 to E_2 . The **endomorphism ring** $\text{End}_k(E)$ is then $\text{End}_k(E) = \text{Hom}_k(E, E)$.
- We note $\text{End}_k^0(E) = \text{End}_k(E) \otimes_{\mathbb{Z}} \mathbb{Q}$ the endomorphism fraction ring.

Remark

- *Every non nul element of $\text{Hom}_k(E_1, E_2)$ is an isogeny (possibly non separable).*
- *$\text{End}_k^0(E_1)$ is a division algebra, and $\text{End}_k(E_1)$ is an order in it.*
- *If $\text{Hom}_k(E_1, E_2) \neq 0$, then $\text{End}_k^0(E_1) = \text{End}_k^0(E_2)$ and $\text{Hom}_k(E_1, E_2)$ is a free \mathbb{Z} -module of the same rank as $\text{End}_k(E_1)$.*
- *If \mathcal{E} is the isogeny class of E , $\text{End}_k^0(E)$ does not depend on the curve $E \in \mathcal{E}$.*
- *$\text{End}_k(E)$ is either commutative of rank 2, or an order of rank 4 in a quaternion algebra.*

The ordinary case

If E is **ordinary**, then

- χ_π is irreducible.
- $K = \text{End}_k^0(E)$ is a quadratic imaginary field.
- K is generated by π : $K = \mathbb{Q}(\pi)$.
- $\text{End}_k(E)$ is an order O in K .
- For any extension k' of k we have $\text{End}_k(E) = \text{End}_{k'}(E) = \text{End}_{\overline{k}}(E)$.

Remark

If k' is an extension of k of degree n , then the Frobenius of $E_{k'}$ seen in K is π^n .

From now on, we assume that E is **ordinary**, and we note $O = \text{End}_k(E)$ and K the quadratic imaginary field $\text{End}_k^0(E)$.

Automorphisms and twist

- The automorphisms of E are the invertible elements in $O = \text{End } E$.
- All invertible elements are roots of unity.
- We usually have $O^* = \{\pm 1\}$ except in the following exceptions:
 - ① $j_E = 1728$ ($p \neq 2, 3$), in this case O is the maximal order in $\mathbb{Q}(i)$ and $\#O^* = 4$;
 - ② $j_E = 0$ ($p \neq 2, 3$), in this case O is the maximal order in $\mathbb{Q}(i\sqrt{3})$ and $\#O^* = 6$;
 - ③ $j_E = 0$ ($p = 3$), in this case E is supersingular and $\#O^* = 12$;
 - ④ $j_E = 0$ ($p = 2$), in this case E is supersingular and $\#O^* = 24$.
- The Frobenius $\pi \in K$ characterizes the isogeny class of E (Tate). A twisted isogeny class will correspond to a Frobenius $\pi' \neq \pi$, where there exist n with $\pi^n = \pi'^n$. This gives a bijection between the twisted isogeny class and the roots of unity in K .
- More generally, there is a bijection between O^* and the twists of E .

Reduction and lifting (see Marco's talk)

- Let O be an order in an imaginary quadratic field K . Then there are h_O (the class number of O) elliptic curves over $\overline{\mathbb{Q}}$ with endomorphism ring O . They are defined over the ray class field H_O of O .
- If $p \nmid \Delta_O$, p is a prime of good reduction. Let \mathfrak{p} be a prime above p in H_O . If p is inert in K , $E_{\mathfrak{p}}$ is supersingular. If p splits, $E_{\mathfrak{p}}$ is ordinary, and its endomorphism ring is the minimal order containing O of index prime to p .
- Reciprocally, if E/\mathbb{F}_q is an ordinary elliptic curve, the couple $(E, \text{End}(E))$ can be lifted over \mathbb{Q}_q .

Corollary

- *If E/\mathbb{F}_q is an ordinary elliptic curve, then $\text{End}(E)$ is an order in $K = \mathbb{Q}(\pi)$ of conductor prime to p . For every order O of K such that $\mathbb{Z}[\pi] \subset O$, there exist an isogenous curve whose endomorphism ring is O .*
- *Reciprocally, for every order O of discriminant a non zero square modulo p ; let n be the order of one of the prime above p in the class group of O . Then there exist an (ordinary) elliptic curve E' over \mathbb{F}_{q^n} with $\text{End}(E') = O$.*

The structure of the rational points

Theorem (Lenstra)

Let E/\mathbb{F}_q be an ordinary elliptic curve. We have as $\text{End}_{\mathbb{F}_q}(E)$ -modules

$$E(\mathbb{F}_{q^n}) \simeq \frac{\text{End}_{\mathbb{F}_q}(E)}{\pi^n - 1}$$

Corollary

- Let $a, m \in \mathbb{Z}$ be such that $\mathcal{O}_K = \mathbb{Z}[\frac{\pi-a}{m}]$.
 - Let γ_E be the index of \mathcal{O} in \mathcal{O}_K .
 - Then $E(\mathbb{F}_q) = \mathbb{Z}/n_1\mathbb{Z} \oplus \mathbb{Z}/n_2\mathbb{Z}$ where $n_1 \mid n_2$ and $n_1 n_2 = \#E(\mathbb{F}_q)$.
 - Explicitly, we have: $n_1 = \gcd(a-1, m/\gamma_E)$.
-
- Exercise: show that $n_1 \mid q-1$ (use the Weil pairing).

Endomorphisms and isogenies

- Let $f : E_1 \rightarrow E_2$ be an isogeny of degree ℓ prime. Then either
 - ① f is an **ascending isogeny**: $O_1 \subset O_2$ with $[O_2 : O_1] = \ell$;
 - ② f is a **descending isogeny**: $O_2 \subset O_1$ with $[O_1 : O_2] = \ell$;
 - ③ f is an **horizontal isogeny**: $O_1 = O_2$.
- The horizontal case can only happen when O_1 is maximal locally in ℓ : $(O_1)_\ell = (O_K)_\ell$.
- Let $\ker f$ be the kernel of f . Let $O_f \subset O_1$ be the subring (of index ℓ) of isogenies fixing $\ker f$. Then f induce an injection $O_f \hookrightarrow O_2$.
- If $\psi \in O_1^*$ is an automorphism, then either ψ fixes $\ker f$ and descends to an automorphism of O_2 , or ψ induce an isogeny equivalent to f .

Isogeny graph: the local picture

- Let E be an ordinary elliptic curve with endomorphism ring O , and $\ell \neq p$ be a prime.
- We note Δ the discriminant of O_K , and $\Delta_\pi = t^2 - 4p$ the discriminant of χ_π .
- We have $\Delta_\pi = \gamma^2 \Delta$, where γ is the conductor of $\mathbb{Z}[\pi] \subset O_K$.
- We note ν the ℓ -adic valuation of γ , and ν_E the ℓ -adic valuation of the conductor γ_E of $O \subset O_K$.

Isogeny graph: horizontal isogenies

If $v = 0$, then every ℓ -isogeny is horizontal, and there are $1 + \frac{\Delta}{\ell}$ such isogeny. More precisely:

- 1 If ℓ splits in \mathcal{O} . In this case Δ_π is a non zero square mod ℓ , and the Frobenius acts on $E[\ell]$ as $\begin{pmatrix} \lambda & 0 \\ 0 & \mu \end{pmatrix}$ where the two eigenvalues λ and μ are distinct. The modular polynomial splits into irreducible factors of degree 1, 1, r, \dots, r where r is the order of $\lambda/\mu \in \mathbb{F}_\ell$. There are **2 horizontal isogenies**.
- 2 If ℓ is inert in \mathcal{O} . Then Δ_π is not a square modulo ℓ . The two eigenvalues λ and μ are conjugate in $\mathbb{F}_{\ell^2} \setminus \mathbb{F}_\ell$. The modular polynomial splits as irreducible factors of degree r , where r is the smallest number such that $\lambda^r \in \mathbb{F}_\ell$ (or equivalently such that π^r acts like a scalar on $E[\ell]$). There are **no horizontal isogenies**.
- 3 If ℓ is ramified in \mathcal{O} . Then $\Delta_\pi \equiv 0 \pmod{\ell}$. In this case π acts on $E[\ell]$ as $\begin{pmatrix} \lambda & 1 \\ 0 & \lambda \end{pmatrix}$. The modular polynomial splits into two irreducible factors of degree 1 and ℓ . There is **one horizontal isogeny**.



Isogeny graph: vertical isogenies

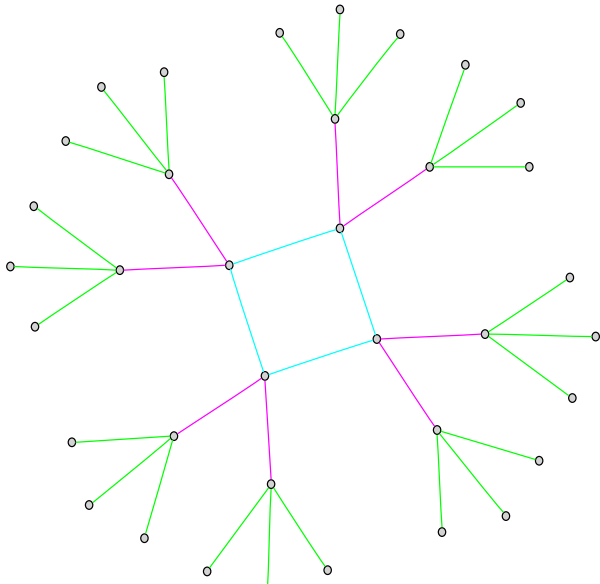
If $\nu \neq 0$. Then

- If $\nu_E = 0$, that is if $O_\ell = (O_K)_\ell$. There are $1 + \frac{\Delta}{\ell}$ horizontal isogenies, and $\ell - \frac{\Delta}{\ell}$ descending isogenies (that is $\ell - 1$, $\ell + 1$ or ℓ whether ℓ splits, is inert or is ramified in O_K).
- If $0 < \nu_E < \nu$, there is one ascending isogeny, and ℓ -descending ones.
- If $\nu_E = \nu$, that is $O_\ell = \mathbb{Z}[\pi]_\ell$, there is only one ascending isogeny.

In the first two cases, π acts as a scalar on $E[\ell]$ (and the modular polynomial splits completely), while in the last case π acts as $\begin{pmatrix} \lambda & 1 \\ 0 & \lambda \end{pmatrix}$ (and the modular polynomial splits into two irreducible factors of degree 1 and ℓ).

Isogeny graph: graphic interpretation of the local picture

The isogeny graph looks like a volcano [FM02]:

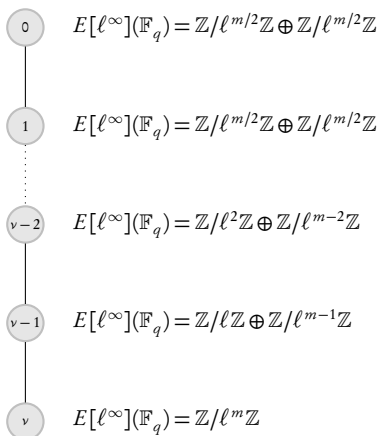


Isogeny graph: graphic interpretation of the local picture

- The volcano has height ν .
- The crater has length:
 - ① 0 if ℓ is inert;
 - ② 1 if ℓ splits;
 - ③ the order of l in the class group of the order of the curves in the crater when ℓ splits as \bar{l} .
- Taking an extension only increase the height of the volcano;
- If the height ν is non 0, then the only extension increasing the height are of degrees d with $\ell \mid d$.
- If $d = \ell$ the height increase only by one (except possibly when $\ell = 2$ and $\nu = 1$).

The structure of the ℓ^∞ -torsion in the volcano

- If E is on the floor, then $E[\ell^\infty](\mathbb{F}_q)$ is cyclic: $E[\ell^\infty](\mathbb{F}_q) = \mathbb{Z}/\ell^m\mathbb{Z}$ (possibly $m = 0$).
- If E is on level $\alpha < m/2$ above the floor, then $E[\ell^\infty](\mathbb{F}_q) = \mathbb{Z}/\ell^\alpha \oplus \mathbb{Z}/\ell^{m-\alpha}$.
- If E is on level $\alpha \geq m/2$, then m is even and $E[\ell^\infty](\mathbb{F}_q) = \mathbb{Z}/\ell^{m/2} \oplus \mathbb{Z}/\ell^{m/2}$.

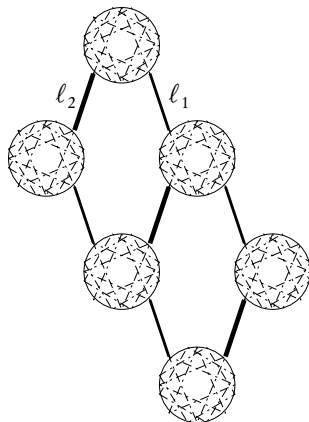
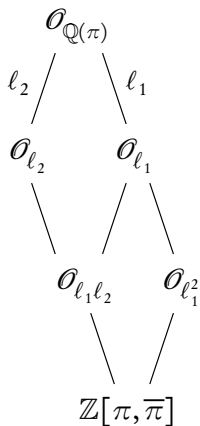


The global structure

Theorem (Complex multiplication)

Let E be an elliptic curve with endomorphism ring \mathcal{O} . Then the set of horizontal isogenies form a principal homogeneous space under the class group of \mathcal{O} .

This yields the following global picture (courtesy of Gaetan Bisson):



Finding the endomorphism ring

- Locally: for each $\ell \mid \gamma$, follow 3 paths in the ℓ -volcano. The first path reaching the floor give us the height of the curve in the volcano.
Since $\gamma \approx \sqrt{q}$, this is exponential.
- Globally, by using relations in the class groups of the orders. If R is a relation in $\text{Cl}(O)$ but the corresponding isogeny path is not cyclic then we know that $O \not\subseteq \text{End}(E)$. This give a subexponential algorithm (under GRH). More details will be given in Gaetan's talk next week.

Cryptographic applications of the endomorphism ring

- It is a finer grained invariant than the number of point.
- It gives an idea of “where we are” in the full isogeny graph.
- It is used by the CRT method to compute class polynomials: from a curve in the isogeny class, we want to find a curve with maximal endomorphism ring.
- The cycle in the crater can be used to compute $\chi_\pi \pmod{\ell^n}$.

Outline

- 1 Isogenies on elliptic curves
- 2 Endomorphisms
- 3 Supersingular elliptic curves**
- 4 Abelian varieties
- 5 References

Isogeny class of supersingular curves

Let $q = p^n$. The isogeny classes of elliptic curves are given by the value of the trace t by Tate's theorem. The possible value of t are:

- t prime to p , in this case the isogeny class is ordinary.
- The other cases give supersingular elliptic curves. The endomorphism fraction ring $\text{End}_k^0(\mathcal{E})$ of the isogeny class is either a quaternion algebra of rank 4, or an imaginary quadratic field. In the latter case, it will become maximal after an extension of degree d , with:
 - ① If n is even:
 - $t = \pm 2\sqrt{q}$, this is the only case where $\text{End}_k^0(\mathcal{E})$ is a quaternion algebra.
 - $t = \pm\sqrt{q}$ when $p \not\equiv 1 \pmod{3}$, here $d = 3$.
 - $t = 0$ when $p \not\equiv 1 \pmod{4}$, here $d = 2$.
 - ② If n is odd:
 - $t = 0$, here $d = 2$.
 - $t = \pm\sqrt{2q}$ when $p = 2$, here $d = 4$.
 - $t = \pm\sqrt{3q}$ when $p = 3$, here $d = 6$.

The commutative case

- If $K = \text{End}_k^0(E)$ is commutative, then χ_π is irreducible and $K = \mathbb{Q}(\pi)$. $\mathbb{Z}[\pi]$ is maximal for every $\ell \neq \{2, p\}$.
- The endomorphism rings of the isogeny class are the orders containing $\mathbb{Z}[\pi]$ maximal at p .
- If O is such an order, the class group $\text{Cl}(O)$ acts principally on the set of elliptic curves in the isogeny class with O as ring of endomorphisms.
- If k' is such that $\text{End}_{k'}^0(E)$ is maximal (i.e. a quaternion algebra), then it can happen that some curves E' in the isogeny class become isomorphic to E over k' .

The maximal case

- If $K = \text{End}_k^0(E)$ is non commutative, then it is the quaternion algebra ramified only at p and ∞ . The frobenius $\pi = p^{m/2} \in \mathbb{Z}$ and χ_π is a square. The endomorphism rings in the isogeny class corresponds to the maximal orders of K .
- If O is any maximal order of K , then the isogeny class of E (up to isomorphism) is of size $\#\text{Cl}(O)$. There is one or two curve in the isogeny class with endomorphism ring O , according to whether \mathfrak{p} is principal or not, where \mathfrak{p} is the ideal such that $\mathfrak{p}^2 = p$.
- If n is even there are two isogeny classes (quadratic twists of each other) with a maximal endomorphism ring.

Remark

Any two supersingular elliptic curves become isogenous after a quadratic extension of degree $2d$ (with d the degree where their endomorphism ring become maximal). But a new maximal class and up to 3 commutative classes appear in this extension.

Supersingular elliptic curves over $\overline{\mathbb{F}}_p$

- In characteristic p , every supersingular curve is defined over \mathbb{F}_{p^2} .
- For every $\ell \neq p$, the isogeny graph of supersingular curves (up to twists) over \mathbb{F}_{p^2} is connected. It has $p/12 + O(1)$ vertices, and diameter $O(\log p)$.
- The absolute endomorphism ring $\text{End}_{\overline{k}}(E)$ of a supersingular curve is a maximal order in the quaternion algebra ramified only at p and ∞ .
- There is a bijection between the set of such orders, and the set of supersingular elliptic curve (up to an action of $\text{Gal}(\mathbb{F}_{p^2}/\mathbb{F}_p)$).

Outline

- 1 Isogenies on elliptic curves
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- 4 Abelian varieties**
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Abelian varieties

Definition

- An **Abelian variety** is a complete connected group variety over a base field k . The group law is abelian.
- A (separable) **isogeny** is a finite surjective (separable) morphism between two Abelian varieties.

Example

- Abelian varieties of dimension 1 are elliptic curves.
- The Jacobian of a curve of genus g is an abelian variety of dimension g .

Non absolutely simple abelian varieties

Definition

- An abelian variety A_k is simple if the only subvariety of A_k are 0_{A_k} and itself.
- A_k is absolutely simple if it is simple over \bar{k} .

Even if an abelian variety A is ordinary, lot of funny things can happen if it is not absolutely simple:

- Not every non zero morphism is an isogeny.
- The endomorphism ring $\text{End}^0(A) = \text{End}(A) \otimes \mathbb{Q}$ may not be a division algebra.
- We can have $\text{End}_{k'}^0(A) \neq \text{End}_k^0(A)$ for extensions k' of k .
- A can be isogenous to another abelian variety A' , isomorphic to it over an extension of k , but not isomorphic to it over k .

Decomposing abelian varieties

Theorem (Poincaré-Weil)

Every abelian variety A is isogenous to a product of simple abelian varieties $A = \prod A_i^{m_i}$. The decomposition is entirely determined by χ_{π_A} .

- $\text{End}^0(A_i)$ is a division algebra.
- $\text{End}^0(A) = \prod M_{m_i}(\text{End}^0(A_i))$.

Theorem (Tate)

$\text{Hom}_k(A, B)$ is free of rank the number of common roots (with multiplicity) of χ_{π_A} and χ_{π_B} .

Endomorphism rings of abelian varieties

Let A be a simple abelian variety of dimension g . Then

- 1 $\chi_\pi = m_A^e$ where m_A is the minimal polynomial of the Frobenius and is irreducible.
- 2 $\text{End}^0(E)$ is a division algebra of center $\mathbb{Q}(\pi)$. The type of $\text{End}^0(E)$ is entirely determined by π .
- 3 We have $2g = de$, where d is the degree of m_A . $\text{End}^0(E)$ is of rank de^2 .

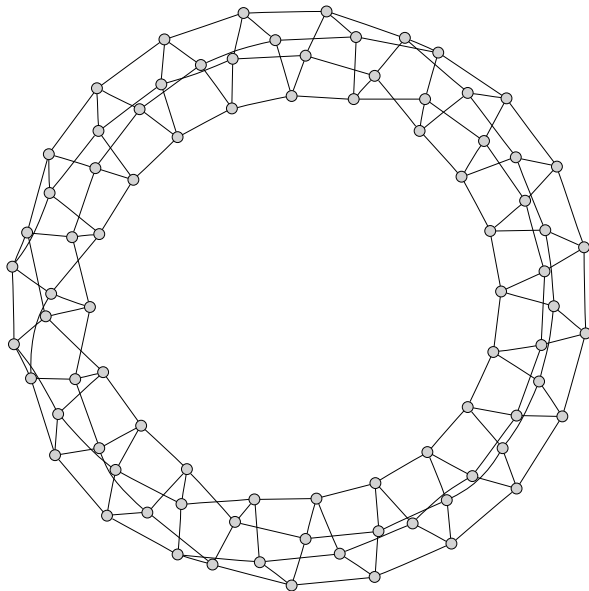
Remark

- If A is ordinary, then $e = 1$, χ_π is irreducible and $K = \text{End}_k^0(E)$ is a CM-field of rank $2g$.
- Moreover if A is absolutely simple, then $K = \mathbb{Q}(\pi) = \mathbb{Q}(\pi^n)$ for every n and $\text{End}_k(A) = \text{End}_{\overline{k}}(A)$.

Computing isogenies and endomorphisms

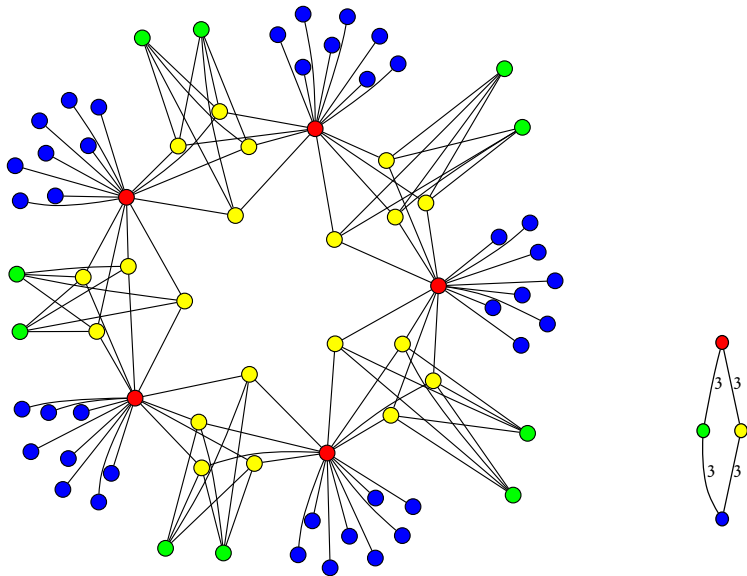
- In dimension 2, one can define modular polynomials using the Igusa invariants [Gau00; Dup06; BL09]. But these are too big to compute even for $\ell \geq 3$.
- We have an equivalent of Vélu's formula for maximally isotropic kernels [LR10; CR11].
- We also have subexponential algorithms to compute the endomorphism ring in dimension 2 [Bis11b].
- See the package AVIsogenies [BCR10] for an implementation of isogenies and endomorphism ring computation (mostly restricted to dimension 2 for now).

Isogeny graph in genus 2: example of horizontal isogenies

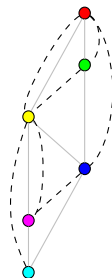
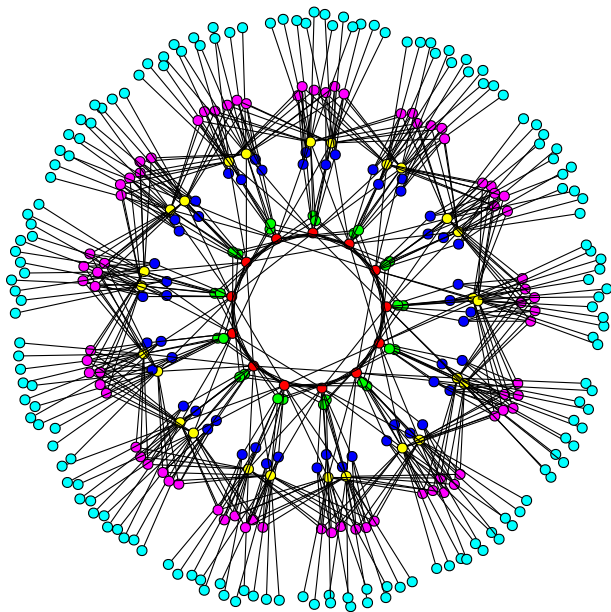


Isogeny graph in genus 2: vertical isogenies

Computations done by Gaetan Bisson using AVIsogenies.



Isogeny graph in genus 2: vertical isogenies



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- 1 Isogenies on elliptic curves
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Elliptic curves

- For a meta look at attacks on elliptic curves using isogenies to transfer the DLP: [KKM09, Section 11.2].
- Computing the modular polynomial: [Eng09a; BLS09].
- Different methods to compute class fields polynomials (the best known methods use the CRT and isogenies): [Eng09b; Sut09; ES10].
- Explicit isogenies in large characteristic: see [Elk92; Elk97]; and [BMS+08] for the best current known algorithm, with a nice history of previous methods.
- Explicit isogenies in small characteristic: [JL06; LS08] for methods based on lifting, [Cou94; Cou96] for Couveigne’s algorithm. The current best implementation of Couveigne’s algorithm is in [Feo10a], a nice summary is in [Feo10b].
- Some papers on SEA point counting algorithm [Sch95; Mor95; Elk97; Ler97].
- About isogenies and isomorphisms descending to the base field, see [Cox89, Proposition 14.19] and [Sch95, Proposition 6.1].
- See [Sil86, Chapter X, Theorem 2.2] for the equivalence between automorphisms and twists.
- An algorithm to compute endomorphism ring was developed in Kohel’s thesis [Koh96]. Some extensions to supersingular curves are in [ML04; Cer04].
- Developing the result of Kohel’s led to the notion of “isogeny volcano” [FM02] and improvements of the computation of the endomorphism ring [Fou01] with applications to the CRT method to compute class polynomials.
- Finally, a subexponential algorithm is developed in [BS09; Bis11a; Bis11b].
- One can also use the cycle given by the crater of the volcano to recover the trace of the Frobenius modulo a power of ℓ [CM94; CDM96; FM02; Fou01].

- Using pairings to go up in the Volcano [IJ10]. The ℓ^∞ -torsion in the volcano is described there, and also in [MMS+06].

Abelian varieties

- For an introduction to abelian variety, see [Mil91]. For more informations, see [Mum70], with [Mil85; Mil86] for simplified proofs using étale cohomology, and [GM07] for a more recent account. For abelian varieties over \mathbb{C} , see [Mum83; Mum84; Mum91] and a more recent account in [BL04].
- Some nice informations on abelian varieties over finite fields (Tate’s theorem, Honda-Tate theory) see [WM71] and [Wat69] for a more complete treatment.
- A description of ordinary abelian variety over a finite field is given by an equivalence of category [Del69], the link is further studied in [How95].
- For algebraic theta functions, see [Mum66; Mum67a; Mum67b], and some new results in [Kem89].
- Computing modular polynomials in genus 2: [Gau00; Dup06; BL09]. Computing a certain modular correspondance using theta functions [FLR11].
- Computing isogenies in abelian varieties using theta functions [LR10; CR11].
- For an introduction to the use of theta functions in cryptography (arithmetic, pairings, isogenies) see [Rob10].
- Computing endomorphism ring see [EL07; FL08; Wag09; Bis11b].

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